

ARM1020E™

Revision: r1p7

Technical Reference Manual

ARM®

ARM1020E

Technical Reference Manual

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Release Information

Change history

Date	Issue	Change
3 July 2001	A	First release
30 January 2002	B	Second release
02 December 2002	C	Third release
13 December 2002	D	Fourth release r1p6
30 June 2003	E	Fifth release r1p7. Defect fixes and corrections for r1p6 errata.

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Preface

This preface introduces the ARM1020E Revision r1p7 Technical Reference Manual. It contains the following sections:

- *About this document* on page xviii
- *Feedback* on page xxii.

About this document

This is the technical reference manual for the ARM1020E processor.

Intended audience

This document is written to help designers develop systems around the ARM1020E processor.

Using this document

This document is organized into the following chapters:

Chapter 1 *Introduction*

Read this chapter to learn about the components of the ARM10 processor.

Chapter 2 *Integer Core*

Read this chapter to learn how the integer core pipeline achieves a throughput approaching one instruction per cycle.

Chapter 3 *Programmer's Model*

Read this chapter for details of the programmer's model and a summary of the ARM and Thumb instruction sets.

Chapter 4 *System Control Coprocessor*

Read this chapter to learn how to use CP15, the system control coprocessor.

Chapter 5 *Memory Management Units*

Read this chapter to learn how to use the address translation process of the memory management units.

Chapter 6 *Caches and Write Buffer*

Read this chapter to learn how to use CP15 to control operation of the instruction and data caches, the write buffer, and the hit-under-miss buffer.

Chapter 7 *Prefetch Unit*

Read this chapter to learn how the ARM10 processor prefetches and buffers instructions, and how to implement an instruction memory barrier to flush the prefetch buffer.

Chapter 8 *Bus Interface*

Read this chapter to learn how to use the bus interface to AMBA™.

Chapter 9 *Coprocessor Interface*

Read this chapter to learn how to integrate one or more coprocessors with the ARM10 processor.

Chapter 10 *JTAG Interface*

Read this chapter to learn how to use the built-in JTAG debug hardware.

Chapter 11 *Debug*

Read this chapter to learn how to use coprocessor 14 to debug application software, operating systems, and hardware systems.

Chapter 12 *Instruction Cycle Summary and Interlocks*

Read this chapter to learn about the cycle counts of ARM and Thumb instructions and how pipeline interlocks resolve data dependencies.

Chapter 13 *Design for Test*

Read this chapter to learn how to use the built-in scan chains, wrapper cells, and memory BIST to test the ARM10 processor.

Chapter 14 *Power Manager*

Read this chapter to learn how to use the power manager to control system power modes.

Chapter 15 *Clock Generator*

Read this chapter to learn how to synthesize the two programmable system clocks.

Appendix A *Signal Descriptions*

Refer to this appendix for a summary of ARM10 signal descriptions.

Product revision status

The *mpn* identifier indicates the revision status of the product described in this document, where:

- rn*** Identifies the major revision of the product.
- pn*** Identifies the minor revision or modification status of the product.

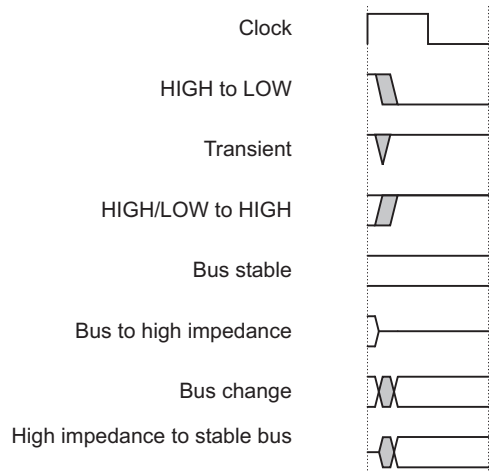
Typographical conventions

The following typographical conventions are used in this book:

<i>italic</i>	Introduces special terminology. Also denotes cross-references.
bold	Denotes signal names. Also used for terms in descriptive lists, where appropriate.
monospace	Denotes text that can be entered at the keyboard, such as commands, file and program names, and source code.
<u>monospace</u>	Denotes a permitted abbreviation for a command or option. The underlined text can be entered instead of the full command or option name.
<i>monospace italic</i>	Denotes arguments to commands and functions where the argument is to be replaced by a specific value.
monospace bold	Denotes language keywords when used outside example code.

Timing diagram conventions

This manual contains one or more timing diagrams. The following key explains the components used in these diagrams. Any variations are clearly labeled when they occur. Therefore, no additional meaning must be attached unless specifically stated.



Key to timing diagram conventions

Shaded bus and signal areas are undefined, so the bus or signal can assume any value within the shaded area at that time. The actual level is unimportant and does not affect normal operation.

Register Notation Conventions

The table shows the terms and abbreviations used in register descriptions. In all cases, reading or writing any fields, including those specified as Unpredictable, Should Be One, or Should Be Zero, does not cause any physical damage to the chip.

Register notation conventions

Term	Description
Unpredictable (UNP)	Data read from this field can have any value. Writing to this field causes unpredictable behavior or an unpredictable change in device configuration.
Undefined (UND)	An instruction that accesses this field in the manner indicated takes the undefined instruction trap.
Should Be Zero (SBZ)	When writing to this field, write only zeros. Writing ones has Unpredictable results.
Should Be One (SBO)	When writing to this field, write only ones. Writing zeros has Unpredictable results.

Further reading

This section lists publications by ARM Limited and by third parties.

ARM periodically provides updates and corrections to its documentation. See <http://www.arm.com> for current errata sheets, addenda, and Frequently Asked Questions.

ARM publications

This document contains information that is specific to the ARM1020E processor. Refer to the following documents for other relevant information:

- *ARM Architecture Reference Manual* (ARM DDI 0100)
- *ARM AMBA Specification* (ARM IHI 0001)
- *ARM10220E Test Chip Implementation Guide* (ARM DXI 0141)
- *VFP10 Technical Reference Manual* (ARM DDI 0106)
- *ETM10 Technical Reference Manual* (ARM DDI 0206).

Other publications

This section lists relevant documents published by third parties:

- IEEE Standard, *Test Access Port and Boundary-Scan Architecture specification* 1149.1-1990 (JTAG).

Feedback

ARM Limited welcomes feedback both on the ARM1020E processor, and on the documentation.

Feedback on the ARM1020E processor

If you have any comments or suggestions about this product, contact your supplier giving:

- the product name
- a concise explanation of your comments.

Feedback on this document

If you have any comments on this document, send email to errata@arm.com giving:

- the document title
- the document number
- the page number(s) to which your comments refer
- a concise explanation of your comments.

General suggestions for additions and improvements are also welcome.

Chapter 1

Introduction

This chapter describes the components and features of the ARM1020E r1p7 processor. It contains the following sections:

- *About the processor* on page 1-2
- *Components of the processor* on page 1-4
- *Silicon revision information* on page 1-9.

1.1 About the processor

The ARM1020E processor incorporates the ARM10E™ integer core, which implements the ARMv5TE architecture. It is a high-performance, low-power, cached processor that provides full virtual memory capabilities. It is designed to run high-end embedded applications and sophisticated operating systems such as JavaOS, Linux, Microsoft WindowsCE, NetBSD, and EPOC-32 from Symbian. It supports the ARM and Thumb instruction sets, and includes EmbeddedICE-RT™ logic and JTAG software debug features.

The ARM1020E processor consists of:

- the ARM10E integer core:
 - load/store unit
 - prefetch unit
 - integer unit
 - EmbeddedICE-RT logic for JTAG-based debug
- external coprocessor interface and coprocessors CP14 and CP15
- *Memory Management Unit* (MMU)
- instruction and data caches
- write-back *Physical Address* (PA) TAG RAM
- write buffer and *Hit-Under-Miss* (HUM) buffer
- *Advanced Micro Bus Architecture* (AMBA) *High-performance Bus* (AHB) bus interface
- *Embedded Trace Macrocell* (ETM) interface.

Features of the ARM1020E processor include:

- a six-stage pipeline
- branch prediction that supports branch folding (zero cycle branches)
- 64KB level 1 cache (32KB instruction, 32KB data)
- full 64-bit interfaces between the integer core and caches, write buffer, and bus interface units on both instruction and data sides, and coprocessors
- multilayer AHB support through independent 64-bit AHB interfaces for instruction and data sides

- parallel execution of data processing instructions under load and store multiple instructions
- a HUM buffer that supports execution of load hits underneath an outstanding load miss
- nonblocking caches that support execution of data processing instructions under load misses
- additional register read and write ports to support reading of up to four registers and writing of three registers in one cycle
- improved power management support
- enhanced debug support.

1.2 Components of the processor

This section introduces the main blocks of the ARM1020E processor and gives references to detailed descriptions of those blocks:

- *Integer core* on page 1-6
- *Memory Management Unit* on page 1-6
- *Instruction and data caches* on page 1-6
- *Cache power-down capabilities* on page 1-7
- *Branch prediction and prefetch unit* on page 1-7
- *AMBA interface* on page 1-7
- *Coprocessor interface* on page 1-7
- *Debug* on page 1-7
- *Instruction cycle summary and interlocks* on page 1-7
- *Design-for-test features* on page 1-7
- *Power management* on page 1-8
- *Clocking and PLL* on page 1-8.

Figure 1-1 on page 1-5 shows the main blocks of the ARM1020E processor.

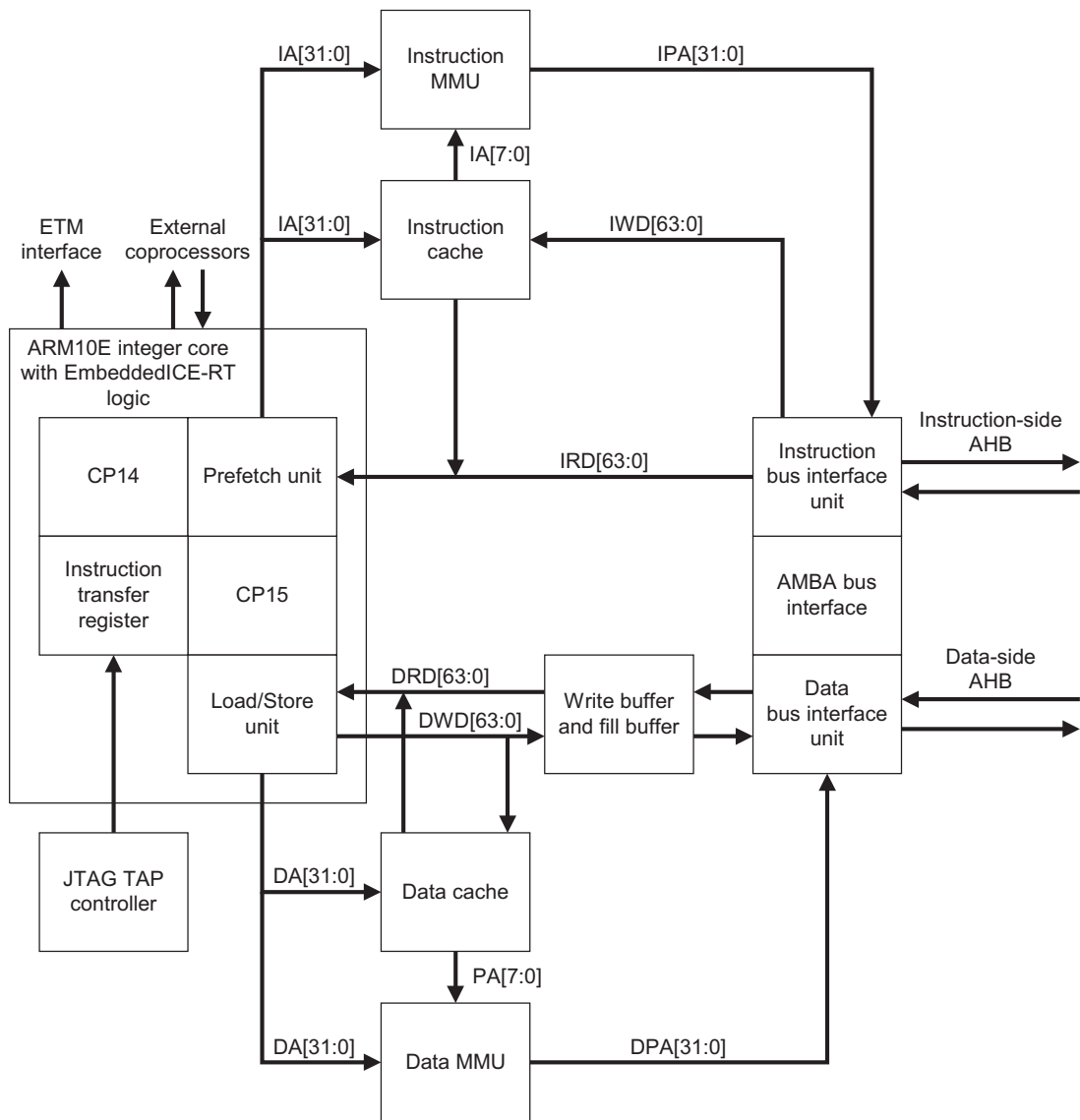


Figure 1-1 ARM1020E processor block diagram

1.2.1 Integer core

This ARM1020E processor is built around the ARM10E integer core in an ARMv5TE implementation that runs the 32-bit ARM and 16-bit compressed Thumb instruction sets. You can balance high performance against code size and extract maximum performance from 8-bit, 16-bit, and 32-bit memory. The processor includes EmbeddedICE-RT logic for JTAG software debugging, and is supported by the Multi-ICE JTAG debug interface.

Refer to Chapter 2 *Integer Core* for details of the pipeline stages and instruction progression.

Refer to Chapter 4 *System Control Coprocessor* for system coprocessor programming information.

1.2.2 Memory Management Unit

The MMU has separate instruction and data *Translation Lookaside Buffers* (TLBs). It is backward-compatible with the ARM v4 architecture MMU of StrongARM and ARM920T. The MMU includes a 1KB tiny page mapping size to enable a smaller RAM and ROM footprint for embedded systems and operating systems such as WindowsCE™ that have many small mapped objects. The ARM1020E processor implements the *Fast Context Switching Extension* (FCSE) and high vectors extension that are required to run Microsoft WindowsCE. Refer to Chapter 5 *Memory Management Units* for more information.

1.2.3 Instruction and data caches

This ARM1020E processor has a 32KB *Instruction Cache* (ICache) and a 32KB *Data Cache* (DCache). The data cache provides *Write-Through* (WT) or *Write-Back* (WB) operation, selected under software control on a per-region basis. The large caches enable you to obtain high performance from commodity memory systems by significantly reducing:

- the read bandwidth required of main memory
- the write bandwidth required of main memory (when write-back caching is used)
- overall system power consumption by reducing accesses to off-chip memory.

The processor provides a write buffer that holds up to eight 64-bit values, each at an independent address.

Refer to Chapter 6 *Caches and Write Buffer* for more information.

1.2.4 Cache power-down capabilities

The power manager provides a software-controlled hardware mechanism to maintain power to the CAM and RAM state element arrays in the caches when the remainder of the device is powered down. Refer to Chapter 6 *Caches and Write Buffer* for more information.

1.2.5 Branch prediction and prefetch unit

The prefetch unit is part of the integer core. It fetches instructions from the ICache or from external memory and issues them to the integer core. To increase performance, it also predicts the outcome of branches in the instruction stream. Refer to Chapter 7 *Prefetch Unit* for more information.

1.2.6 AMBA interface

The bus interface unit provides a multimaster AHB interface to memory and peripherals. The AHB is an on-chip bus with two unidirectional 64-bit data buses and one 32-bit address bus. Refer to Chapter 8 *Bus Interface* for more information.

1.2.7 Coprocessor interface

Chapter 9 *Coprocessor Interface* describes the interface for on-chip coprocessors such as floating-point units or application-specific hardware acceleration units.

1.2.8 Debug

The debug coprocessor, CP14, implements a full range of debug features described in Chapter 10 *JTAG Interface* and Chapter 11 *Debug*.

1.2.9 Instruction cycle summary and interlocks

Chapter 12 *Instruction Cycle Summary and Interlocks* describes instruction cycle times and gives examples of interlock timing.

1.2.10 Design-for-test features

The ARM1020E processor is designed to be embedded into large *System-on-Chip* (SoC) designs. The EmbeddedICE-RT logic debug facilities, AMBA on-chip system bus, and test methodology are all designed for efficient use of the processor when integrated into a larger IC. Refer to Chapter 13 *Design for Test* for details of testing.

1.2.11 Power management

Power management features are described in Chapter 14 *Power Manager*.

1.2.12 Clocking and PLL

The ARM10 processor has two clock inputs:

- **GCLK**
- **HCLK**.

The design is fully static. When both these clocks are stopped, the internal state of the processor is preserved indefinitely. **GCLK** drives the internal logic in the processor. **HCLK** drives the bus interface. Most input and output timings are specified with respect to **HCLK**.

Refer to Chapter 15 *Clock Generator* and Chapter 8 *Bus Interface* for details.

———— **Note** —————

Typically, **GCLK** frequency is higher than that of **HCLK**. The two clocks must have a fixed phase relationship. **HCLK** is usually derived by dividing down the source of **GCLK**.

—————

1.2.13 ETM interface logic

An optional external ETM can be connected to the ARM1020E processor to provide real-time tracing of instructions and data in an embedded system. The processor includes the logic and interface to enable you to trace program execution and data transfers using the ETM10. Further details are in *Embedded Trace Macrocell Specification*. See Table A-11 on page A-14 for descriptions of ETM-related signals.

1.3 Silicon revision information

This manual is for revision r1p7 of the ARM1020E processor, see *Product revision status* on page xix for details of revision numbering. Revision r1p7 of the ARM1020E processor contains fixes for revision r1p6 errata.

Chapter 2

Integer Core

This chapter describes the ARM10 integer core. It contains the following sections:

- *About the integer core* on page 2-2
- *Pipeline* on page 2-4
- *Prefetch unit* on page 2-5
- *Typical operations* on page 2-6
- *Load/store unit* on page 2-8
- *Instruction progression* on page 2-9.

2.1 About the integer core

By overlapping the various stages of operation, the integer core maximizes the clock rate achievable to execute each instruction. Because it has multiple execution units, the integer core enables multiple instructions to exist in the same pipeline stage, enabling simultaneous execution of some instructions. As a result, it delivers a peak throughput approaching one instruction per cycle. The integer core consists of:

Prefetch unit

The prefetch unit fetches instructions from instruction cache or external memory. To reduce the number of pipeline refills, it predicts the outcome of branches whenever it can.

Integer unit

The integer unit decodes instructions sent from the prefetch unit. It contains the barrel shifter, ALU, and multiplier, and executes dataprocessing instructions such as MOV, ADD, and MUL. The integer unit helps the load/store unit to execute loads, stores, and coprocessor transfer instructions such as LDR, STM, LDC, and MCRR. It also contains the main instruction sequencer that takes care of multicycle data processing instructions, mode changes, exceptions, and debug events.

Load/store unit

The *Load/Store Unit* (LSU) can load or store two registers (64 bits) per cycle, if the data address is 64-bit aligned. After the first access of a load or store multiple instruction (LDM or STM) the LSU can decouple from the integer unit and complete the instruction autonomously.

While the LSU is decoupled, the integer unit can run data processing instructions if there are no dependencies on the LSU or on the loaded or stored data.

The LSU also supports *Hit-Under-Miss* (HUM) operation. If a load misses in the data cache, the outstanding request is moved into the HUM buffer. Other instructions, including loads, can continue to execute unless a second miss occurs or a dependency on the outstanding data is detected.

These components are shown in Figure 2-1 on page 2-3.

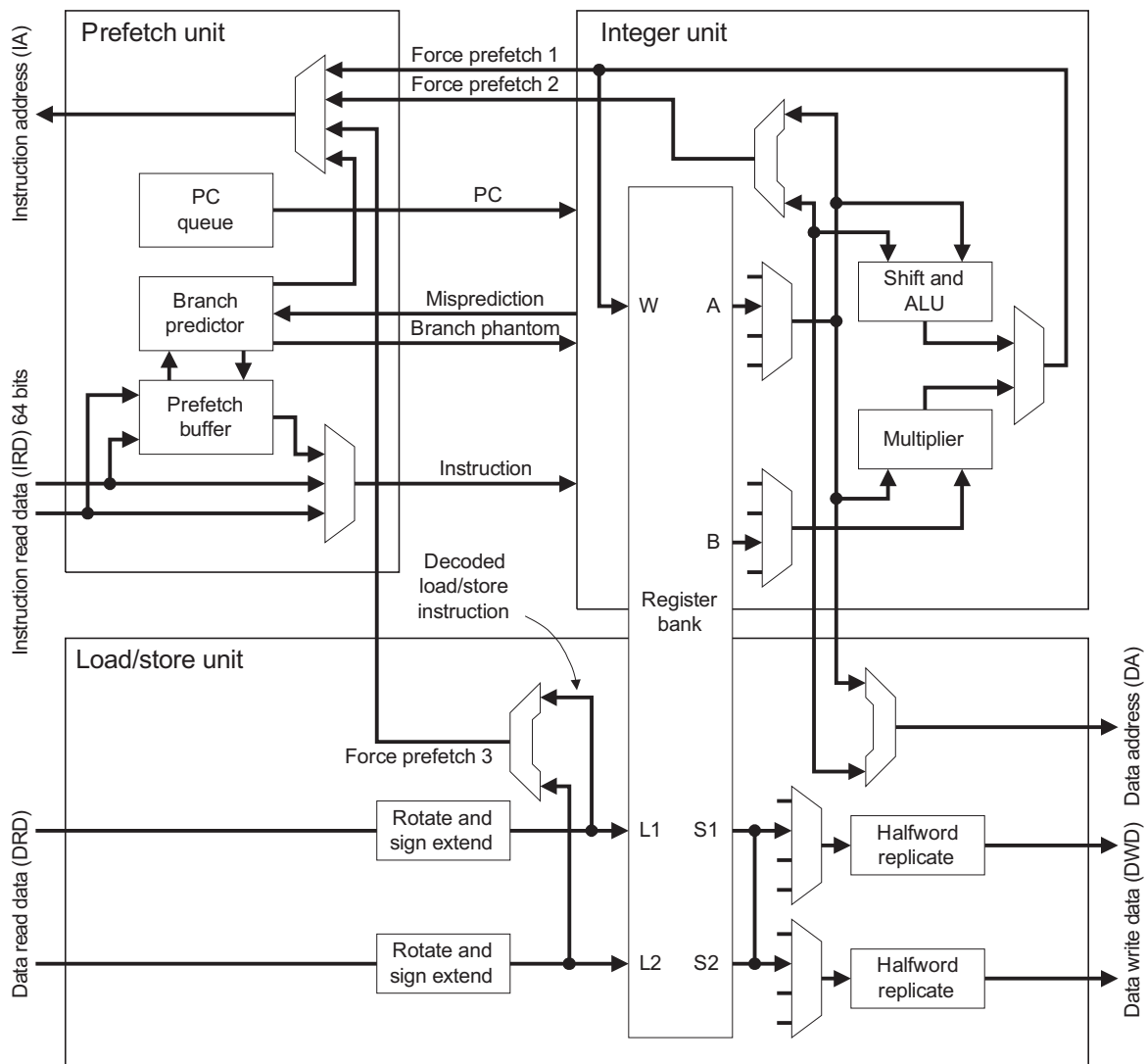


Figure 2-1 Integer core components

2.2 Pipeline

The ARM10 pipeline consists of six stages to maximize instruction throughput:

Fetch	Instruction cache access. Branch prediction for instructions that have already been fetched.
Issue	Initial instruction decode.
Decode	Final instruction decode, register reads for <i>Arithmetic/Logic Unit (ALU)</i> operation, forwarding, and initial interlock resolution.
Execute	Data access address calculation, data processing shift, shift and saturate, ALU operation, first stage of multiplications, flag setting, condition code check, branch mispredict detection, and store data register read.
Memory	Data cache access, second stage of multiplications, and saturations.
Write	Register writes, instruction retirement.

The Fetch stage uses a *First-In-First-Out buffer (FIFO)* prefetch buffer that can hold up to three instructions. Here a path to fetch along is predicted ahead of execution of branch instructions.

The Issue and Decode stages can contain a predicted branch in parallel with one instruction.

The Execute, Memory, and Write stages can simultaneously contain all of the following:

- a predicted branch
- an ALU or multiply instruction
- ongoing multicycle load or store multiple instructions
- ongoing multicycle coprocessor instructions.

2.3 Prefetch unit

The prefetch unit and branch prediction are described in detail in Chapter 7 *Prefetch Unit*.

The prefetch unit operates in the Fetch stage of the pipeline. It can fetch 64 bits every cycle from the instruction-side cache. It can only issue one 32-bit instruction per cycle to the integer unit. Because it can fetch more instructions than it can issue, the prefetch unit puts pending instructions in the prefetch buffer. While an instruction is in the prefetch buffer, the branch prediction logic can decode it to see if it is a predictable branch.

Where possible, the branch prediction logic removes branches from the instruction stream. If the branch is predicted to be taken, then the instruction address is redirected to the branch target address. If the branch is predicted not to be taken, then the instruction address continues to progress through the instructions following the branch instruction. Often in these cases, if the instruction following the branch is already in the prefetch buffer, it can be issued in place of the branch and the branch effectively takes no cycles. When there is not enough time to completely remove the branch, the fetch address is redirected anyway, because this still helps to reduce the branch penalty.

The integer unit executes unpredicted or unpredictable branches. To get the address out quickly, it uses a dedicated fast branch adder whose inputs do not pass through the barrel shifter.

A multiplexor in the LSU sends loaded data straight to the prefetch unit. This updates the fetch address after loads to the *Program Counter* (PC).

There is also a path from the ALU output to the prefetch unit. This is used for data processing instructions that write to the PC. Because the path through the barrel shifter and ALU is slower than that through the dedicated adder, these instructions usually take one more cycle than branches. The one exception is a simple move that does not require a shift, for example, `MOV PC R14`. For optimum performance, this uses the fast branch adder rather than the ALU.

2.4 Typical operations

Figure 2-2 shows in the six stages of a typical operation.

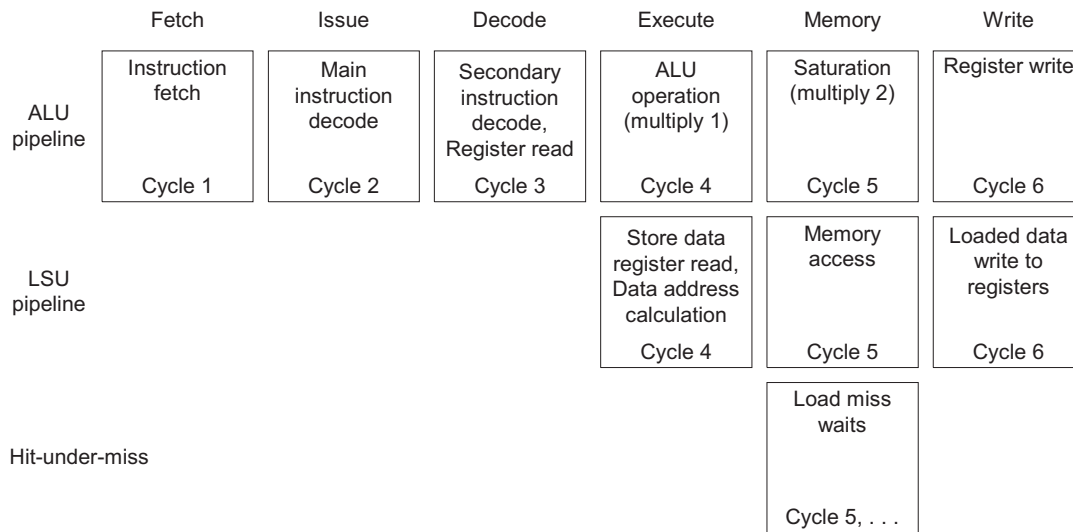


Figure 2-2 Pipeline stages of a typical operation

Figure 2-3 on page 2-7 shows the stages of a typical data processing operation.

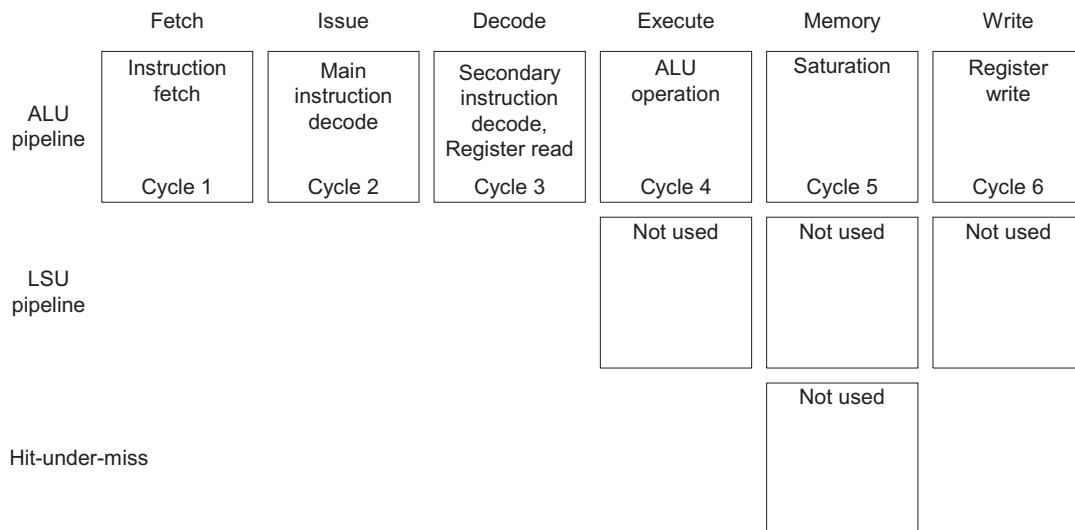


Figure 2-3 Pipeline stages of a typical ALU operation

Figure 2-4 shows the stages of a typical multiply operation. The MUL loops in the Execute stage until it passes through the first part of the multiplier array enough times. Then it progresses to the Memory stage where it passes once through the second half of the array to produce the final result.

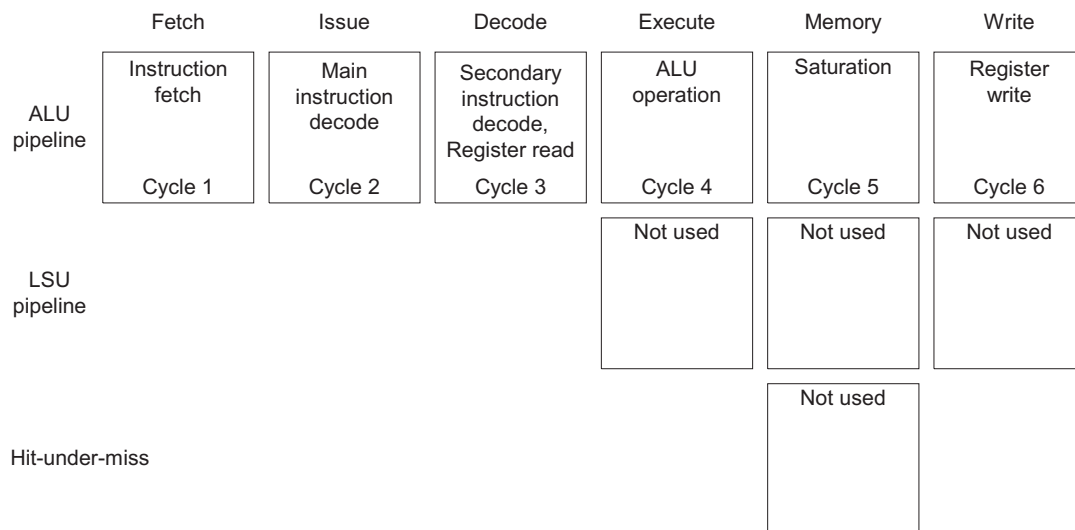


Figure 2-4 Pipeline stages of a typical multiply operation

2.5 Load/store unit

If the data address is 64-bit aligned, the LSU can load or store two registers (64 bits) per transfer. This does not speed up single load or store instructions (LDR, STR) but it does considerably speed up load and store multiple instructions (LDM, STM). Load and store double instructions (LDRD, STRD) also take advantage of the available bandwidth.

Accesses that are not 64-bit aligned have to take place over two cycles. If an LDM or STM address is not 64-bit aligned, then only one register (32 bits) is transferred on the first access. After that, two registers per cycle can be transferred each cycle.

Single loads and stores work in cooperation with the integer unit. The first cycle of multiple loads and stores works in cooperation with the integer unit, but the LSU can finish ongoing multiple loads and stores autonomously.

The LSU calculates the address for the data access using a dedicated adder. This adder evaluates in parallel with the adder in the ALU. The adder in the ALU calculates a base register write-back value if it is required.

The A and B register ports of the integer unit read the operands for both adders. For complex (scaled-register) addressing modes that require the barrel shifter, the ALU has to calculate data addresses. This costs one extra cycle.

The LSU has two dedicated register bank read ports (S1 and S2) and two dedicated write ports (L1 and L2). These are used to read data to be stored and to write data that is loaded.

2.6 Instruction progression

Figure 2-5 shows a simple LDR/STR operation that hits in the data cache.

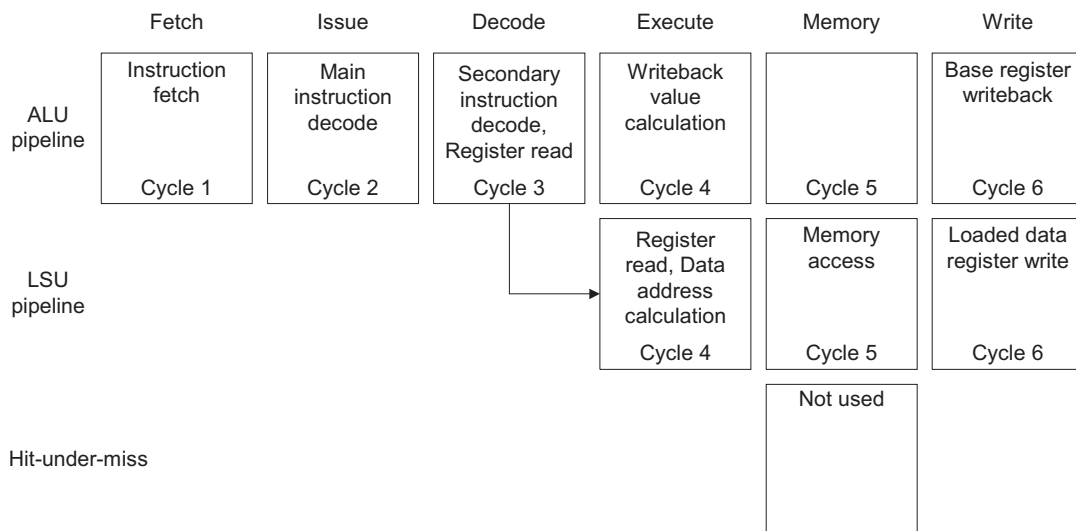


Figure 2-5 Pipeline stages of a load or store operation

Figure 2-6 on page 2-10 shows the progression of an LDM/STM operation using the load/store pipeline to complete. Other instructions can use the ALU pipeline at the same time as the LDM/STM completes in the LSU pipeline.

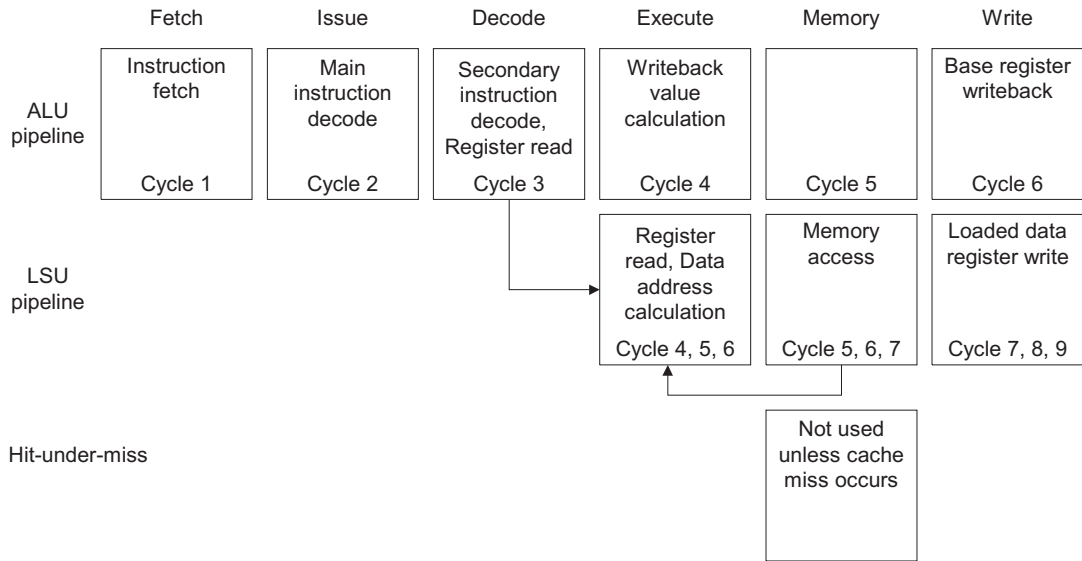


Figure 2-6 Pipeline stages of a load multiple or store multiple operation

Figure 2-7 shows the progression of an LDR that misses. When the LDR is in the HUM stage, other instructions, including independent loads that hit in the cache, can run under it.

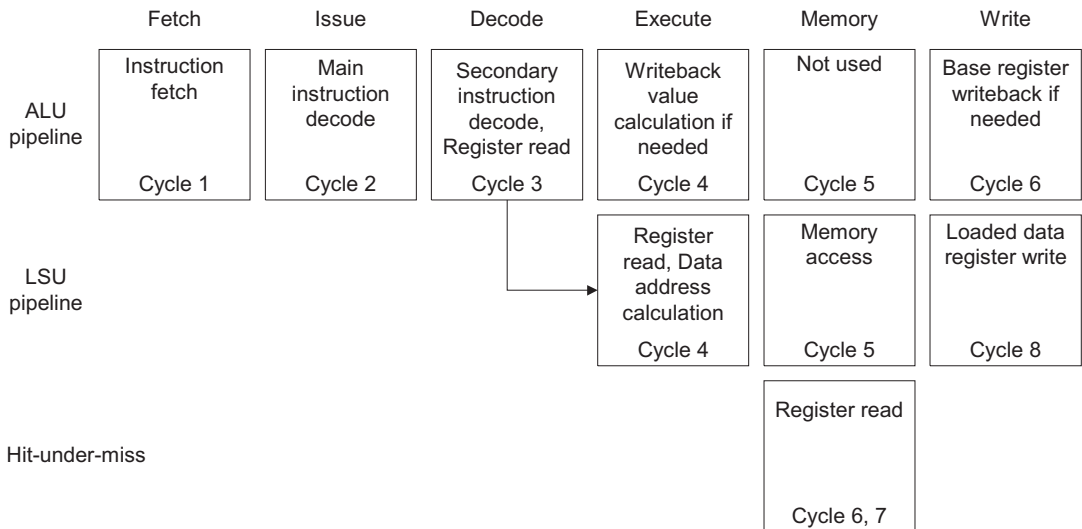


Figure 2-7 Pipeline stages of an LDR operation that misses

Refer to Chapter 12 *Instruction Cycle Summary and Interlocks* for further details of instruction cycles and timings.

Chapter 3

Programmer's Model

This chapter describes the ARM10E programmer's model and provides a summary of the ARM and Thumb instruction sets. It contains the following sections:

- *About the programmer's model* on page 3-2
- *Instruction set summary* on page 3-3.

3.1 About the programmer's model

The ARM10E programmer's model, including a detailed instruction set specification, is described in the *ARM Architecture Reference Manual*. The programmer's model of the ARM1020E processor is the same as the programmer's model of the ARM10E core, but extended in the following ways:

- The system control coprocessor (CP15) is integrated into the ARM10 processor and provides additional registers for configuring and controlling caches, MMU, protection system, power-down, and clocking mode.
- The MMU page tables define the virtual-to-physical address mapping, page and section access permissions, cache, and write buffer configuration. These are created by the operating system software and accessed automatically by the MMU hardware whenever an instruction read or data access causes a TLB miss.

3.2 Instruction set summary

This section provides a summary of the ARM and Thumb instruction sets:

- *ARM instruction summary* on page 3-4
- *Thumb instruction summary* on page 3-12.

The ARM1020E processor is an implementation of the ARM architecture version 5TE. For a complete description of both instruction sets, refer to the *ARM Architecture Reference Manual*.

Table 3-1 is a key to the notation used in the instruction set tables.

Table 3-1 Key to instruction set table notation

Notation	Description
{cond}	Table 3-11 on page 3-11 defines the condition notation.
<Oprnd2>	Table 3-9 on page 3-10 gives examples of Oprnd2.
{field}	Table 3-10 on page 3-11 defines the field notation.
S	Sets condition codes (optional).
B	Byte operation (optional).
H	Halfword operation (optional).
T	Forces address translation. Cannot be used with preindexed addresses.
<a_mode2>	Table 3-3 on page 3-7 describes addressing mode 2.
<a_mode2P>	Table 3-4 on page 3-8 describes addressing mode 2 (privileged).
<a_mode3>	Table 3-5 on page 3-9 describes addressing mode 3.
<a_mode4L>	Table 3-6 on page 3-9 describes addressing mode 4 (load).
<a_mode4S>	Table 3-7 on page 3-10 describes addressing mode 4 (store).
<a_mode5>	Table 3-8 on page 3-10 describes addressing mode 5.
#32bit_Imm	A 32-bit constant, formed by right-rotating an 8-bit value by an even number of bits.
<reglist>	A comma-separated list of registers, enclosed in braces ({ and }).

3.2.1 ARM instruction summary

Table 3-2 summarizes the ARM instructions. Asterisks in the Operation column denote ARMv5TE instructions.

Table 3-2 ARM instruction summary

Operation		Assembler
Move	Move	MOV{cond}{S} Rd, <Oprnd2>
	Move NOT	MVN{cond}{S} Rd, <Oprnd2>
	Move SPSR to register	MRS{cond} Rd, SPSR
	Move CPSR to register	MRS{cond} Rd, CPSR
	Move register to SPSR	MSR{cond} SPSR_{field}, Rm
	Move register to CPSR	MSR{cond} CPSR_{field}, Rm
	Move immediate to SPSR flags	MSR{cond} SPSR_f, #32bit_Imm
	Move immediate to CPSR flags	MSR{cond} CPSR_f, #32bit_Imm
Arithmetic	Add	ADD{cond}{S} Rd, Rn, <Oprnd2>
	Add with carry	ADC{cond}{S} Rd, Rn, <Oprnd2>
	* Saturating add	QADD{cond} Rd, Rm, Rn
	* Saturating add	QDADD{cond} Rd, Rm, Rn
	Subtract	SUB{cond}{S} Rd, Rn, <Oprnd2>
	* Saturating subtract	QSUB{cond} Rd, Rm, Rn
	* Saturating subtract	QDSUB{cond} Rd, Rm, Rn
	Subtract with carry	SBC{cond}{S} Rd, Rn, <Oprnd2>
	Subtract reverse subtract	RSB{cond}{S} Rd, Rn, <Oprnd2>
	Subtract reverse subtract with carry	RSC{cond}{S} Rd, Rn, <Oprnd2>
	Multiply	MUL{cond}{S} Rd, Rm, Rs
	Multiply accumulate	MLA{cond}{S} Rd, Rm, Rs, Rn
	Multiply unsigned long	UMULL{cond}{S} RdLo, RdHi, Rm, Rs
	Multiply unsigned accumulate long	UMLAL{cond}{S} RdLo, RdHi, Rm, Rs

Table 3-2 ARM instruction summary (continued)

Operation	Assembler	
	Multiply signed long	SMULL{cond}{S} RdLo, RdHi, Rm, Rs
*	Multiply signed, 16-bit operands	SMUL<x><y>{cond}Rd, Rm, Rs, Rn
	Multiply signed, Word and 16-bit operand	SMULW<y>{cond}Rd, Rm, Rs, Rn
*	Multiply signed accumulate, 16-bit operands	SMLA<x><y>{cond} Rd, Rs, Rm, Rn
	Multiply signed accumulate long	SMLAL{cond}{S} RdLo, RdHi, Rm, Rs
*	Multiply signed accumulate long, 16-bit operands	SMLAL<x><y>{cond}{S} RdLo, RdHi, Rm, Rs
*	Multiply signed accumulate, Word and 16-bit operand	SMLAW<y>{cond} Rd, Rs, Rm, Rn
	Compare	CMP{cond} Rd, <Oprnd2>
	Compare negative	CMN{cond} Rd, <Oprnd2>
Logical	Test	TST{cond} Rn, <Oprnd2>
	Test equivalence	TEQ{cond} Rn, <Oprnd2>
	AND	AND{cond}{S} Rd, Rn, <Oprnd2>
	EOR	EOR{cond}{S} Rd, Rn, <Oprnd2>
	ORR	ORR{cond}{S} Rd, Rn, <Oprnd2>
	Bit clear	BIC{cond}{S} Rd, Rn, <Oprnd2>
Branch	Branch	B{cond} label
	Branch with link	BL{cond} label
	Branch and exchange instruction set	BX{cond} Rn
Load	Word	LDR{cond} Rd, <a_mode2>
	Word with user-mode privilege	LDR{cond}T Rd, <a_mode2P>
	Byte	LDR{cond}B Rd, <a_mode2>
	Byte with user-mode privilege	LDR{cond}BT Rd, <a_mode2P>
	Byte signed	LDR{cond}SB Rd, <a_mode3>
	Halfword	LDR{cond}H Rd, <a_mode3>
	Halfword signed	LDR{cond}SH Rd, <a_mode3>

Table 3-2 ARM instruction summary (continued)

Operation	Assembler	
*	Pair of registers	LDR{cond}D Rd, <a_mode3>
Load multiple	Increment before	LDM{cond}IB Rd{!}, <reglist>{^}
	Increment after	LDM{cond}IA Rd{!}, <reglist>{^}
	Decrement before	LDM{cond}DB Rd{!}, <reglist>{^}
	Decrement after	LDM{cond}DA Rd{!}, <reglist>{^}
	Stack operations	LDM{cond}<a_mode4L> Rd{!}, <reglist>
	Stack operations and restore CPSR	LDM{cond}<a_mode4L> Rd{!}, <reglist+pc>^
	User registers	LDM{cond}<a_mode4L> Rd{!}, <reglist>^
Store	Word	STR{cond} Rd, <a_mode2>
	Word with User mode privilege	STR{cond}T Rd, <a_mode2P>
	Byte	STR{cond}B Rd, <a_mode2>
	Byte with User mode privilege	STR{cond}BT Rd, <a_mode2P>
	Halfword	STR{cond}H Rd, <a_mode3>
*	Pair of registers	STR{cond}D Rd, <a_mode3>
Store multiple	Increment before	STM{cond}IB Rd{!}, <reglist>{^}
	Increment after	STM{cond}IA Rd{!}, <reglist>{^}
	Decrement before	STM{cond}DB Rd{!}, <reglist>{^}
	Decrement after	STM{cond}DA Rd{!}, <reglist>{^}
	Stack operations	STM{cond}<a_mode4S> Rd{!}, <reglist>
	User registers	STM{cond}<a_mode4S> Rd{!}, <reglist>^
Swap	Word	SWP{cond} Rd, Rm, [Rn]
	Byte	SWP{cond}B Rd, Rm, [Rn]
CP operations	Data operations	CDP{cond} p<cpnum>, <op1>, CRd, CRn, CRm, <op2>
	Move to ARM register from coprocessor	MRC{cond} p<cpnum>, <op1>, Rd, CRn, CRm, <op2>
	Move to coprocessor from ARM register	MCR{cond} p<cpnum>, <op1>, Rd, CRn, CRm, <op2>

Table 3-2 ARM instruction summary (continued)

Operation	Assembler
* Move two coprocessor registers into ARM registers	MRRC{cond} <coproc>, <opcode>, Rd, <Rm>, <CRm>
* Move two ARM registers into coprocessor registers	MCCR{cond} <coproc>, <opcode>, <Rd>, <Rn>, <CRm>
Load	LDC{cond} p<cpnum>, CRd, <a_mode5>
Store	STC{cond} p<cpnum>, CRd, <a_mode5>
Software interrupt	SWI 24bit_Imm
*Soft preload	PLD <a_mode2>

Table 3-3 shows addressing mode 2 operations.

Table 3-3 Addressing mode 2

Addressing mode 2	
Immediate offset	[Rn, #+/-12bit_Offset]
Register offset	[Rn, +/-Rm]
Scaled register offset	[Rn, +/-Rm, LSL #5bit_shift_imm]
	[Rn, +/-Rm, LSR #5bit_shift_imm]
	[Rn, +/-Rm, ASR #5bit_shift_imm]
	[Rn, +/-Rm, ROR #5bit_shift_imm]
	[Rn, +/-Rm, RRX]
Preindexed offset	
Immediate	[Rn, #+/-12bit_Offset]!
Register	[Rn, +/-Rm]!
Scaled register	[Rn, +/-Rm, LSL #5bit_shift_imm]!
	[Rn, +/-Rm, LSR #5bit_shift_imm]!
	[Rn, +/-Rm, ASR #5bit_shift_imm]!
	[Rn, +/-Rm, ROR #5bit_shift_imm]!
	[Rn, +/-Rm, RRX]!

Table 3-3 Addressing mode 2 (continued)

Addressing mode 2	
Postindexed offset	
Immediate	[Rn], #+/-12bit_Offset
Register	[Rn], +/-Rm
Scaled register	[Rn], +/-Rm, LSL #5bit_shift_imm
	[Rn], +/-Rm, LSR #5bit_shift_imm
	[Rn], +/-Rm, ASR #5bit_shift_imm
	[Rn], +/-Rm, ROR #5bit_shift_imm
	[Rn], +/-Rm, RRX]

Table 3-4 shows privileged addressing mode 2 operations.

Table 3-4 Addressing mode 2, privileged

Addressing mode 2 Privileged	
Immediate offset	[Rn, #+/-12bit_Offset]
Register offset	[Rn, +/-Rm]
Scaled register offset	[Rn, +/-Rm, LSL #5bit_shift_imm]
	[Rn, +/-Rm, LSR #5bit_shift_imm]
	[Rn, +/-Rm, ASR #5bit_shift_imm]
	[Rn, +/-Rm, ROR #5bit_shift_imm]
	[Rn, +/-Rm, RRX]
Postindexed offset	
Immediate	[Rn], #+/-12bit_Offset
Register	[Rn], +/-Rm
Scaled register	[Rn], +/-Rm, LSL #5bit_shift_imm
	[Rn], +/-Rm, LSR #5bit_shift_imm

Table 3-4 Addressing mode 2, privileged (continued)

Addressing mode 2 Privileged	
	[Rn], +/-Rm, ASR #5bit_shift_imm
	[Rn], +/-Rm, ROR #5bit_shift_imm
	[Rn, +/-Rm, RRX]

Table 3-5 shows addressing mode 3 operations.

Table 3-5 Addressing mode 3

Addressing mode 3 Signed byte, and halfword data transfer	
Immediate offset	[Rn, #+/-8bit_Offset]
Preindexed	[Rn, #+/-8bit_Offset]!
Postindexed	[Rn], #+/-8bit_Offset
Register	[Rn, +/-Rm]
Preindexed	[Rn, +/-Rm]!
Postindexed	[Rn], +/-Rm

Table 3-6 shows addressing mode 4 (load) operations.

Table 3-6 Addressing mode 4, load

Addressing mode 4 Load	Stack type
IA increment after	FD full descending
IB increment before	ED empty descending
DA decrement after	FA full ascending
DB decrement before	EA empty ascending

Table 3-7 shows addressing mode 4 (store) operations.

Table 3-7 Addressing mode 4, store

Addressing mode 4 Store	Stack type
IA increment after	EA empty ascending
IB increment before	FA full ascending
DA decrement after	ED empty descending
DB decrement before	FD full descending

Table 3-8 shows addressing mode 5 (load) operations.

Table 3-8 Addressing mode 5

Addressing mode 5 Coprocesor data transfer	
Immediate offset	[Rn, #+/- (8bit_Offset*4)]
Preindexed	[Rn, #+/- (8bit_Offset*4)]!
Postindexed	[Rn], #+/- (8bit_Offset*4)

Table 3-9 shows example uses of Oprnd2.

Table 3-9 Oprnd2 examples

Oprnd2	Example
Immediate value	#32bit_Imm
Logical shift left	Rm LSL #5bit_Imm
Logical shift right	Rm LSR #5bit_Imm
Arithmetic shift right	Rm ASR #5bit_Imm
Rotate right	Rm ROR #5bit_Imm
Register	Rm
Logical shift left	Rm LSL Rs
Logical shift right	Rm LSR Rs

Table 3-9 Oprnd2 examples (continued)

Oprnd2	Example
Arithmetic shift right	Rm ASR Rs
Rotate right	Rm ROR Rs
Rotate right extended	Rm RRX

Table 3-10 shows the suffixes to set fields in MSR operations.

Table 3-10 Suffixes to set fields

Suffix	Sets
_c	Control field mask bit (bit 3)
_x	Extension field mask bit (bit 2)
_s	Status field mask bit (bit 1)
_f	Flags field mask bit (bit 0)

Table 3-11 shows the condition code extensions.

Table 3-11 Condition fields

Extension	Description
EQ	Equal
NE	Not equal
CS	Unsigned higher or same
CC	Unsigned lower
MI	Negative
PL	Positive or zero
VS	Overflow
VC	No overflow
HI	Unsigned higher
LS	Unsigned lower, or same

Table 3-11 Condition fields (continued)

Extension	Description
GE	Greater, or equal
LT	Less than
GT	Greater than
LE	Less than, or equal
AL	Always

3.2.2 Thumb instruction summary

Table 3-12 summarizes the Thumb instruction set.

Table 3-12 Thumb instruction summary

Operation		Assembler
Move	Immediate	MOV Rd, #8bit_Imm
	High to low	MOV Rd, Hs
	Low to high	MOV Hd, Rs
	High to high	MOV Hd, Hs
Arithmetic	Add	ADD Rd, Rs, #3bit_Imm
	Add low and low	ADD Rd, Rs, Rn
	Add high to low	ADD Rd, Hs
	Add low to high	ADD Hd, Rs
	Add high to high	ADD Hd, Hs
	Add immediate	ADD Rd, #8bit_Imm
	Add value to SP	ADD SP, #7bit_Imm
	Add with carry	ADC Rd, Rs
	Subtract	SUB Rd, Rs, Rn SUB Rd, Rs, #3bit_Imm
	Subtract immediate	SUB Rd, #8bit_Imm
	Subtract with carry	SBC Rd, Rs

Table 3-12 Thumb instruction summary (continued)

Operation	Assembler
Negate	NEG Rd, Rs
Multiply	MUL Rd, Rs
Compare low and low	CMP Rd, Rs
Compare low and high	CMP Rd, Hs
Compare high and low	CMP Hd, Rs
Compare high and high	CMP Hd, Hs
Compare negative	CMN Rd, Rs
Compare immediate	CMP Rd, #8bit_Imm
Logical	AND Rd, Rs
EOR	EOR Rd, Rs
OR	ORR Rd, Rs
Bit clear	BIC Rd, Rs
Move NOT	MVN Rd, Rs
Test bits	TST Rd, Rs
Shift/rotate	Logical shift left LSL Rd, Rs, #5bit_shift_imm LSL Rd, Rs
	Logical shift right LSR Rd, Rs, #5bit_shift_imm LSR Rd, Rs
	Arithmetic shift right ASR Rd, Rs, #5bit_shift_imm ASR Rd, Rs
	Rotate right ROR Rd, Rs
Branch	Conditional
	If Z set BEQ label
	If Z clear BNE label
	If C set BCS label
	If C clear BCC label
	If N set BMI label

Table 3-12 Thumb instruction summary (continued)

Operation	Assembler
If N clear	BPL <i>label</i>
If V set	BVS <i>label</i>
If V clear	BVC <i>label</i>
If C set and Z clear	BHI <i>label</i>
If C clear and Z set	BLS <i>label</i>
If N set and V set, or if N clear and V clear	BGE <i>label</i>
If N set and V clear, or if N clear and V set	BLT <i>label</i>
If Z clear and N or V set, or if Z clear, and N or V clear	BGT <i>label</i>
If Z set, or N set and V clear, or N clear and V set	BLE <i>label</i>
Unconditional	B <i>label</i>
Long branch with link	BL <i>label</i>
Optional state change	
To address held in Lo reg	BX <i>Rs</i>
To address held in Hi reg	BX <i>Hs</i>
Load	
With immediate offset	
Word	LDR <i>Rd</i> , [<i>Rb</i> , # <i>7bit_offset</i>]
Halfword	LDRH <i>Rd</i> , [<i>Rb</i> , # <i>6bit_offset</i>]
Byte	LDRB <i>Rd</i> , [<i>Rb</i> , # <i>5bit_offset</i>]
With register offset	
Word	LDR <i>Rd</i> , [<i>Rb</i> , <i>Ro</i>]
Halfword	LDRH <i>Rd</i> , [<i>Rb</i> , <i>Ro</i>]
Signed halfword	LDRSH <i>Rd</i> , [<i>Rb</i> , <i>Ro</i>]
Byte	LDRB <i>Rd</i> , [<i>Rb</i> , <i>Ro</i>]
Signed byte	LDRSB <i>Rd</i> , [<i>Rb</i> , <i>Ro</i>]
PC-relative	LDR <i>Rd</i> , [<i>PC</i> , # <i>10bit_Offset</i>]

Table 3-12 Thumb instruction summary (continued)

Operation	Assembler
SP-relative Address	LDR Rd, [SP, #10bit_Offset]
Using PC	ADD Rd, PC, #10bit_Offset
Using SP	ADD Rd, SP, #10bit_Offset
Multiple	LDMIA Rb!, <reglist>
Store	
With immediate offset	
Word	STR Rd, [Rb, #7bit_offset]
Halfword	STRH Rd, [Rb, #6bit_offset]
Byte	STRB Rd, [Rb, #5bit_offset]
With register offset	
Word	STR Rd, [Rb, Ro]
Halfword	STRH Rd, [Rb, Ro]
Byte	STRB Rd, [Rb, Ro]
SP-relative	STR Rd, [SP, #10bit_offset]
Multiple	STMIA Rb!, <reglist>
Push/pop	
Push registers onto stack	PUSH <reglist>
Push LR and registers onto stack	PUSH <reglist, LR>
Pop registers from stack	POP <reglist>
Pop registers and PC from stack	POP <reglist, PC>
Software interrupt	SWI 8bit_Imm

Chapter 4

System Control Coprocessor

This chapter describes the registers of the *system control coprocessor*. It contains the following sections:

- *About the system control coprocessor* on page 4-2
- *Register descriptions* on page 4-5.

4.1 About the system control coprocessor

The ARM10 programmer's model, including a detailed instruction set specification, is described in the *ARM Architecture Reference Manual*. The programmer's model of the ARM1020E processor is the same as the programmer's model of the ARM10 integer unit, but extended in the following ways:

- The system control coprocessor (CP15) provides additional registers for configuring and controlling caches, MMU, protection system, power-down, and clocking mode.
- The MMU page tables define the virtual-to-physical address mapping, page and section access permissions, cache, and write buffer configuration. These are created by the operating system software and accessed automatically by the MMU hardware whenever an instruction read or data access causes a TLB miss.

4.1.1 Accessing CP15 registers

CP15 registers can be accessed only with MRC and MCR instructions in a privileged mode. Figure 4-1 and Figure 4-2 show the MCR and MRC instruction formats.

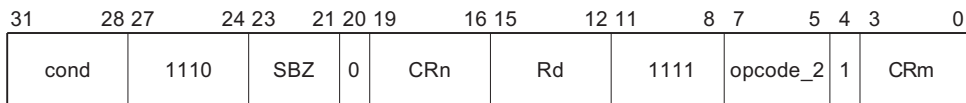


Figure 4-1 CP15 MCR instruction format

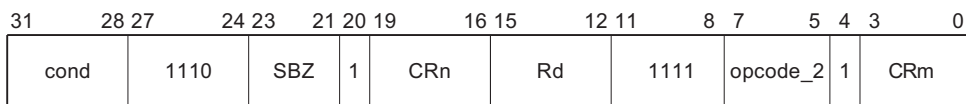


Figure 4-2 CP15 MRC instruction format

The assembler for these instructions is:

```
MCR{cond} P15, opcode_1, Rd, CRn, CRm, opcode_2
MRC{cond} P15, opcode_1, Rd, CRn, CRm, opcode_2
```

Other CP15 instructions (CDP, LDC, and STC), with MRC and MCR instructions executed in User mode, are Undefined. Any MCR or MRC instruction that is not executed in a privileged mode takes the Undefined instruction trap. The CRn field of MRC and MCR instructions

specifies the coprocessor register to access. The CRm fields, opcode_1, and opcode_2, specify a particular action when addressing registers. Refer to the *ARM Architecture Reference Manual* for details of these fields.

4.1.2 Summary of CP15 registers

CP15 contains 16 registers. Table 4-1 shows their read and write functions.

Table 4-1 CP15 register summary

Register	Register name	Reads	Writes
CP15 c0	Device ID register Cache type register	Cache ID and type information	-
CP15 c1	Control register 1	Control	Control
CP15 c2	Translation table base register	Translation table base	Translation table base
CP15 c3	Domain access control register	Domain access control	Domain access control
CP15 c4	-	Unpredictable	Unpredictable
CP15 c5	Fault status register	Fault status	Fault status
CP15 c6	Fault address register	Fault address	Fault address
CP15 c7	Index cache operations register VA cache operations register	Unpredictable	Cache operations
CP15 c8	TLB operations register	Unpredictable	MMU operations
CP15 c9	Cache lockdown register	Cache lockdown	Cache lockdown
CP15 c10	TLB lockdown register	TLB lockdown	TLB lockdown
CP15 c11	-	Undefined	Undefined
CP15 c12	-	Undefined	Undefined
CP15 c13	Process ID register Context ID register	Process ID Context ID	Process ID and context ID
CP15 c14	-	Undefined	Undefined
CP15 c15	PLL configuration register Power manager status register Power manager receive data register Power manager transmit data register Control register 2	PLL configuration Power manager status Power manager receive data Power manager transmit data Cache and soft TLB control	PLL configuration Power manager status Power manager receive data Power manager transmit data Cache and soft TLB control

All CP15 register bits that are defined and contain state are cleared by reset except:

- the V bit in CP15 c1, which takes the value of input signal **HIVECSINIT**
- the B bit in CP15 c1, which takes the value of input signal **BIGENDINIT**.

4.1.3 Address types

The ARM processor uses three address types:

- *Virtual Address (VA)*
- *Modified Virtual Address (MVA)*
- *Physical Address (PA)*.

Table 4-2 shows the address types.

Table 4-2 Address types

	Integer unit	Caches and TLBs	AMBA bus
Address type	Virtual address	Modified virtual address	Physical address

Figure 1-1 on page 1-5 shows paths for these addresses. When the integer core requests an instruction, the following address manipulation occurs:

1. The integer unit issues the VA of the instruction.
2. The VA is translated using the process ID to the MVA. The instruction cache and MMU perform a lookup using the MVA.
3. If the protection check carried out by the MMU on the MVA does not abort, and the MVA tag is in the instruction cache, then the instruction data is returned to the integer unit.
4. If the MVA tag is not in the instruction cache, causing an instruction cache miss, then the MMU performs a translation to produce the *Instruction PA (IPA)*.
5. The PA is passed to the AMBA bus interface to perform an external access.

4.2 Register descriptions

This section describes the CP15 registers:

- *CP15 c0, Device ID and Cache Type Registers*
- *CP15 c1, Control Register 1* on page 4-7
- *CP15 c2, Translation Table Base Register* on page 4-10
- *CP15 c3, Domain Access Control Register* on page 4-11
- *CP15 c4* on page 4-12
- *CP15 c5, Fault Status Register* on page 4-12
- *CP15 c6, Fault Address Register* on page 4-15
- *CP15 c7, Index and VA Cache Operations Registers* on page 4-16
- *CP15 c8, TLB Operations Register* on page 4-19
- *CP15 c9, Cache Lockdown Register* on page 4-20
- *CP15 c10, TLB Lockdown Register* on page 4-22
- *CP15 c11* on page 4-23
- *CP15 c12* on page 4-23
- *CP15 c13, Process ID and Context ID Registers* on page 4-23
- *CP15 c14* on page 4-26
- *CP15 c15* on page 4-26.

4.2.1 CP15 c0, Device ID and Cache Type Registers

The device ID and cache type registers are read-only. Depending on the value of `opcode_2`, reading CP15 c0 returns one of the following:

- When `opcode_2` is 0, reading CP15 c0 returns the device ID value `0x4105A20r`, where `r` is the revision.
- When `opcode_2` is 1, reading CP15 c0 returns the cache information value `0x0D1B21B2`, which reflects the type, size, associativity, and line length of the ICache and DCache.

The CRm field Should Be Zero when reading CP15 c0. Writing to CP15 c0 is Unpredictable.

Table 4-3 shows the instructions for using the device ID and cache type registers.

Table 4-3 Device ID and Cache Type Register instructions

Function	Data	Instruction
Read device ID	ARM processor device ID	MRC p15, 0, Rd, c0, c0, 0
Read cache information	ICache and DCache type	MRC p15, 0, Rd, c0, c0, 1

Device ID Register

Figure 4-3 shows the device ID register bit fields.

31	28	27	24	23	20	19	16	15	12	11	8	7	4	3	0	
0	1	0	0	0	0	0	1	0	0	0	1	0	0	0	0	Revision

Figure 4-3 Device ID Register

Table 4-4 describes the bit fields of the Device ID Register.

Table 4-4 Encoding of the Device ID Register

Bits	Meaning
[31:24]	Implementor = 0x41 (ARM Limited).
[23:20]	Variant 0x1.
[19:16]	Architecture. 0x5 = ARM architecture version 5TE.
[15:4]	Contain the three-digit part number, 0xA20.
[3:0]	Contain the revision number for the ARM processor.

Cache Type Register

Figure 4-4 on page 4-7 shows the cache type register bit fields.

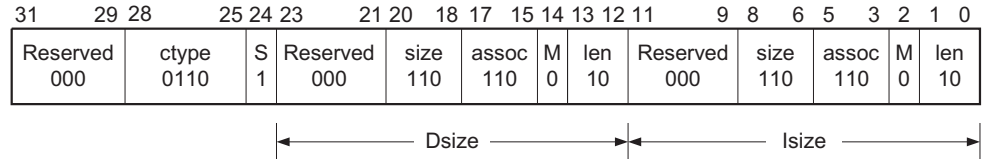
**Figure 4-4 Cache Type Register**

Table 4-5 describes the bit fields of the Cache Type Register.

Table 4-5 Encoding of the Cache Type Register

Bits	Meaning	Value	Notes
[31:29]	Reserved	000	-
[28:25]	Cache class	0110	Cache-clean-step operation Cache-invalidate-step operation Lock-down facilities
[24]	Harvard architecture	1	-
[23:21]	Reserved	000	-
[20:18]	Data cache sizes	110	32KB
[17:15]	Data cache associativity	110	64-way associative
[14]	Data cache parameters	0	Associativity and size are equal
[13:12]	Data cache line length	10	Eight words per line
[11:9]	Reserved	000	-
[8:6]	Instruction cache size	110	32KB
[5:3]	Instruction cache associativity	110	64-way set associative
[2]	Instruction cache parameters	0	Associativity and size are equal
[1:0]	Instruction cache line length	10	Eight words per line

4.2.2 CP15 c1, Control Register 1

The read/write Control Register 1:

- enables fast interrupts

- selects the T bit after a load PC operation
- selects random or round-robin victim replacement
- selects high-address or low-address vector locations
- enables the ICache, DCache, and write buffer
- enables branch prediction
- enables ROM protection and MMU protection
- selects big-endian or little-endian operation
- enables fault checking of address alignment
- enables the MMU.

Use a read-modify-write sequence to access control register 1. For both reading and writing, the CRm and opcode_2 fields should be zero. Table 4-6 shows the instructions for using control register 1.

Table 4-6 Control Register 1 instructions

Operation	Data	Instruction
Read configuration	Configuration data	MCR p15, 0, Rd, c1, c0, 0
Write configuration	Configuration data	MCR p15, 0, Rd, c1, c0, 0

All defined control bits are cleared on reset except:

- The V bit is cleared at reset if the **HIVECSINIT** signal is LOW, or set if the **HIVECSINIT** signal is HIGH.
- The B bit is cleared at reset if the **BIGENDINIT** signal is LOW, or set if the **BIGENDINIT** signal is HIGH.

Figure 4-5 shows the control register 1 bit fields.

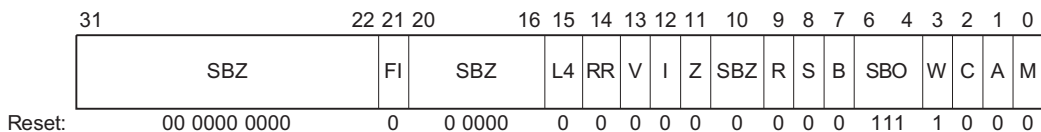


Figure 4-5 Control Register 1

Using a read-modify-write sequence when changing control register 1 provides the greatest future compatibility. Table 4-7 describes the control register 1 bit fields.

Table 4-7 Encoding of Control Register 1

Bits	Name	Meaning
[31:22]	-	Reading returns an Unpredictable value. When written, Should Be Zero, or a value read from bits [31:18] on the same processor.
[21]	FI	Fast interrupt bit. Disables HUM and reduces write buffer to half depth, or four doublewords. Reset clears FI. 1 = write buffer is four slots, HUM disabled, streaming disabled, core is blocking 0 = write buffer is eight slots, HUM and streaming enabled, core is nonblocking.
[20:16]	-	Should Be Zero.
[15]	L4	Configures if load instructions to the PC set the T bit: 0= loads to PC set the T bit 1= loads to PC do not set T bit (ARMv4 behavior). For more details, see the <i>ARM Architecture Reference Manual</i> .
[14]	RR	ICache and DCache round-robin replacement bit. Reset clears RR. 1 = round-robin replacement enabled 0 = random replacement.
[13]	V	Exception vector location bit. Reset clears V. 1 = vectors start at 0xFFFF 0000 0 = vectors start at 0x0000 0000.
[12]	I	Instruction cache enable bit. Reset clears I. 1 = ICache enabled 0 = ICache disabled.
[11]	Z	Branch prediction enable bit. Reset clears Z. 1 = branch prediction enabled 0 = branch prediction disabled.
[10]	-	Should Be Zero.
[9]	R	ROM protection enable bit. Reset clears R. 1 = ROM protection enabled 0 = ROM protection disabled.
[8]	S	System protection enable bit. Reset clears S. 1 = IMMU and DMMU protection enabled 0 = IMMU and DMMU protection disabled.

Table 4-7 Encoding of Control Register 1 (continued)

Bits	Name	Meaning
[7]	B	Big-endian bit. Reset clears B. 1 = big-endian operation 0 = little-endian operation.
[6:4]	-	Reading returns 111. When written, Should Be One.
[3]	W	Write buffer enable bit. Reset sets W. 1 = write buffer enabled 0 = write buffer disabled.
[2]	C	DCache enable bit. Reset clears C. 1 = DCache enabled 0 = DCache disabled.
[1]	A	Address alignment fault checking enable bit. Reset clears A. 1 = fault checking of address alignment enabled 0 = fault checking of address alignment disabled.
[0]	M	MMU enable bit. Reset clears M. 1 = IMMU and DMMU enabled 0 = IMMU and DMMU disabled.

———— **Note** —————

Be careful with the address mapping of the code sequence used to enable the MMU (see *Enabling the MMU* on page 5-36).

See *DCache and write buffer enable/disable* on page 6-8 for restrictions, and for effects of having caches enabled when the MMU is disabled.

4.2.3 CP15 c2, Translation Table Base Register

The Translation Table Base Register, TTBR, contains the *Translation Table Base* (TTB) of the level 1 translation table.

When read, bits [31:14] return the pointer to the level 1 translation table, and bits [13:0] return an Unpredictable value.

Writing to TTBR updates the pointer to the level 1 translation table in bits [31:14]. Bits [13:0] Should Be Zero.

The CRm and opcode_2 fields Should Be Zero when writing to TTBR.

Table 4-8 shows the instructions for using TTBR.

Table 4-8 Translation Table Base Register instructions

Operation	Data	Instruction
Read TTBR	TTBR address	MRC p15, 0, Rd, c2, c0, 0
Write TTBR	TTBR address	MCR p15, 0, Rd, c2, c0, 0

Figure 4-6 shows the TTBR bit fields.

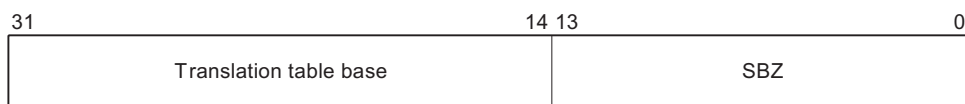


Figure 4-6 Translation Table Base Register

4.2.4 CP15 c3, Domain Access Control Register

The Domain Access Control Register, DACR, contains 16 discrete 2-bit domain access control fields, each of which defines the access permissions for one of the 16 domains, D15-D0.

Reading DACR returns the value of the domain access control bit fields. Writing to DACR writes the value of the domain access control bit fields.

The CRm and opcode_2 fields Should Be Zero when writing to DACR.

Table 4-9 shows the instructions for using DACR.

Table 4-9 Domain Access Control Register instructions

Operation	Data	Instruction
Read domain access	Domain 15 to 0 access control	MRC p15, 0, Rd, c3, c0, 0
Write domain access	Domain 15 to 0 access control	MCR p15, 0, Rd, c3, c0, 0

Figure 4-7 on page 4-12 shows the DACR bit fields.

31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	13	12	11	10	9	8	7	6	5	4	3	2	1	0	
D15	D14	D13	D12	D11	D10	D9	D8	D7	D6	D5	D4	D3	D2	D1	D0																	

Figure 4-7 Domain Access Control Register

Table 4-10 describes the DACR bit fields.

Table 4-10 Encoding of the Domain Access Control Register

Bits	Meaning
D15-D0	Domain access control: 00 = no access; access generates domain fault. 01 = client access; access permissions are checked. 10 = reserved; behaves as <i>no access</i> domain. 11 = manager; access permissions are not checked.

Any write to DACR causes all unlocked TLB entries to be invalidated. If you change the domain access control field corresponding to a locked TLB entry, you must invalidate that entry in the TLB using the *invalidate single entry* operation and reload it. Ideally, a program that locks entries in the TLB maps those locked entries to unmodified DAC fields.

4.2.5 CP15 c4

Reading or writing CP15 register c4 is Undefined.

4.2.6 CP15 c5, Fault Status Register

The Data Fault Status Register, DFSR, contains the source of the last data fault. It indicates the domain and type of access being attempted when an abort occurred. The DFSR also indicates the domain where the abort occurred.

Data Fault Status Register

The DFSR is for checking all Data Aborts and watchpoints. The data FSR maps the debug event to a watchpoint. The instruction FSR is used to check all Prefetch Aborts and breakpoints. The instruction FSR maps the debug event to a breakpoint.

Table 4-13 Priority of data fault types (continued)

Priority	Source	Status	Domain	FAR
	Translation section	0101	Invalid	Valid
	Translation page	0111	Valid	Valid
	Domain section	1001	Valid	Valid
	Domain page	1011	Valid	Valid
	Permission section	1101	Valid	Valid
	Permission page	1111	Valid	Valid
	External abort	1010	Valid	Valid
Lowest	Debug event	0010	Valid	Valid

Reading DFSR returns the value of the DFSR.

Writing to DFSR changes the DFSR to the value of the data written. This is useful for a debugger to restore the value of the DFSR. The register must be written using a read-modify-write sequence. Bits [31:8] Should Be Zero.

The CRm field Should Be Zero when reading or writing DFSR.

Instruction Fault Status Register

The IFSR is for checking all Prefetch Aborts and breakpoints. The IFSR maps the debug event to a breakpoint.

Table 4-14 shows the instructions for using the IFSR.

Table 4-14 Instruction Fault Status Register instructions

Operation	Data	Instruction
Read IFSR	IFSR instruction	MCR p15, 0, Rd, c5, c0, 1
Write IFSR	IFSR instruction	MCR p15, 0, Rd, c5, c0, 1

Figure 4-9 on page 4-15 shows the IFSR bit fields.

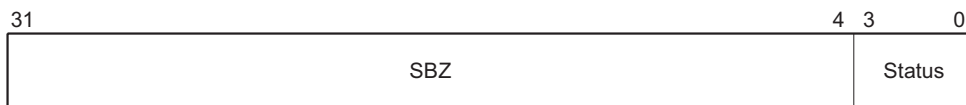
**Figure 4-9 Instruction Fault Status Register**

Table 4-15 describes the IFSR bit fields.

Table 4-15 Encoding of the Instruction Fault Status Register

Bits	Meaning
[31:4]	Should Be Zero
[3:0]	Type of fault generated (see Table 4-13 on page 4-13)

Reading IFSR returns the value of the IFSR.

Writing to IFSR changes the IFSR to the value written. This is useful for a debugger to restore the value of the IFSR. The register must be written using a read-modify-write sequence. Bits [31:8] Should Be Zero.

The CRm field Should Be Zero when reading or writing IFSR.

4.2.7 CP15 c6, Fault Address Register

The Fault Address Register, FAR, holds the VA of the access that was attempted when a fault occurred.

Table 4-16 shows the instructions for using the FAR.

Table 4-16 Fault Address Register instructions

Operation	Data	Instruction
Read data FAR	FAR data	MRC p15, 0, Rd, c6, c0, 0
Write data FAR	FAR data	MCR p15, 0, Rd, c6, c0, 0
Read instruction FAR	FAR data	MRC p15, 0, Rd, c6, c0, 1
Write instruction FAR	FAR data	MCR p15, 0, Rd, c6, c0, 1

Reading FAR returns the value of either the data FAR or the instruction FAR as specified by the opcode_2 value.

Writing to FAR changes the FAR to the value of the data written. This is useful for a debugger to restore the value of the FAR.

The CRm fields Should Be Zero when reading or writing FAR.

Figure 4-10 shows the FAR bit field.



Figure 4-10 Fault Address Register

The data FAR contains the address of the memory access that caused either a watchpoint or a Data Abort. The instruction FAR contains the address (PC + 8 for ARM, PC + 4 for Thumb) of the instruction that triggered said memory access.

4.2.8 CP15 c7, Index and VA Cache Operations Registers

The Index and VA Cache Operations Registers are write-only registers for managing the ICache and DCache.

Table 4-17 shows the instructions for performing index and VA cache operations.

Table 4-17 Cache Operations Register instructions

Function	Data	Instruction
Invalidate caches	Should Be Zero	MCR p15, 0, Rd, c7, c7, 0
Invalidate ICache	Should Be Zero	MCR p15, 0, Rd, c7, c5, 0
Invalidate ICache single entry using VA	Virtual address	MCR p15, 0, Rd, c7, c5, 1
Prefetch ICache line	Virtual address	MCR p15, 0, Rd, c7, c13, 1
Invalidate DCache	Should Be Zero	MCR p15, 0, Rd, c7, c6, 0
Invalidate DCache single entry using VA	Virtual address	MCR p15, 0, Rd, c7, c6, 1
Clean DCache single entry using VA	Virtual address	MCR p15, 0, Rd, c7, c10, 1
Clean and invalidate DCache single entry using VA	Virtual address	MCR p15, 0, Rd, c7, c14, 1
Clean DCache single entry using index	Index, segment format	MCR p15, 0, Rd, c7, c10, 2

Table 4-17 Cache Operations Register instructions (continued)

Function	Data	Instruction
Clean and invalidate DCache entry using index	Index, segment format	MCR p15, 0, Rd, c7, c14, 2
Empty write buffer	Should Be Zero	MCR p15, 0, Rd, c7, c10, 4
Wait for interrupt	Should Be Zero	MCR p15, 0, Rd, c7, c0, 4

The opcode_2 and CRm fields in the MCR instruction select the cache operation. Writing opcode_2 or CRm values other than those shown in Table 4-17 on page 4-16 is Unpredictable.

Reading the index and VA cache operations registers is Unpredictable.

Table 4-18 describes the cache operations in more detail.

———— **Note** —————

Dirty data is data that has been modified in the cache but not yet copied back to main memory.

Table 4-18 Cache operation descriptions

Function	Description
Invalidate cache	Invalidates all cache data, including any dirty data. Use with caution.
Invalidate single entry using VA	Invalidates a single cache line, including any dirty data. Use with caution.
Clean single DCache entry using either index or VA	Writes the specified cache line to main memory if the line is marked valid and dirty and is from a write-back memory region and marks the line as not dirty. The valid bit is unchanged.
Clean and invalidate single DCache entry using either index or VA	Writes the specified cache line to main memory if the line is marked valid and dirty, and is from a write-back memory region. The line is marked not valid.
Prefetch cache line	Performs an ICache lookup of the specified address. If the cache misses, and the region is cachable, a linefill is performed.

Index Cache Operation Register

The operations that act on a single cache line identify the line using the contents of Rd as the address, passed in the MCR instruction. Figure 4-11 shows the index cache operation register.



Figure 4-11 Index Cache Operations Register

Table 4-19 describes the bit fields of the index cache operation register.

Table 4-19 Encoding of the Index Cache Operations Register

Bits	Meaning
[31:26]	Index in segment being accessed
[25:9]	Should Be Zero
[8:5]	Segment being accessed
[4:3]	64-bit double word being accessed
[2:0]	Should Be Zero

The index tag format of Example 4-1 is for accessing a specific line in the cache. Example 4-1 shows the command clean D single entry (using index).

Example 4-1 Clean D single entry (using index)

```

;code is specific to the ARM1020E macrocell with 32KB caches
MOV R0, #0:SHL:5 ;select segment
seg_loop
MOV R1, #0:SHL:26 ;select index
line_loop
ORR R2,R1,R0
MCR p15,0,R2,c7,c10,2
ADD R1,R1,#1:SHL:26 ;increment index
CMP R1,#0 ;check for index overflow
BNE line_loop
ADD R0,R0,#1:SHL:5 ;increment segment

```

```

CMP R0,#1:SHL:8      ;check for segment overflow
BNE seg_loop

```

VA Cache Operations Register

The VA Cache Operations Register is useful for invalidating a particular address or range of addresses in the caches. Figure 4-12 shows the bit fields of the VA cache operations register.

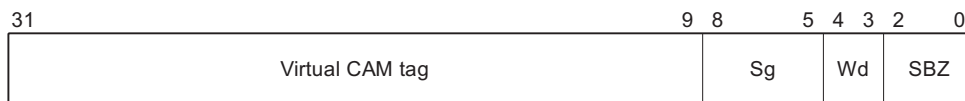


Figure 4-12 VA Cache Operations Register

Table 4-20 describes the bit fields of the VA cache operations register.

Table 4-20 Encoding of the VA Cache Operations Register

Bits	Meaning
[31:9]	Virtual <i>Content Addressable Memory</i> (CAM) tag
[8:5]	Segment being accessed
[4:3]	64-bit double word being accessed
[2:0]	Should Be Zero

Note

ICache prefetch operations and DCache clean operations are performed requested-word-first.

4.2.9 CP15 c8, TLB Operations Register

The TLB Operations Register is a write-only register for managing the instruction TLB and the data TLB. Reading from the TLB operations register is Unpredictable.

Table 4-21 shows the instructions for performing TLB operations.

Table 4-21 TLB Operations Register instructions

Operation	Data	Instruction
Invalidate instruction and data TLBs	Should Be Zero	MCR p15, 0, Rd, c8, c7, 0
Invalidate instruction TLB	Should Be Zero	MCR p15, 0, Rd, c8, c5, 0
Invalidate instruction TLB single entry (using VA)	Virtual address	MCR p15, 0, Rd, c8, c5, 1
Invalidate data TLB	Should Be Zero	MCR p15, 0, Rd, c8, c6, 0
Invalidate data TLB single entry (using VA)	Virtual address	MCR p15, 0, Rd, c8, c6, 1

The opcode_2 and CRm fields in the MCR instruction select the TLB operation. Writing opcode_2 or CRm values other than those shown in Table 4-21 is Unpredictable.

Figure 4-13 shows the TLB operations register bit fields.



Figure 4-13 TLB Operations Register

Note

Invalidating the full TLB invalidates all the unlocked entries in the TLB. Invalidating TLB single entry functions invalidates any TLB entry corresponding to the VA given in Rd, regardless of its locked state (see *CP15 c10, TLB Lockdown Register* on page 4-22).

4.2.10 CP15 c9, Cache Lockdown Register

The Cache Lockdown Register enables software to:

- control which line ICache or DCache line is loaded for a linefill by changing the base value or the victim counter value respectively
- prevent ICache or DCache lines from being replaced during a linefill, locking them into the cache.

Table 4-22 shows the instructions for using the cache lockdown register.

Table 4-22 Cache Lockdown Register instructions

Operation	Data	Instruction
Read DCache lockdown base	Base	MRC p15, 0, Rd, c9, c0, 0
Write DCache victim and lockdown base	Victim = base	MCR p15, 0, Rd, c9, c0, 0
Read ICache lockdown base	Base	MRC p15, 0, Rd, c9, c0, 1
Write ICache victim and lockdown base	Victim = base	MCR p15, 0, Rd, c9, c0, 1

Reading the cache lockdown register returns the value of the cache lockdown register, which is the base pointer for all cache segments. Reset clears the cache lockdown register.

Note

Only bits [31:26] are returned. Bits [25:0] are zero.

Figure 4-14 shows the bit fields of the cache lockdown register

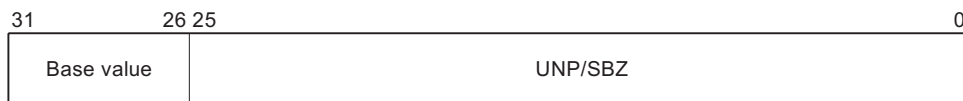


Figure 4-14 Cache Lockdown Register

Writing to the cache lockdown register updates the base pointer and the current victim counter value for all cache segments. Bits [25:0] Should Be Zero. The next linefill uses the victim counter value, then increments the victim counter. The victim counter continues incrementing on linefills and wraps around to the base pointer. For example, setting the base pointer to $0x3$ prevents the victim counter from selecting entries $0x0$ to $0x2$, locking them into the cache.

The victim counter specifies the cache line to be used as the victim for the next linefill. The counter is incremented using either a random or round-robin replacement policy, determined by the state of the RR bit in control register 1, CP15 c1. The victim counter generates values from base to base + 63. This locks lines with index values from 0 to base - 1, with an upper limit of 63 locked entries in the DCache. If base = 0 there are no locked lines.

Example 4-2 shows how to load a single entry into line 0 and lock it down.

Example 4-2 Updating the base pointer and current victim pointer

```
MCR to CP15 c9, Victim=Base=0x0
MCR to cause an I prefetch, LDR/LDM, depending on whether it is ICache or
DCache. Assuming the appropriate cache misses, a linefill occurs to line 0.
MCR to CP15 c9, Victim=Base=0x1
```

Further linefills now occur into lines 1 to 63.

4.2.11 CP15 c10, TLB Lockdown Register

There is a TLB Lockdown Register for each TLB. Reading the TLB lockdown register returns the value of the TLB lockdown counter base register, the current victim counter value, and the preserve bit. The TLB lockdown register is cleared at reset.

Writing to the TLB lockdown register updates the TLB lockdown counter base register, the current victim counter value, and the state of the preserve bit. Bits [19:1] Should Be Zero. Table 4-23 shows the instructions for using the TLB lockdown register.

Table 4-23 TLB Lockdown Register instructions

Operation	Data	Instruction
Read data TLB lockdown	TLB lockdown	MRC p15, 0, Rd, c10, c0, 0
Write data TLB lockdown	TLB lockdown	MCR p15, 0, Rd, c10, c0, 0
Read instruction TLB lockdown	TLB lockdown	MRC p15, 0, Rd, c10, c0, 1
Write instruction TLB lockdown	TLB lockdown	MCR p15, 0, Rd, c10, c0, 1

Figure 4-15 shows the bit fields of the TLB lockdown register.

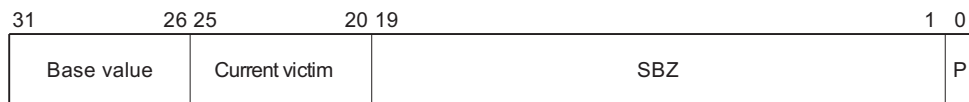


Figure 4-15 TLB Lockdown Register

The entries in the TLBs are replaced using a round-robin replacement policy. This is implemented using a victim counter that counts up continuously from entry 0 at the base value to entry 63, wrapping back from 63 to the base value each time.

There are two mechanisms to ensure that entries are not removed from the TLB:

- Locking an entry down prevents it from being selected for overwriting during a table walk. This is achieved by programming the base value to which the victim counter reloads. For example, if the bottom three entries (0 to 2) are to be locked down, the base counter must be programmed to 3.
- An entry can also be preserved during an *invalidate all* instruction. This is done by ensuring the P bit is set when the entry is loaded into the TLB.

Example 4-3 shows how to load a single entry into location 0, make it immune to *invalidate all*, and lock it down.

Example 4-3 Ensuring an entry is not removed from the TLB

```
MCR to CP15 c10, Base Value = 0, Current Victim = 0, Preserve = '1'
MCR to cause prefetch, assuming a miss occurs in the TLB then entry 0 is loaded.
MCR to CP15 c10, Base Value = 1, Current Victim = 1, Preserve = '0'
```

4.2.12 CP15 c11

Reading or writing CP15 register c11 takes the Undefined instruction trap.

4.2.13 CP15 c12

Reading or writing CP15 register c12 takes the Undefined instruction trap.

4.2.14 CP15 c13, Process ID and Context ID Registers

The process ID and context ID registers are read/write registers. Reset clears the process ID register.

Table 4-24 shows the instructions for using the process ID and context ID registers.

Table 4-24 Process ID and Context ID Register instructions

Operation	Instruction
Read process ID	MRC p15, 0, Rd, c13, c0, 0
Write process ID	MCR p15, 0, Rd, c13, c0, 0
Read context ID	MRC p15, 0, Rd, c13, c0, 1
Write context ID	MCR p15, 0, Rd, c13, c0, 1

Reading the process ID register returns the value of the process ID.

Writing to the process ID register updates the process ID. Bits [24:0] Should Be Zero. Figure 4-16 shows the bit fields of the process ID register.

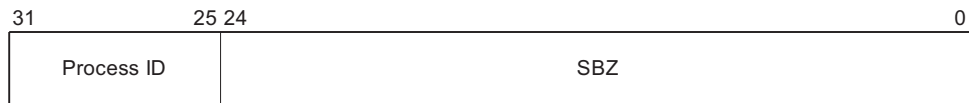


Figure 4-16 Process ID Register

The context ID register is a holding register for storing the current context of the program. Reading the context ID register returns the context ID. Writing to the context ID register updates the context ID.

Figure 4-17 shows the bit fields of the context ID register.

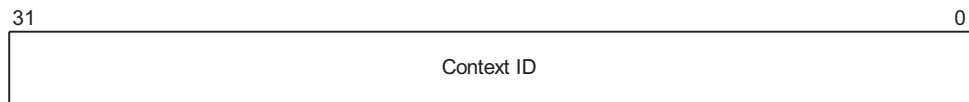


Figure 4-17 Context ID Register

Using the process ID

Addresses issued by the integer unit in the range 0 - 32MB are translated by the process ID. Address A becomes $A + (\text{process ID} \times 32\text{MB})$. This translated address is used by both the caches and MMU. Addresses above 32MB are not translated. This is shown in Figure 4-18. The process ID is a seven-bit field, enabling $128 \times 32\text{MB}$ processes to be mapped.

Note

If the process ID is zero, as it is on reset, then a flat mapping exists between the integer unit, the caches, and the MMU.

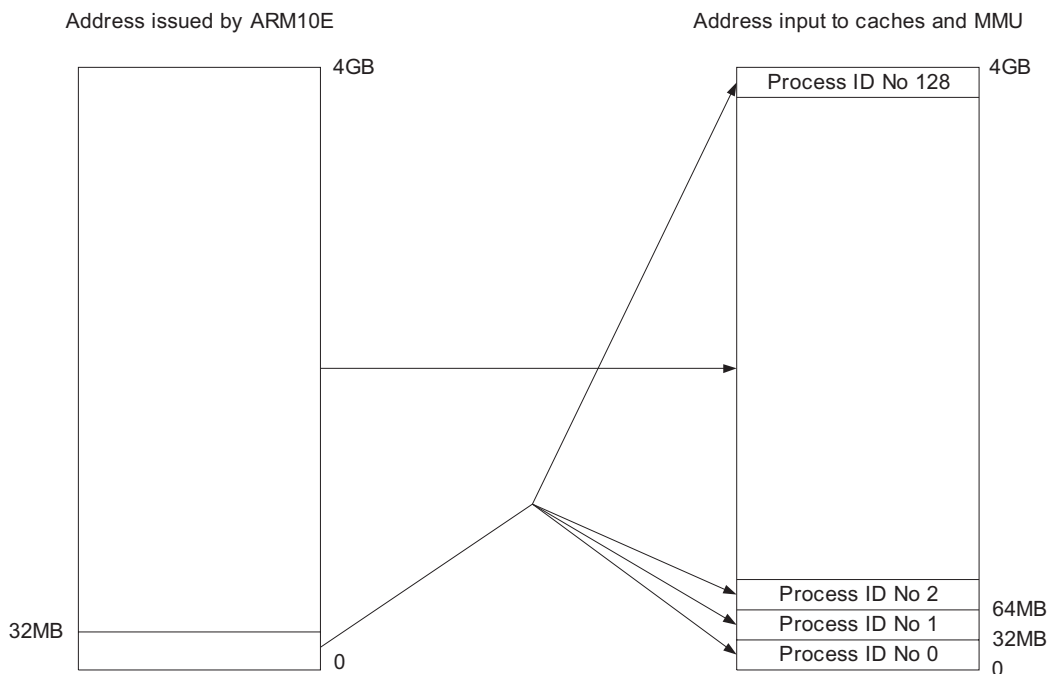


Figure 4-18 Address mapping using CP15 R13

A fast context switch is performed by writing to the context ID register. The contents of the caches and TLBs do not have to be invalidated after a fast context switch because they still hold valid address tags. From two to five instructions can be fetched with the old process ID after the MCR that writes to the process ID:

Example 4-4 Changing the process ID and performing a fast context switch

```

{procID = 0}
MOV r0, #1; Fetched with procID = 0
MCR p15,0,r0,c13,c0,0           ; Fetched with procID = 0
A0    (any instruction)         ; Fetched with procID = 0
A1    (any instruction)         ; Fetched with procID = 0
A2    (any instruction)         ; Fetched with procID = 0/1
A3    (any instruction)         ; Fetched with procID = 0/1
A4    (any instruction)         ; Fetched with procID = 0/1
A5    (any instruction)         ; Fetched with procID = 1

```

4.2.15 CP15 c14

Reading or writing CP15 register c14 is Undefined.

4.2.16 CP15 c15

CP15 register c15 is used for test purposes. Reading or writing CP15 c15 in normal operation is Unpredictable.

The functions of CP15 register c15 are described in:

- *PLL Configuration Register*
- *Power Manager Status Register* on page 4-28
- *Power Manager Receive Data Register* on page 4-29
- *Power Manager Transmit Data Register* on page 4-30
- *Transmission protocol* on page 4-31
- *Control Register 2* on page 4-32.

PLL Configuration Register

The PLL configuration register is for reprogramming the core clock frequency or AHB bus frequency. The register has a defined reset value as shown in Figure 4-19 on page 4-27. Refer to Chapter 15 *Clock Generator* for more details.

Table 4-25 shows the instructions for using the PLL configuration register.

Table 4-25 PLL Configuration Register instructions

Operation	Instruction
Read status	MRC p15, 0, Rd, c15, c12, 0
Write configuration	MCR p15, 0, Rd, c15, c12, 0

Figure 4-19 shows the bit fields of the PLL configuration register.

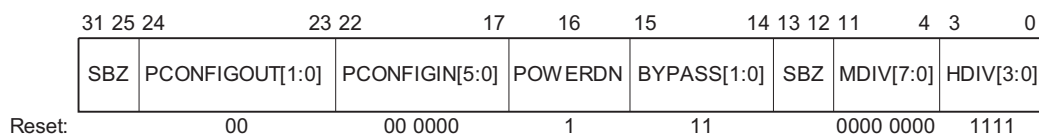


Figure 4-19 PLL Configuration Register

Table 4-26 describes the bit fields of the PLL configuration register.

Table 4-26 Encoding of the PLL Configuration Register

Bits	Meaning
[31:25]	Should Be Zero.
[24:23]	Bit 24 is for a partner-defined PLL function. Bit 23 is for a lock-detect signal.
[22:17]	Partner-specific PLL functions.
[16]	POWERDN PLL draws only minimum leakage current due to VCO being clamped. Lock is lost.
[15:14]	BYPASS [1:0] Controls the selects for the GCLK , HCLK , and VCO multiplexors. See Chapter 15 <i>Clock Generator</i> .
[13:12]	Should Be Zero.
[11:4]	PLL feedback divider (MCLK) MDIV[7:0] .
[3:0]	AHB clock divider (HCLK) HDIV[3:0] .

Power Manager Status Register

The Power Manager Status Register, PMSR, contains the version number of the power manager. It also indicates when the receive channel is available to check the last state of the system, and when the transmit channel is available to send new data.

Table 4-27 shows the instructions for using PMSR.

Table 4-27 Power Manager Status instructions

Operation	Instruction
Read status	MRC p15, 0, Rd, c15, c14, 0
Check receive channel	MRC p15, 0, Rd, c15, c14, 1
Write transmit channel	MCR p15, 0, Rn, c15, c14, 1

Figure 4-20 shows the PMSR bit fields.



Figure 4-20 Power Manager Status Register

Table 4-28 describes the PMSR bit fields.

Table 4-28 Encoding of the Power Manager Status Register

Bits	Meaning
[31:28]	Version = 0x0001
[27:2]	Should Be Zero
[1]	Denotes transmit channel is ready: 1 = Idle 0 = Busy
[0]	Denotes receive channel is full: 1 = Full 0 = Empty

Power Manager Receive Data Register

When the R flag in power manager status register is set, valid data can be read from the Power Manager Receive Data Register, PMRDR. An acknowledgement is sent to the power manager to indicate data acceptance. When the R flag is clear, reading PMRDR is Unpredictable. Writing to PMRDR is Unpredictable. Figure 4-21 shows the PMRDR bit fields.

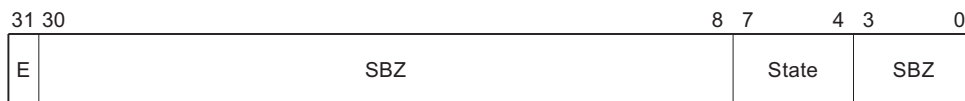


Figure 4-21 Power Manager Receive Data Register

Table 4-29 describes the PMRDR bit fields.

Table 4-29 Encoding of the Power Manager Receive Data Register

Bits	Meaning
[31]	Emulation flag. When exiting a reset sequence, E reflects the last programmed state of the system. 1 = power manager issued a command in emulation mode 0 = power manager issued a command in normal mode
[30:8]	Reserved. Reads as zero.
[7:4]	System power state. When exiting a reset sequence, this field reflects the last programmed state of the system. 1111 = TURBO 1110 = NORMAL 110x = SLOW 100x = IDLE 01xx = NAP 0011 = SLEEP 0010 = COMA 0001 = HIBERNATE 0000 = OFF
[3:0]	Reserved. Reads as zero.

Power Manager Transmit Data Register

When the W flag in power manager status register is set, new data can be written to the Power Manager Transmit Data Register, PMTDR. An acknowledgement following the write is sent to the power manager to indicate that new data is available. Writing to PMTDR clears W. Writing to PMTDR when W is clear is Unpredictable. Reading PMTDR is Unpredictable. Figure 4-22 shows the PMTDR bit fields.

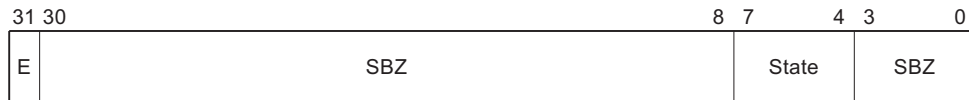


Figure 4-22 Power Manager Transmit Data Register

Table 4-30 describes the PMTDR bit fields.

Table 4-30 Encoding of the Power Manager Transmit Data Register

Bits	Name	Meaning
[31]	E	1 = power manager issued a command in emulation mode 0 = power manager issued a command in normal mode.
[30:8]	-	Should Be Zero.
[7:4]	State	System power state. When exiting a reset sequence, this value reflects the last programmed state of the system. 1111 = TURBO 1110 = NORMAL 110x = SLOW 100x = IDLE 01xx = NAP 0011 = SLEEP 0010 = COMA 0001 = HIBERNATE 0000 = OFF.
[3:0]	-	Should Be Zero.

Transmission protocol

When issuing commands to the power manager, a specific protocol must be followed:

1. By reading the W and R flags, software checks to see that both transmit data and receive data bit fields are empty.
2. When transmitting, software must write a command to the transmit data register. This clears the W flag. Hardware then performs a handshake with the power manager, waiting for acceptance of the command using a double-ended handshake.
3. When the handshake for the transmit data is done, hardware sets the W flag.

When receiving data, software must wait until the R flag is set. When set, new data is valid in the receive data register.

Data Transmit Code

To transmit data to the power manager, software must always perform the code sequence shown below. The command is sent using ARM register r1, while ARM register r0 reflects the status register contents:

tx_command:

```
MRC CP15, 0, R0, c15, c14, 0 ; check for outstanding commands
TST R0, #W_flag ; 'W' flag clear indicates active command
BNE tx_command ; if command active, loop again
MCR CP15, 0, R1, c15, c14, 1 ; write new command to controller
```

———— Note ————

The W flag is polled until it is one. When W is set, the command can be sent to the power manager.

Data Receive Code

To wait until data has been received in the receive data register, software must always perform the code sequence shown below. The command is received into ARM register r1, while ARM register r0 reflects the status register contents:

rx_status:

```
MRC CP15, 0, R0, c15, c14, 0 ; check for incoming data
TST R0, #R_flag ; 'R' flag clear indicates no data
BNE rx_status ; if no data, loop again
MRC CP15, 0, R0, c15, c14, 1 ; read in 'previous-state'
```

Note

The R flag is polled until it is cleared. When R is cleared, the command can be read.

Control Register 2

The read/write control register 2 is primarily useful when in debug mode. When not in debug mode, it is also useful to define the behavior of accesses when the caches are on and the DMMU is off.

Table 4-31 shows the instructions for using control register 2.

Table 4-31 Control Register 2 instructions

Operation	Instruction
Read status	MRC p15, 0, Rd, c15, c11, 0
Write	MCR p15, 0, Rd, c15, c11, 0

Figure 4-23 shows control register 2.

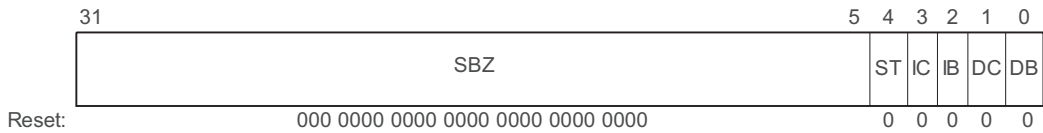


Figure 4-23 Control Register 2

Table 4-32 describes the bit fields of control register 2.

Table 4-32 Encoding of Control Register 2

Bits	Meaning
[31:5]	Should Be Zero.
[4]	ST, CP15 soft TLB enable bit. Reset clears ST. 1 = soft TLB enabled 0 = soft TLB disabled.
[3]	IC, CP15 instruction cachable bit, used only in debug mode with the ICache on. Reset clears IC. 1 = instructions cachable 0 = instructions not cachable.
2	IB, CP15 instruction bufferable bit. Used only in debug mode with the ICache on. Included for compatability with the B bit in the MMU descriptors. Reset clears IB. 1 = instructions bufferable 0 = instructions not bufferable.
1	DC, CP15 data cachable bit. Used in debug mode or when the DMMU is off and the DCache is on. Reset clears DC. 1 = data cachable 0 = data not cachable.
0	DB, CP15 data bufferable bit; used in debug mode or when the DMMU is off and the DCache is on. Reset clears DB. 1 = data bufferable 0 = data not bufferable.

Chapter 5

Memory Management Units

This chapter describes the ARMv5 *Memory Management Units* (MMUs). It contains the following sections:

- *About the MMUs* on page 5-2
- *MMU software-accessible registers* on page 5-3
- *Address translation* on page 5-5
- *MMU memory access control* on page 5-24
- *MMU cachable and bufferable information* on page 5-26
- *MMU and write buffer* on page 5-27
- *MMU aborts* on page 5-28
- *MMU fault checking sequence* on page 5-29
- *CPU aborts on MMU faults* on page 5-32
- *Fault priority* on page 5-33
- *External aborts* on page 5-34
- *Interaction of the MMU, caches, and write buffer* on page 5-36
- *Soft page table support* on page 5-37.

5.1 About the MMUs

The MMUs control external memory accesses and translate *Virtual Addresses* (VAs) to *Physical Addresses* (PAs).

The *Instruction MMU* (IMMU) and *Data MMU* (DMMU) provide address translation and access permission checks for the instruction and data ports of the integer unit. They control the descriptor fetch hardware that accesses page table descriptors in main memory. To support sections and pages, there are two levels of page tables. The finished VA-to-PA translations are put into separate instruction-side and data-side *Translation Lookaside Buffers* (TLBs).

MMU features include:

- standard MMU mapping sizes, domains, and access protection
- 1KB, 4KB, 64KB, and 1MB mapping sizes
- access permissions for 1MB sections
- separate access permissions for one-quarter page subpages of 64KB large pages and 4KB small pages
- 16 domains
- separate 64-entry instruction and data TLBs
- independent lockdown of instruction and data TLBs
- hardware page table descriptor fetches
- round-robin replacement algorithm
- support for soft page tables.

5.2 MMU software-accessible registers

The CP15 registers shown in Table 5-1, along with the page table descriptors stored in memory, control MMU operation.

Table 5-1 CP15 register MMU functions

CP15 register	Bits	Register description
c1 Control register 1	M	Bit 0, MMU enable bit: 1 = IMMU and DMMU enabled 0 = IMMU and DMMU disabled.
	A	Bit 1, address alignment fault checking enable bit: 1 = fault checking of address alignment enabled 0 = fault checking of address alignment disabled.
	S	Bit 8, system protection enable bit: 1 = IMMU and DMMU protection enabled 0 = IMMU and DMMU protection disabled.
	R	Bit 9, ROM protection enable bit: 1 = ROM protection enabled 0 = ROM protection disabled.
c2 Translation table base register	[31:14]	Holds PA of base of translation table in main memory. Base must reside on a 16KB boundary and is common to both IMMU and DMMU.
c3 Domain access control register	[31:0]	Has 16 2-bit fields. Each field defines the access control attributes for one of 16 domains (D15-D0). See Table 5-5 on page 5-24.
c5 Fault status register	[31:8]	Indicates domain number and cause of Data Abort. Should Be Zero.
	[7:4]	Indicate domain (D15-D0) in which fault occurred.
	[3:0]	Indicate type of access attempted. See Table 5-8 on page 5-33.
c6 Fault address register	[31:0]	Holds VA associated with access that caused Abort. See <i>CP15 c6, Fault Address Register</i> on page 4-15 for FAR access instructions. See Table 5-8 on page 5-33 for details of the address stored for each type of fault.

Table 5-1 CP15 register MMU functions (continued)

CP15 register	Bits	Register description
c8 TLB operations register	[31:5]	Writing to c8 causes the MMU to perform TLB maintenance operations, invalidating one or all unpreserved TLB entries.
c10 TLB lockdown register	[31:20], [0]	Allows specific page table entries to be locked into a TLB and the TLB victim counter to be read/written. Locking entries in a TLB guarantees that accesses to the locked page or section can proceed without incurring the time penalty of a TLB miss. This enables the execution latency for time-critical pieces of code such as IRQ handlers to be minimized.
c15 Control register 2	[4]	Allows the MMU to be configured for soft TLB support.

Note

All the CP15 MMU registers, except CP15 c8, contain state and can be read using MRC instructions and written to using MCR instructions. CP15 c5 and CP15 c6 are also written by the MMU. Reading CP15 c8 is Unpredictable.

5.3 Address translation

The address translation process begins when the integer unit requests access to an address that has no VA-to-PA translation in the TLB, causing a TLB miss. The MMU then fetches a page table descriptor.

5.3.1 TLBs

Each TLB caches 64 translated entries. If, during a memory access, the TLB contains a translated entry for the VA, the MMU reads the protection data to determine if the access is permitted:

- If the access is permitted, and off-chip access is required, the MMU produces the PA.
- If the access is permitted, and off-chip access is not required, the cache services the access.
- If the access is not permitted, the MMU signals the CPU to abort.

If a TLB miss occurs, the page table descriptor fetch hardware retrieves the translation information from a translation table in main memory. The retrieved information is written into the TLB, possibly overwriting an existing value.

The entry to be written is usually chosen by cycling sequentially through the TLB locations. To enable use of TLB locking features, the location to be written can be specified using the TLB lockdown register, CP15 c10.

When the MMU is turned off, as happens at reset, no address mapping occurs, and all regions are marked as noncacheable and nonbufferable.

5.3.2 Page table descriptor fetches

A page table descriptor fetch occurs whenever there is a TLB miss. The descriptor fetch begins with the formation of a level 1 descriptor.

———— **Note** —————

If the DMMU is performing an external memory operation for the load/store unit, the write buffer is emptied before the descriptor fetch. This guarantees that memory remains coherent. The DMMU then performs the operation as noncacheable and nonbufferable.

IMMU activity does not cause the write buffer to be emptied.

5.3.3 Translation routes for sections and pages

The MMU translates VAs from the integer unit to PAs for an external memory access. The two types of memory blocks, sections and pages, require a specific translation process to occur.

Figure 5-1 on page 5-7 shows the translation process. A section requires only a level 1 descriptor fetch. A page requires both a level 1 and level 2 descriptor fetch.

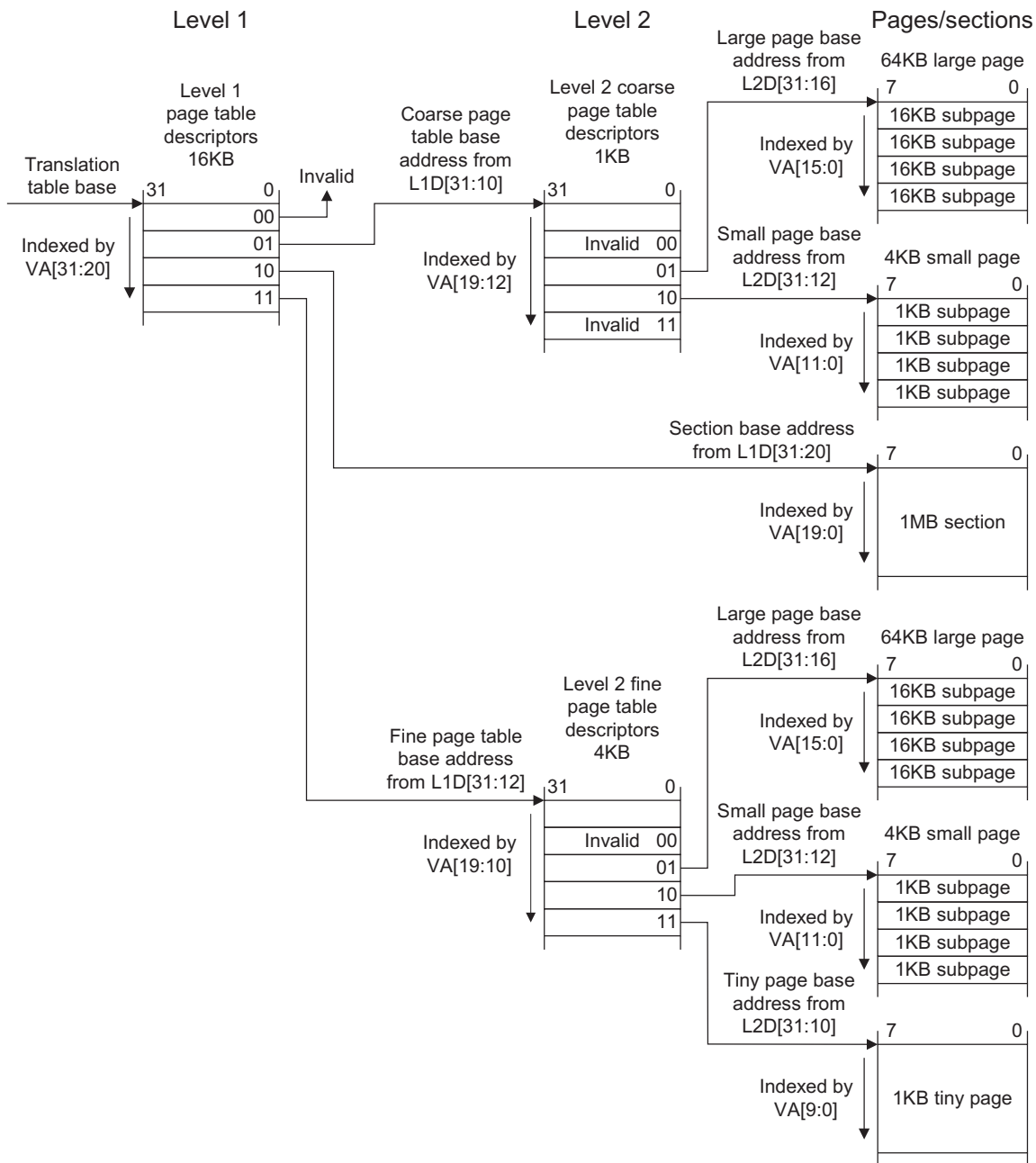


Figure 5-1 Address translation overview

5.3.4 Level 1 descriptor address

Figure 5-2 shows how the MMU uses the translation table base field in CP15 c2 and the VA from the integer unit to create the level 1 descriptor address.

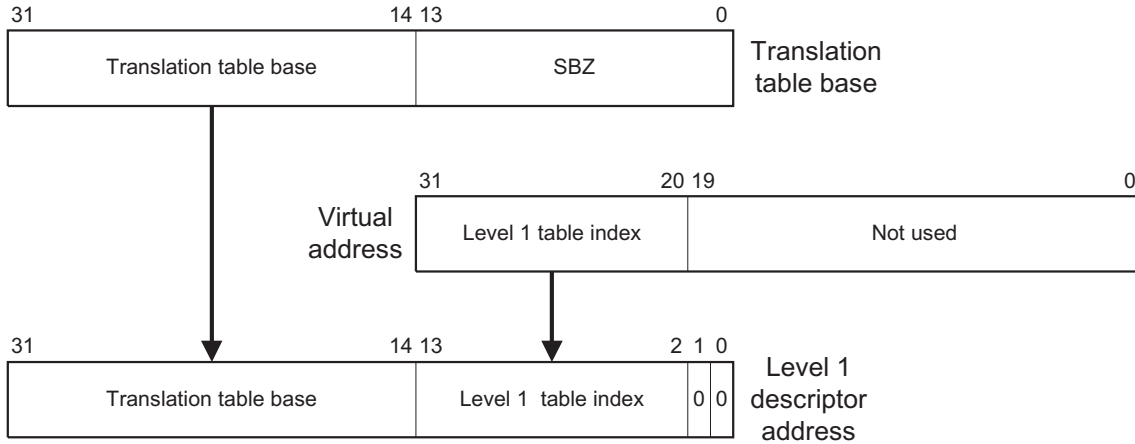


Figure 5-2 Translating a level 1 descriptor address

5.3.5 Level 1 page table descriptors

The level 1 descriptor indicates whether the access is:

- a translation fault
- an access to a level 2 coarse page table
- an access to a 1MB section of external memory
- an access to a level 2 fine page table.

Bits [1:0] of the level 1 descriptor determine the type of access. Figure 5-3 on page 5-9 shows the level 1 descriptor formats for the three access types.

	31		20	19		12	11	10	9	8		5	4	3	2	1	0
Translation fault	Ignore															0	0
Coarse page table	Level 2 coarse page table base address										SBZ	Domain selector		1	SBZ	0	1
1MB section	Section base address					SBZ	AP	SBZ	Domain selector		1	C	B	1	0		
Fine page table	Level 2 fine page table base address							SBZ	Domain selector		1	SBZ	1	1			

Figure 5-3 Level 1 descriptor formats

Using the level 1 descriptor address, the MMU makes a request to external memory. This returns the level 1 descriptor. Bits [1:0] of the level 1 descriptor indicate the access type as Table 5-2 shows.

Table 5-2 Access types in a level 1 descriptor

Bits [1:0]	Access type
00	Translation fault
01	Coarse page table base address
10	Section base address
11	Fine page table base address

Level 1 translation fault

If bits [1:0] of the level 1 descriptor are 00, a translation fault is generated. This causes either a Prefetch Abort or Data Abort in the integer unit. A Prefetch Abort occurs in the IMMU. A Data Abort occurs in the DMMU.

Level 1 coarse page table address

If bits [1:0] of the level 1 descriptor are 01, then a descriptor fetch from a coarse page table is required. Figure 5-6 on page 5-12 shows how the MMU generates a coarse page table address.

Level 1 section base address

If bits [1:0] of the level 1 descriptor are 10, a request to access a 1MB memory section is requested. Figure 5-4 shows the translation process for a 1MB section.

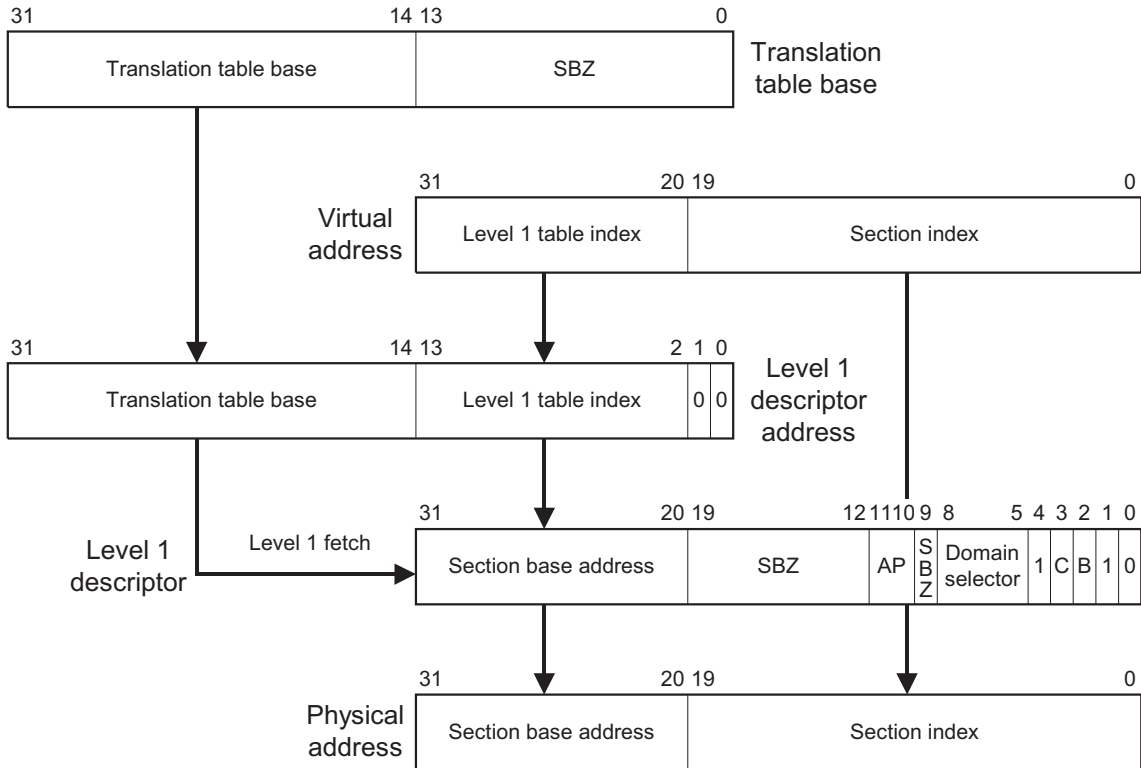


Figure 5-4 Translating a section address

Following the level 1 descriptor translation, the the MMU uses the PA to transfer the requested data between external memory and the integer unit. This is done only after the domain and access permission checks are performed on the level 1 descriptor for the section. These checks are described in *MMU memory access control* on page 5-24.

Level 1 fine page table base address

If bits [1:0] of the level 1 descriptor are 11, then a descriptor fetch from a fine page table is required. This is shown in Figure 5-9 on page 5-17.

5.3.6 Level 2 descriptor

If the level 1 descriptor points to a page table, the MMU determines the page table type, coarse or fine, and fetches a level 2 descriptor. The level 2 descriptor indicates whether the access is:

- a translation fault
- an access from a coarse page table to a large page with 64K 8-bit entries
- an access from a coarse page table to a small page with 4K 8-bit entries
- an access from a fine page table to a large page, a small page, or a tiny page with 1K 8-bit entries.

Figure 5-5 shows the level 2 descriptor formats for selecting page types.

	31		16	15		12	11	10	9	8	7	6	5	4	3	2	1	0	
Translation fault	Ignore																0	0	
64KB large page	Large page base address				SBZ	AP3	AP2	AP1	AP0	C	B	0	1						
4KB small page	Small page base address					AP3	AP2	AP1	AP0	C	B	1	0						
1KB tiny page	Tiny page base address							SBZ	AP	C	B	1	1						

Figure 5-5 Level 2 descriptor formats

Bits [1:0] of the level 2 descriptor indicate the page type. A large page can be divided four 16KB subpages with different access permissions. Bits [15:14] of the VA page index select the subpages of a large page.

A small page can be divided into four 1KB subpages with different access permissions. Bits[11:10] of the VA page index select the subpages of a small page.

Level 2 coarse page table descriptor fetch

When the level 1 descriptor bits [1:0] indicate a descriptor fetch from a coarse page table is required, the MMU requests the address of the level 2 coarse page table from external memory. Figure 5-6 on page 5-12 shows how the address is generated.

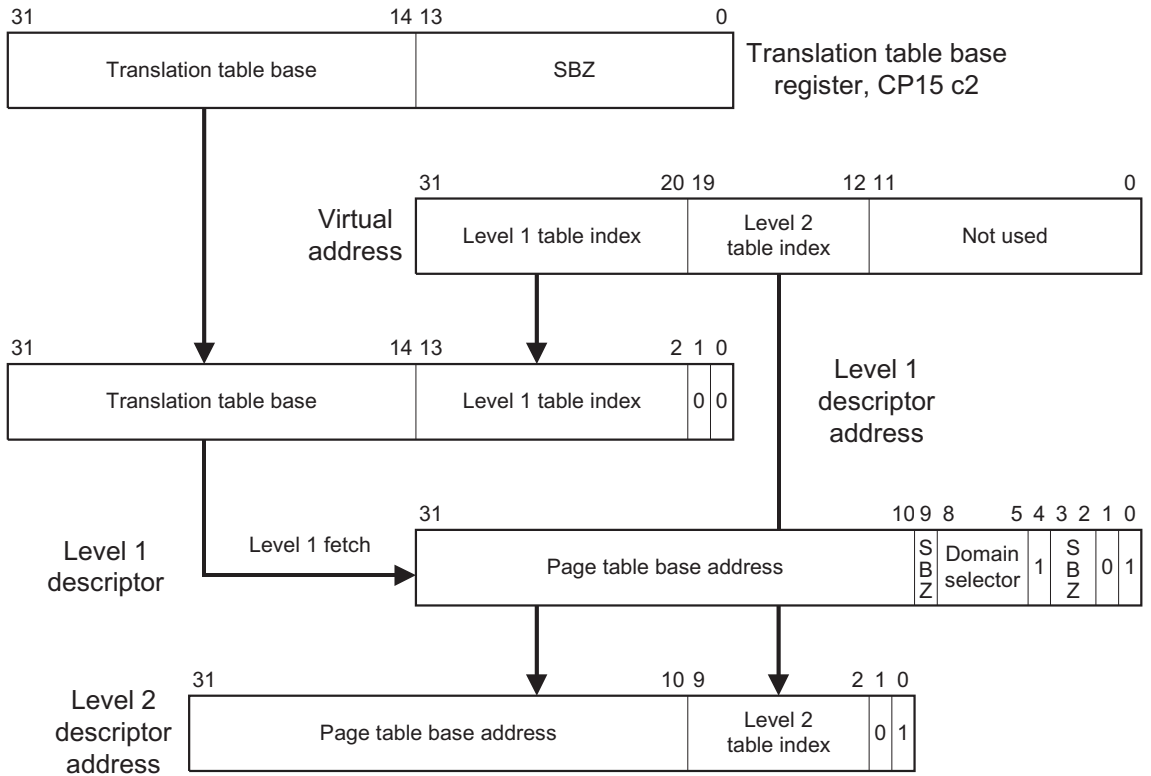


Figure 5-6 Translating a coarse page table address

When the coarse page table address is generated, a request is made to external memory for the level 2 coarse page table descriptor. Bits [1:0] of the level 2 coarse page table descriptor indicate the access type as shown in Table 5-3.

Table 5-3 Access types in a coarse page table descriptor

Bits[1:0]	Access type
00	Translation fault
01	64KB large page base address
10	4KB small page base address
11	Translation fault

Level 2 coarse translation fault

If bits [1:0] of the level 2 coarse page table descriptor are 00 or 11, then a translation fault is generated. This generates an abort to the integer unit, either a Prefetch Abort for the instruction side or a Data Abort for the data side.

Level 2 coarse large page base address

If bits [1:0] of the level 2 coarse page table descriptor are 01, then a descriptor fetch from a coarse large page table is required. Figure 5-7 on page 5-14 shows the translation process for a 64KB large page or a 16KB subpage of a large page.

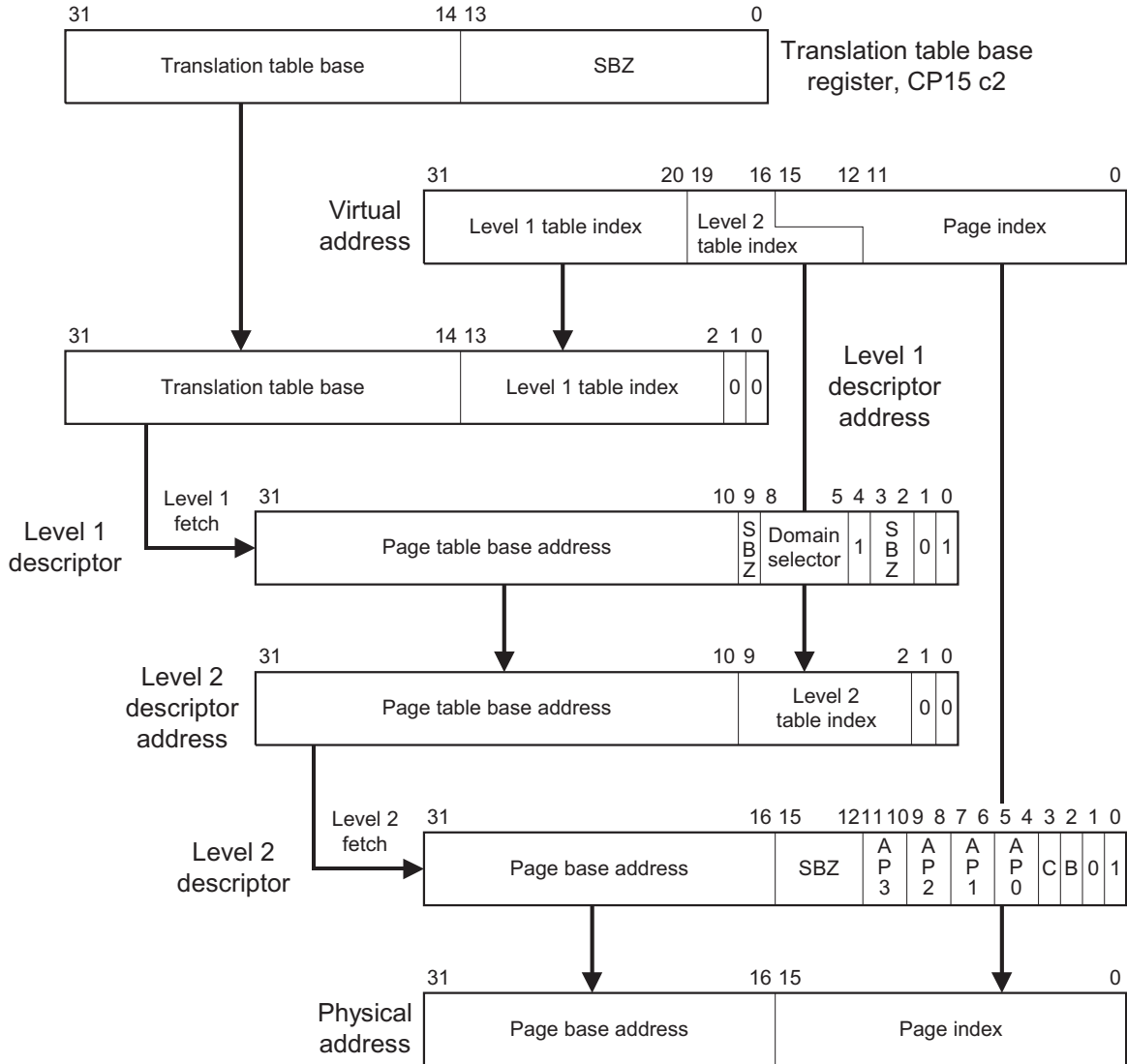


Figure 5-7 Translating a large page or subpage address from a coarse page table

The 64KB large page is generated by setting all of the AP bit pairs to the same values, AP3 = AP2 = AP1 = AP0. If any one of the pairs is different, then the 64KB large page is converted into four 16KB subpages.

Level 2 coarse small page base address

If bits [1:0] of the level 2 coarse page table descriptor are 10, then a descriptor fetch from a coarse small page table is required. Figure 5-8 on page 5-16 shows the translation process for a 4KB small page or a 1KB subpage of a small page.

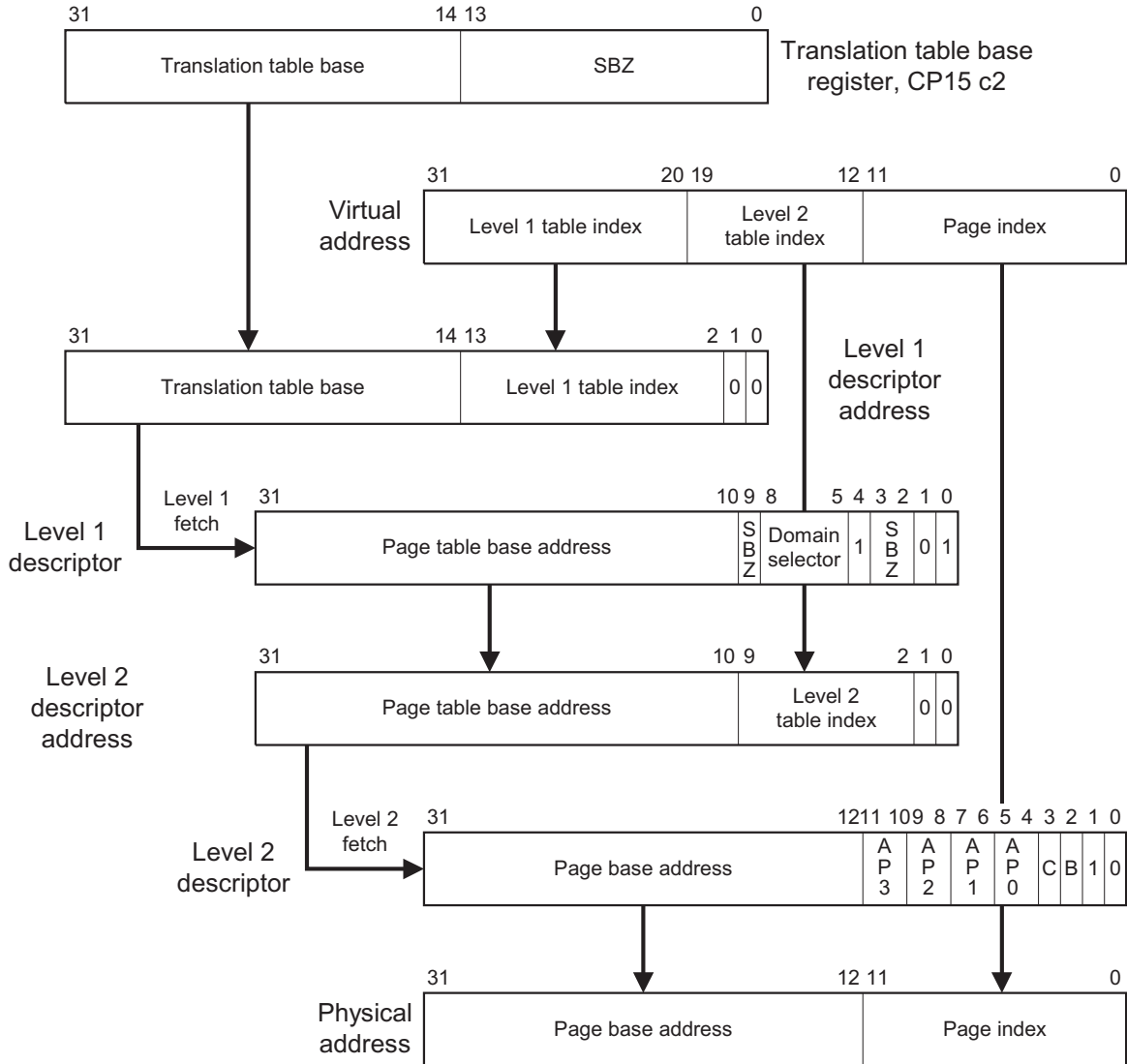


Figure 5-8 Translating a small page or subpage address from a coarse page table

The 4KB small page is generated by setting all of the AP bit pairs to the same values, AP3 = AP2 = AP1 = AP0. If any one of the pairs are different, then the 4KB small page is converted into four 1KB small page subpages.

Level 2 fine page table descriptor fetch

When the level 1 descriptor bits [1:0] indicate that a descriptor fetch from a fine page table is required, the MMU requests the level 2 fine page table address from external memory. Figure 5-9 shows how the address is generated.

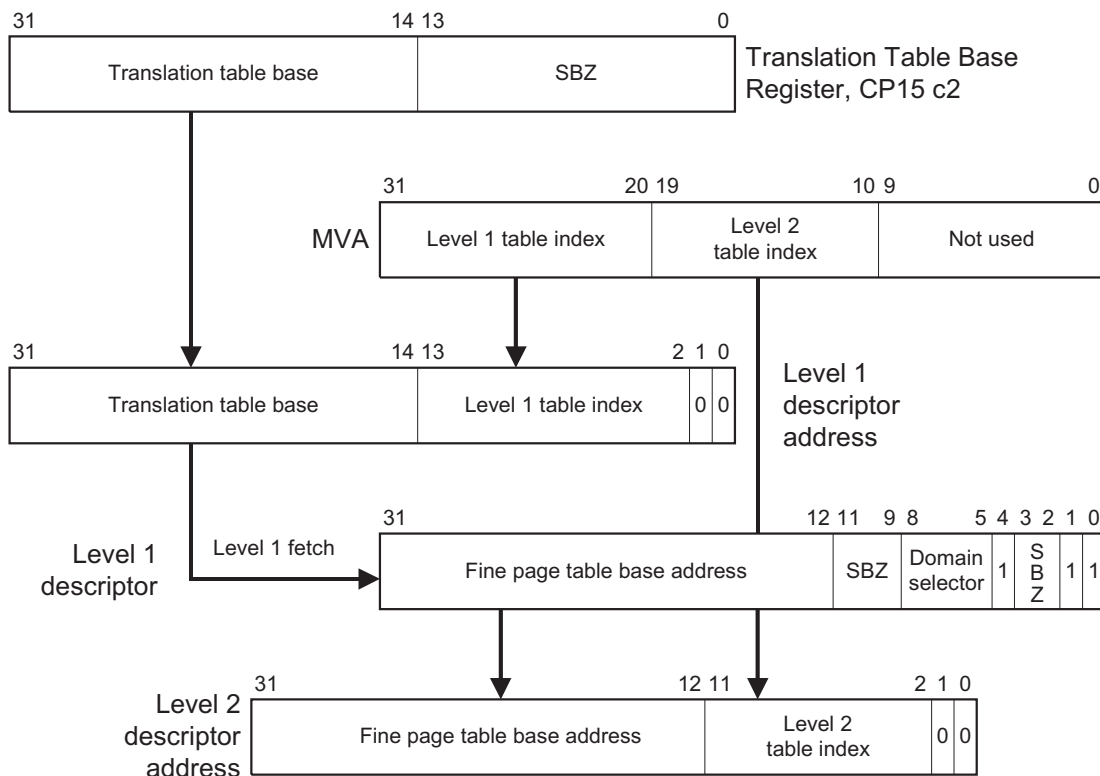


Figure 5-9 Translating a fine page table address

When the fine page table address is generated, a request is made to external memory for the level 2 fine page table descriptor. Bits [1:0] of the level 2 fine page table descriptor indicate the access type as shown in Table 5-4.

Table 5-4 Access types in a fine page table descriptor

Bits [1:0]	Access type
00	Translation fault
01	Large page table base address
10	Small page base address
11	Tiny page table base address

Level 2 fine translation fault

If bits [1:0] of the level 2 fine page table descriptor are 00, then a translation fault is generated. This causes either a Prefetch Abort or a Data Abort in the integer unit. A Prefetch Abort occurs on the instruction side, while a Data Abort occurs on the data side.

Level 2 fine large page base address

If bits [1:0] of the level 2 fine page table descriptor are 01, then a descriptor fetch from a fine large page table is required. Figure 5-10 on page 5-19 shows the translation process for a 64KB large page or a 16KB subpage of a large page.

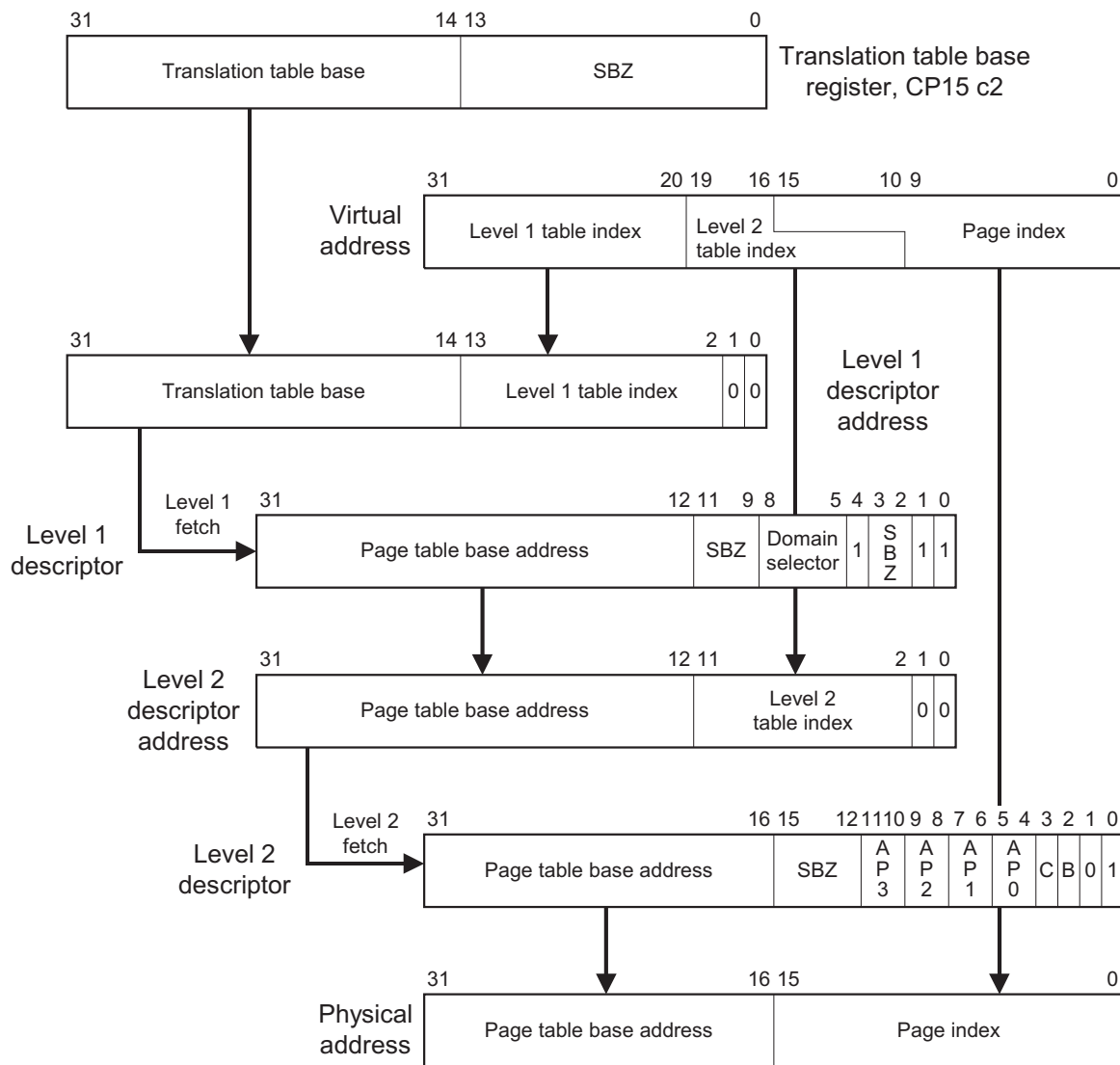


Figure 5-10 Translating a large page or subpage address from a fine page table

The 64KB large page is generated by setting all of the AP bit pairs to the same values, AP3 = AP2 = AP1 = AP0. If any of the pairs is different, then the 64KB large page is converted into four 16KB subpages.

Level 2 fine small page base address

If bits [1:0] of the level 2 fine page table descriptor are 10, then a descriptor fetch from a fine small page table is required. Figure 5-11 on page 5-21 shows the translation process for a 4KB small page or a 1KB subpage of a small page.

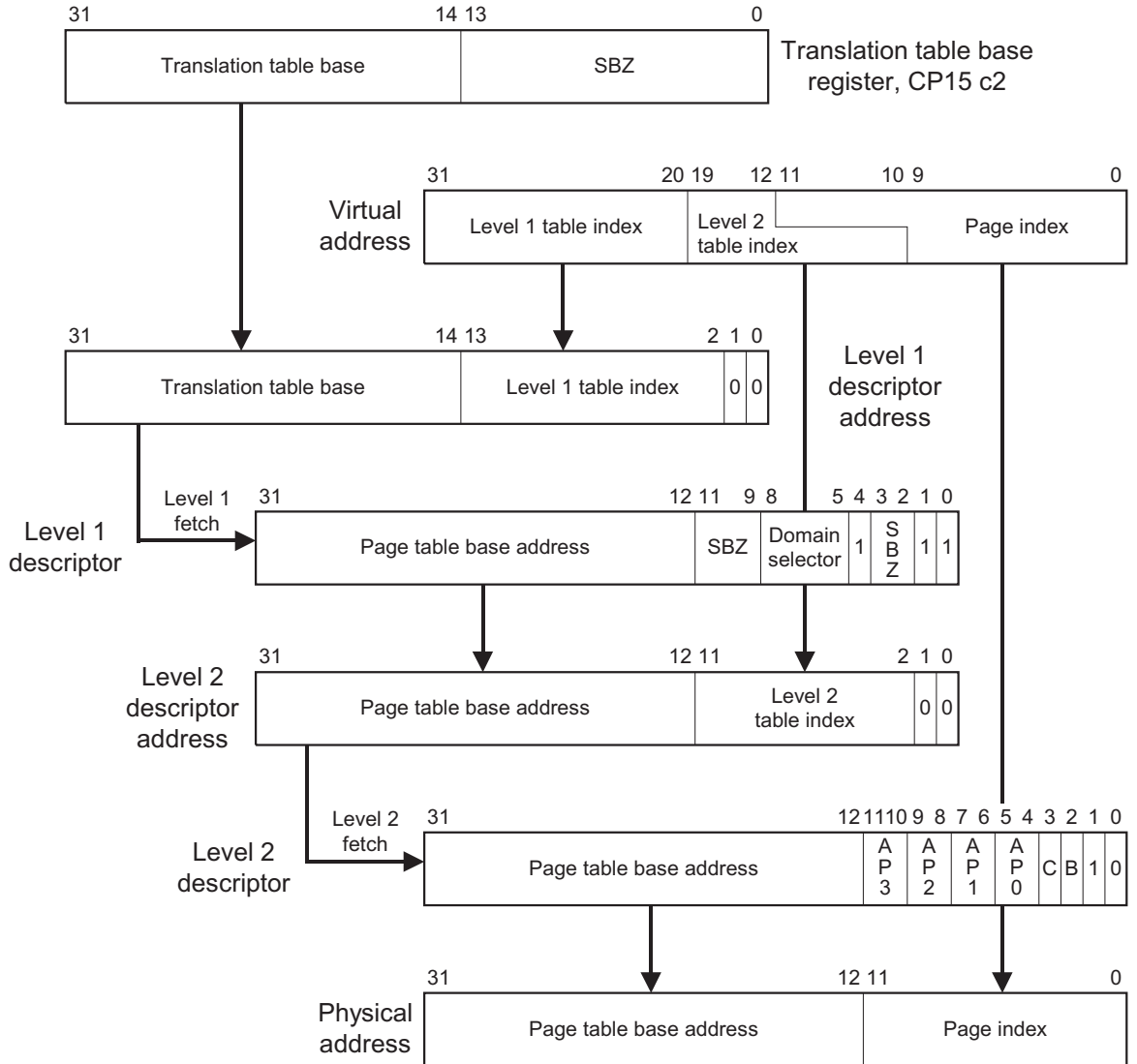


Figure 5-11 Translating a small page or subpage address from a fine page table

Level 2 fine tiny page base address

If bits [1:0] of the level 2 fine page table descriptor are 11, then a descriptor fetch from a fine tiny page table is required. Figure 5-12 on page 5-23 shows the translation process for a 1KB tiny page.

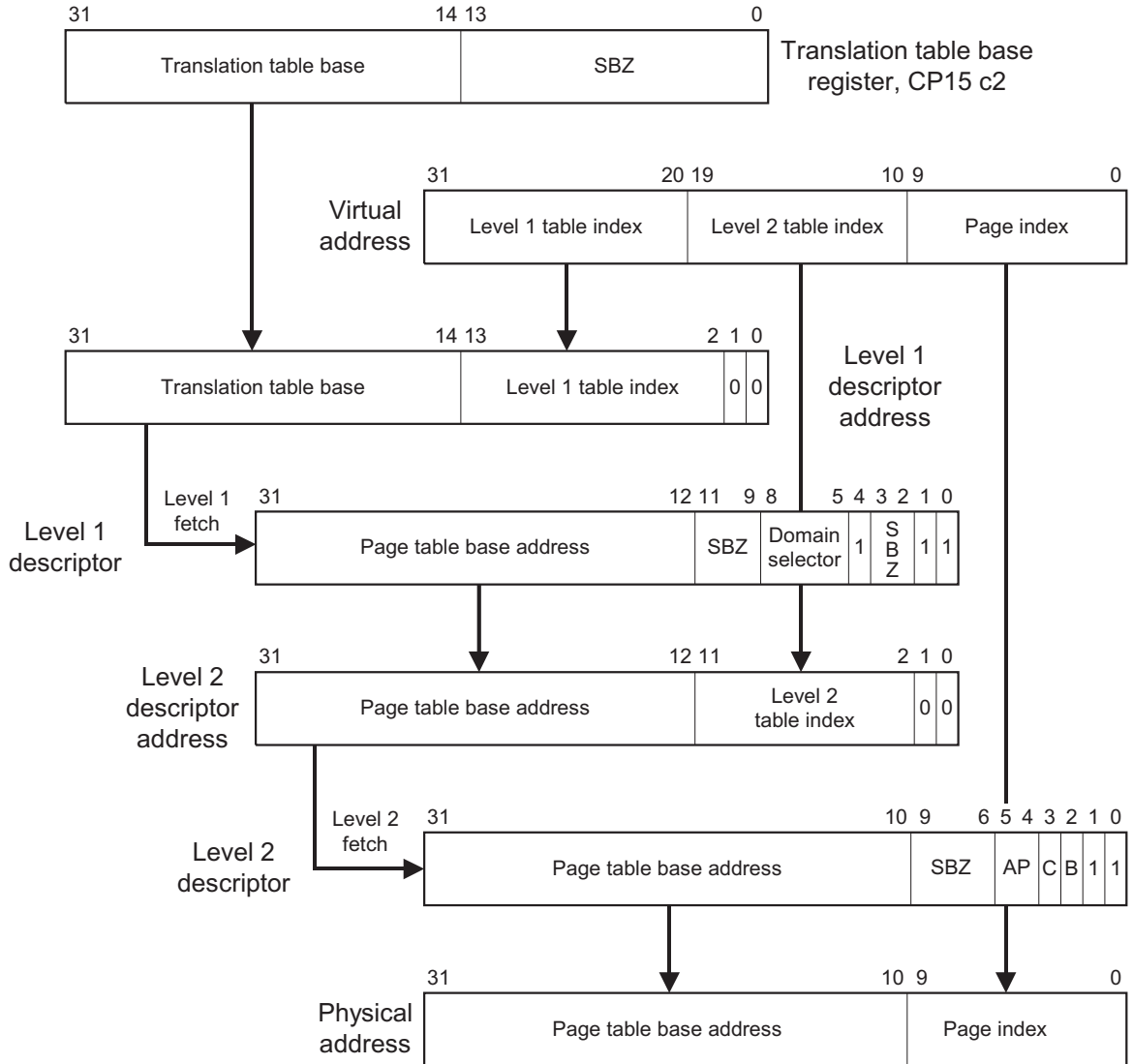


Figure 5-12 Translating a tiny page address

5.4 MMU memory access control

Memory domains support multiuser operating systems. All regions of memory have an associated domain. Domains are the primary memory access control mechanism and define the conditions in which an access can proceed. Each domain determines whether:

- access is qualified to proceed as shown in Table 5-6 on page 5-25
- access is unconditionally enabled to proceed
- access is unconditionally aborted.

In the latter two cases, the access permission attributes are ignored. There are 16 domains, D15-D0, that are configured in the domain access control register.

The domain definition provides access for two types of users, manager and client. The two-bit D15-D0 fields in CP15 c3 control access to both the IMMU and the DMMU domains. Table 5-5 shows the encoding for of the domain access control fields.

Table 5-5 Domain access encoding

D15-D0	User	Notes
00	No access	Access generates a domain fault.
01	Client	Access permissions are checked.
10	Reserved	Behaves as a <i>no access</i> domain.
11	Manager	Access permissions are not checked.

A manager access has to be checked only against the access permissions for the domain. A client access has to be checked against the access permissions for the domain and the system protection bit, S, and the ROM protection bit, R, in CP15 c1. Table 5-6 on page 5-25 shows the effect of the S and R bits.

Table 5-6 S and R bit encoding

D15-D0	S	R	Supervisor permissions	User permissions	Notes
00	0	0	No access	No access	Any access generates a permission fault.
00	1	0	Read-only	No access	Supervisor read-only permitted.
00	0	1	Read-only	Read-only	Writing generates a permission fault.
00	1	1	Reserved	-	-
01	-	-	Read/write	No access	Supervisor mode access only
10	-	-	Read/write	Read-only	Writes in User mode cause permission fault.
11	-	-	Read/write	Read/write	All access types permitted in both modes.
-	1	1	Reserved	-	-

5.5 MMU cachable and bufferable information

The *Cachable* (C) and *Bufferable* (B) bits in the level 1 and level 2 descriptors control the operation of memory accesses to external memory. Table 5-7 indicates how the MMU and cache interpret the C and B bits.

Table 5-7 C and B bit access control

C	B	Notes
0	0	Uncached, unbuffered
0	1	Uncached, buffered
1	0	Write-through cached, buffered
1	1	Write-back cached, buffered

Refer to *Cache coherence* on page 6-15 for information on how cache coherence is maintained.

5.6 MMU and write buffer

During any descriptor fetch, the IMMU or DMMU has access to external memory. The integer unit is stalled during any descriptor fetch.

Before a DMMU descriptor fetch, the write buffer has to be emptied to preserve memory coherency. If the write buffer contains any page table entries that have been modified, those entries are forced to external memory as a result of the descriptor fetch.

When either the IMMU or DMMU contains valid TLB entries that are being modified, these TLB entries must be invalidated before the new section or page is accessed. This also applies to any data that resides in the ICache or DCache. The ICache lines must be invalidated, and the DCache line or lines must be cleaned and invalidated (see *Cache coherence* on page 6-15).

5.7 MMU aborts

During any translation process, the integer unit stops executing instructions whenever an MMU fault is generated or an external abort occurs:

- If the abort is from the IMMU, a Prefetch Abort is indicated to the integer unit.
- If the abort is from the DMMU, then a Data Abort is indicated to the integer unit.

The fault status and fault address registers in CP15 log both the status and address for any fault that occurs.

In the case of an external abort, the *Bus Interface Unit* (BIU) ignores the abort unless one of the following is true:

- it was caused by a write to a *NonCached NonBuffered* (NCNB) region
- it was caused by a read from a *Noncached Buffered* (NCB) region
- it occurred during a descriptor fetch.

5.8 MMU fault checking sequence

During the processing of a section or page, the MMU behaves differently while it is checking for faults. This section describes the following conditions:

- *Alignment fault*
- *Translation fault*
- *Domain fault* on page 5-31
- *Permission fault* on page 5-31.

Figure 5-13 on page 5-30 shows the fault checking sequence.

5.8.1 Alignment fault

An alignment fault occurs whenever the integer unit indicates a particular data memory access size and the address does not comply with that size. If $MAS[1:0] = 10$ indicating a 32-bit access, and the VA bits $[1:0] \neq 00$, then an alignment fault occurs. If $MAS[1:0] = 01$ indicating a 16-bit access, and the VA bit $0 \neq 0$, then an alignment fault occurs. No check is performed for $MAS[1:0] = 00$.

Alignment checks are performed with the MMU both on and off.

5.8.2 Translation fault

Two types of translation faults occur:

- section
- page.

A section translation fault results from an invalid level 1 descriptor. Bits $[1:0]$ of the descriptor are 00.

A page translation fault results from an invalid level 2 descriptor. Bits $[1:0]$ of the coarse page table descriptor are 00 or 11, or bits $[1:0]$ of the fine page table descriptor are 00.

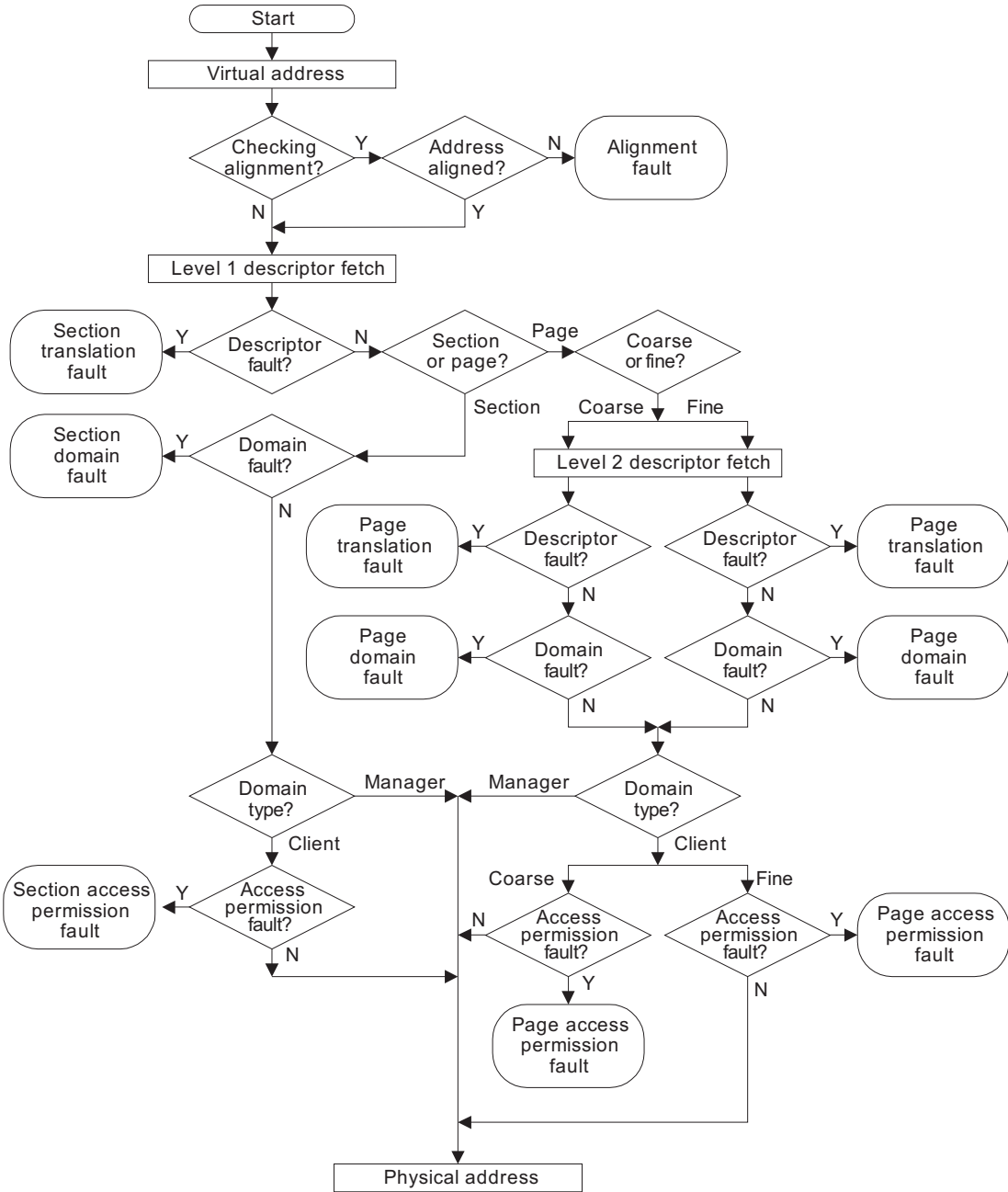


Figure 5-13 Fault checking flowchart

5.8.3 Domain fault

Three types of domain faults occur:

- section
- coarse page
- fine page.

For each type, the level 1 descriptor indicates which domain to select in the domain access control register, CP15 c3. If bit 0 of the selected domain is zero, indicating either *No access* or *Reserved*, then a domain fault occurs. A section domain fault occurs when the level 1 descriptor is returned. Both the coarse and fine page domain faults are checked whenever the level 2 descriptor is returned.

The MMU empties any unlocked TLB entry following a write to the domain access control register. To guarantee the behavior, all locked TLB entries must not modify their DACR entry. If the DACR entry is modified, it must be unlocked and invalidated.

5.8.4 Permission fault

There are three types of access permission faults:

- section
- coarse page
- fine page.

Whenever the domain indicates that a client has accessed a region of memory, an access permission check follows. If the access does not comply with the access permission table, then a fault corresponding to the access type occurs. A section permission fault check occurs when the level 1 descriptor is returned and is designated as a client. Both the coarse and fine page permission faults are checked whenever the level 2 descriptor is returned and is designated as a client.

5.9 CPU aborts on MMU faults

The MMU generates an abort on the following types of faults:

- alignment faults (data accesses only)
- translation faults
- domain faults
- permission faults.

In addition, an external abort can be raised on some types of external data access.

Alignment fault checking is enabled by the A bit in CP15 c1. Alignment fault checking is independent of the MMU being enabled. Translation, domain, and permission faults are only generated when the MMU is enabled.

The access control mechanisms of the MMU detect the conditions that produce these faults. If a fault is detected as the result of a memory access, the MMU aborts the access and signals the fault condition to the CPU. The MMU retains status and address information about faults generated by the data accesses in the fault status register and fault address register. The MMU does not retain status about faults generated by instruction fetches.

An access violation for a given memory access inhibits any corresponding external access, with an abort returned to the integer unit.

5.9.1 Fault address registers and fault status registers

Both the IMMU and DMMU have a fault address register and a fault status register. In the IMMU, a Prefetch Abort updates bits [3:0] of the IMMU fault status register and is pipelined to the Execute stage. This is only used if the Prefetch Abort exception is taken.

The DMMU updates bits [3:0] of the DMMU fault status register with the domain number. It also loads the VA of the Data Abort into the data fault address register. If an access violation simultaneously generates more than one source of abort, they are encoded in the priority given in Table 5-8 on page 5-33. The DMMU fault address register and DMMU fault status register are not updated by faults caused by instruction fetches.

5.10 Fault priority

Table 5-8 lists MMU faults in order of priority, from highest to lowest.

Table 5-8 Priority encoding of MMU faults

Priority	Source	Status	Domain	FAR
Highest	Alignment	0001	Invalid	Valid
	TLB miss	0000	Invalid	Valid
	External abort on level 1 translation	1100	Invalid	Valid
	External abort on level 2 translation	1110	Valid	Valid
	Translation section	0101	Invalid	Valid
	Translation page	0111	Valid	Valid
	Domain section	1001	Valid	Valid
	Domain page	1011	Valid	Valid
	Permission section	1101	Valid	Valid
	Permission page	1111	Valid	Valid
	External abort	1010	Valid	Valid
Lowest	Debug event	0010	Valid	Valid

The values in the domain field are invalid when the fault occurs before the MMU reads the domain field from a page table description. Any abort masked by the priority encoding can be regenerated by fixing the primary abort and restarting the instruction.

5.11 External aborts

The smallest page size the MMU TLB supports is 1KB. This page size is used to filter external aborts.

For an LDM or STM access that does not cross a 1KB page boundary, an external abort is indicated only during the first access of the LDM or STM. For an LDM or STM access that crosses a 1KB page boundary, an external abort can be indicated during the first access of the LDM or STM as well as during the first access that crosses the 1KB page boundary.

In this example, an external abort is possible only on the first access. Table 5-9 shows the sequence.

```
STMIA/LDMIA r0, {r1-r10} r0=0x000000FC
```

Table 5-9 First-access-only external abort

Time	Address	Contents	Comments
t1	0x000000FC	R1	External abort access is possible only on first access.
t2	0x00000100	R2, R3	-
t3	0x00000108	R4, R5	-
t4	0x00000110	R6, R7	-
t5	0x00000118	R8, R9	-
t6	0x00000120	R10	-

In the next example, external aborts are possible on the first access and on page boundary crossings. Table 5-10 shows the sequence.

```
STMIA/LDMIA r0, {r1-r10} r0=0x000003F8
```

Table 5-10 First-access and page-boundary external aborts

Time	Address	Contents	Comments
t1	0x000003F8	R1, R2	External abort is possible on first access
t2	0x00000400	R3, R4	External abort is possible on page boundary crossing
t3	0x00000408	R5, R6	-
t4	0x00000410	R7, R8	-
t5	0x00000418	R9, R10	-

In the next example, external aborts are possible on the first access and on page cross access (last access). Table 5-11 shows the sequence.

STMIA/LDMIA r0, {r1-r10} r0=0x000003E0

Table 5-11 First-access and last-access external aborts

Time	Address	Contents	Comments
t1	0x000003E0	R1, R2	External abort is possible on first access
t2	0x000003E8	R3, R4	-
t3	0x000003F0	R5, R6	-
t4	0x000003F8	R7, R8	-
t5	0x00000400	R9, R10	External abort is possible on page cross access (last access)

In addition to the MMU-generated aborts, the AMBA bus can externally abort the ARM10 processor, which can be used to flag an error on an external memory access. However, not all accesses can be aborted in this way, and the BIU ignores external aborts that cannot be handled.

The following accesses might be aborted:

- a noncached read
- an unbuffered write
- a page table descriptor fetch
- a read-lock-write sequence to noncachable memory.

In the case of a read-lock-write (SWP) sequence in which the read aborts, the write is always canceled.

5.12 Interaction of the MMU, caches, and write buffer

Bit 0 of CP15 c1 enables and disables the MMU.

5.12.1 Enabling the MMU

To enable the MMU:

1. Program the translation table base and domain access control registers.
2. Program level 1 and level 2 descriptor page tables as required.
3. Enable the MMU by setting bit 0 in CP15 c1.

Note

You must take care if the translated address differs from the untranslated address because several instructions following the enabling of the MMU might have been prefetched with the MMU off (using PA = VA *flat translation*), and enabling the MMU might be considered as a branch with delayed execution. A similar situation occurs when the MMU is disabled. Consider the following code sequence:

```
MRC p15, 0, R1, c1, c0, 0 ; Read control register
ORR R1, R1, #0x1
MCR p15, 0, R1, c1, c0, 0 ; Enable MMUs
Fetch Flat
Fetch Flat
Fetch Translated
```

The ICache DCache can be enabled simultaneously with the MMU using a single MCR instruction (see *CP15 c1, Control Register 1* on page 4-7).

5.12.2 Disabling the MMU

To disable the MMU, clear bit 0 in CP15 c1. The data cache must be disabled prior to, or at the same time as the MMU being disabled, by clearing bit 2 for the control register (see *Enabling the MMU* regarding prefetch effects).

Note

If the MMU is enabled, then disabled and subsequently reenabled the contents of the TLBs are preserved. If these are now invalid, you must invalidate the TLBs before the MMU is reenabled (see *CP15 c8, TLB Operations Register* on page 4-19).

5.13 Soft page table support

The soft TLB structure of the MMU requires that TLB entries be written from within the Prefetch Abort handler or Data Abort handler. Because the ARM processor contains both an IMMU and DMMU, there are separate instructions for writing entries into the instruction TLB and the data TLB. The instructions for writing to the instruction TLB are as follows:

```
MCR p15, 0, r2, c15, c8, 1 ; write r2 into I-TLB CAM holding reg
MCR p15, 0, r3, c15, c8, 4 ; write r3 into I-TLB protection RAM holding reg
MCR p15, 0, r4, c15, c8, 6 ; write r4 into I-TLB phys.address RAM holding reg
MCR p15, 0, r1, c15, c0, 3 ; write holding regs into I-TLB at index r1
```

The instructions for writing into the data TLB are as follows:

```
MCR p15, 0, r2, c15, c10, 1 ; write r2 into D-TLB CAM holding reg
MCR p15, 0, r3, c15, c10, 4 ; write r3 into D-TLB protection RAM holding reg
MCR p15, 0, r4, c15, c10, 6 ; write r4 into D-TLB phys. address RAM holding reg
MCR p15, 0, r1, c15, c0, 5 ; write holding regs into D-TLB at index r1
```

Figure 5-14 shows the instruction TLB bit fields.



Figure 5-14 Instruction TLB bit fields

Table 5-12 describes the instruction TLB bit fields.

Table 5-12 Encoding of instruction TLB bit fields

Bits	Name	Meaning
[31:28]	-	Should Be Zero
[27]	V	Valid bit: 1 = valid entry 0 = invalid entry
[26]	L	Lock bit: 1 = locked 0 = not locked
[25:22]	Mask bits	Mask bits [3:0]: 0111 = 64KB page (check CAM tag bits [31:16] against new address) 0011 = 16KB page (check CAM tag bits [31:14] against new address) 0001 = 4KB page (check CAM tag bits [31:12] against new address) 0000 = 1KB page (check CAM tag bits [31:10] against new address)
[21:0]	Tag bits	Tag bits [31:10]

Figure 5-15 shows the protected RAM bit fields.

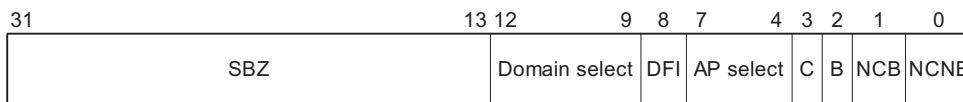


Figure 5-15 Protected RAM bit fields

Table 5-13 describes the protected RAM bit fields.

Table 5-13 Protected RAM bit field values

Bit	Name	Meaning
[31:13]	-	Should Be Zero
[12:9]	Domain select	Domain select bits
[8]	DFI	Domain fault indicator bit: 1 = fault 0 = no fault
[7:4]	AP select	Access index bits [3:0] (2-to-4 encoded): 0000 if not a client of the domain 0001 if client and AP = 00 0010 if client and AP = 01 0100 if client and AP = 10 1000 if client and AP = 11
[3]	C	Cachable bit
[2]	B	Bufferable bit
[1]	NCB	Noncachable bufferable bit
[0]	NCNB	Noncachable nonbufferable bit

Figure 5-16 shows the physical address RAM bit fields.

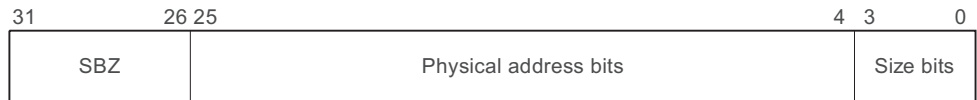


Figure 5-16 Physical address RAM bit fields

Table 5-14 describes the physical address bit fields.

Table 5-14 TLB physical address bit fields and meanings

Bits	Meaning
[31:26]	Should Be Zero
[25:4]	Physical address bits [31:10]: 1111 = 1MB page (address constructed by PA RAM[31:20] + VA[19:0]) 0111 = 64KB page (address constructed by PA RAM[31:16] + VA[15:0]) 0011 = 16KB page (address constructed by PA RAM[31:14] + VA[13:0]) 0001 = 4KB page (address constructed by PA RAM[31:12] + VA[11:0]) 0000 = 1KB page (address constructed by PA RAM[31:10] + VA[9:0])
[3:0]	Size bits

5.13.1 Locked entry requirements

To properly service the IMMU and DMMU aborts when using soft TLB support, the MMUs must have the following entries locked prior to being enabled to guarantee that an infinite abort loop is not entered:

- all exception handler entry points
- any support code for the exception handlers
- any exception handler literal pool accessed areas
- all soft TLB abort handling routines
- any support code for the soft TLB abort handling routines
- any literal pool accessed areas required by soft TLB routines.

5.13.2 Prefetch Abort and Data Abort handling routines

This section gives examples of:

- a Prefetch Abort handler
- a Data Abort handler.

Example 5-1 is a Prefetch Abort handler routine.

Example 5-1 Prefetch Abort handler routine

```
I_softTLB_abort_handler
    < other abort code here >

    MRC p15, 0, r6, c5, c0, 1      ; read instruction FSR
    AND r6, r6, #0xf              ; mask out all but bits 3:0
```

```

CMP r6, #0x0          ; should be 0b0000 if soft TLB abort
BEQ I_softTLB_abort_handler_fix

< other abort code here >

B I_softTLB_abort_handler_fix_end
LTOrg

I_softTLB_abort_handler_fix
MRC p15, 0, r11, c1, c0, 0      ; read CP15 register 1
BIC r6, r11, #0x1              ; disable the MMU
MCR p15, 0, r6, c1, c0, 0      ; reprogram CP15 register 1

MOV r6, r14, lsr #10           ; r14 contains VA
ORR r6, r6, #0x08000000        ; mark valid, map to 1KB page
MCR p15, 0, r6, c15, c8, 1     ; write I-TLB CAM holding register

MOV r6, #0x08c                 ; domain 0, cachable, bufferable
MCR p15, 0, r6, c15, c8, 4     ; write I-TLB protection RAM holding register

MOV r6, r14, lsr #6            ; r14 contains PA
BIC r6, r6, #0x0f              ; map to 1KB page
MCR p15, 0, r6, c15, c8, 6     ; write I-TLB physical address RAM holding register

MRC p15, 0, r4, c10, c0, 1     ; read I-TLB lockdown register
MOV r6, r4, lsl #6             ; shift victim into position
MCR p15, 0, r6, c15, c0, 3     ; write holding regs into I-TLB entry

TST r6, #0x03f00000            ; check for last entry in victim (0 to 63)
MOVEQ r6, r4, lsr #6           ; if last entry, victim=base
ADDNE r6, r6, #0x00100000      ; otherwise increment victim pointer
BIC r4, r4, #0x03f00000        ; clear out old victim
ORR r4, r4, r6                  ; insert new victim
MCR p15, 0, r4, c10, c0, 1     ; write I-TLB lockdown register

MCR p15, 0, r11, c1, c0, 0     ; restore CP15 register 1

I_softTLB_abort_handler_fix_end

< other abort code here >

SUBS pc, r14, #4                ; return to aborted instruction

```

Example 5-2 on page 5-42 is a Data Abort handler routine.

Example 5-2 Data Abort handler routine

D_softTLB_abort_handler

< other abort code here >

```
MRC p15, 0, r6, c5, c0, 0 ; read data FSR
AND r6, r6, #0xf          ; mask out all but bits 3:0
CMP r6, #0x0             ; should be 0b0000 if soft TLB abort
BEQ D_softTLB_abort_handler_fix
```

< other abort code here >

```
B D_softTLB_abort_handler_fix_end
LTCOR
```

D_softTLB_abort_handler_fix

```
MRC p15, 0, r11, c1, c0, 0 ; read CP15 register 1
BIC r6, r11, #0x1         ; disable the MMU
MCR p15, 0, r6, c1, c0, 0 ; reprogram CP15 register 1

MRC p15, 0, r7, c6, c0, 0 ; read the data fault address

MOV r6, r7, lsr #10       ; r7 contains VA
ORR r6, r6, #0x08000000   ; mark valid, map to 1KB page
MCR p15, 0, r6, c15, c8, 1 ; write D-TLB CAM holding register

MOV r6, #0x08c           ; domain 0, cachable, bufferable
MCR p15, 0, r6, c15, c8, 4 ; write D-TLB protection RAM holding register

MOV r6, r7, lsr #6       ; r7 contains PA
BIC r6, r6, #0x0f        ; map to 1KB page
MCR p15, 0, r6, c15, c8, 6 ; write D-TLB physical address RAM holding register

MRC p15, 0, r4, c10, c0, 1 ; read D-TLB lockdown register
MOV r6, r4, lsl #6       ; shift victim into position
MCR p15, 0, r6, c15, c0, 3 ; write holding regs into D-TLB entry

TST r6, #0x03f00000      ; check for last entry in victim (0 to 63)
MOVEQ r6, r4, lsr #6     ; if last entry, victim=base
ADDNE r6, r6, #0x00100000 ; otherwise increment victim pointer
BIC r4, r4, #0x03f00000  ; clear out old victim
ORR r4, r4, r6           ; insert new victim
MCR p15, 0, r4, c10, c0, 1 ; write D-TLB lockdown register

MCR p15, 0, r11, c1, c0, 0 ; restore CP15 register 1
```

D_softTLB_abort_handler_fix_end

< other abort code here >

SUBS pc, r14, #8 ; return to aborted instruction

Chapter 6

Caches and Write Buffer

This chapter describes the *Instruction Cache* (ICache), the *Data Cache* (DCache), and the write buffer. It contains the following sections:

- *About the caches and write buffer* on page 6-2
- *ICache* on page 6-3
- *DCache and write buffer* on page 6-7
- *Cache coherence* on page 6-15
- *Portability issues* on page 6-17.

6.1 About the caches and write buffer

The ARM processor includes:

- an ICache
- a DCache
- a write buffer
- a *Hit-Under-Miss* (HUM) buffer.

The 32KB ICache and 32KB DCache have the following features:

- Sixteen segments, each containing 64 lines.
- Virtually-addressed 64-way associativity.
- Eight words per line (32 bytes per line) with one valid bit, one dirty bit, and one write-back bit per line.
- Write-through and write-back (copy-back) DCache operation, selected per memory region by the C and B bits in the MMU translation tables.
- Pseudorandom or round-robin replacement, selectable by the RR bit in CP15 c1.
- Low-power CAM-RAM implementation.
- Independently lockable caches with granularity of $1/64^{\text{th}}$ of the cache, that is 128 words (512 bytes) to a maximum of $63/64^{\text{th}}$ of the cache.
- For compatibility with Microsoft WindowsCE, and to reduce interrupt latency, the physical address corresponding to each DCache entry is stored in the DCache PA tag RAM for use during cache line write-backs, in addition to the VA tag stored in the cache CAMs. This means that the MMU is not involved in cache write-back operations, removing the possibility of MMU misses related to the write-back address.
- Cache maintenance operations to provide efficient cleaning of the entire DCache, and to provide efficient cleaning and invalidation of small regions of virtual memory. The latter enables ICache coherency to be efficiently maintained when small code changes occur, for example, self-modifying code and changes to exception vectors.

The write buffer can hold eight 64-bit packets of data, each with an associated address element.

6.2 ICache

The 32KB ICache has 1024 lines of 32 bytes. It is arranged as a 64-way set-associative cache and uses virtual addresses from the integer unit.

The ICache uses *allocate-on-read-miss* linefills. The RR bit in CP15 c1 selects random or round-robin replacement. After reset, replacement is random.

You can also lock instructions in the ICache so that they cannot be overwritten by a linefill. Lockdown operates with a granularity of $1/64^{\text{th}}$ of the cache, which is 128 words (512 bytes), to a maximum of $63/64^{\text{ths}}$ of the cache.

All instruction accesses are subject to MMU permission and translation checks. Instruction fetches that are aborted by the MMU do not cause linefills or instruction fetches to appear on the AHB.

The following sections describe the ICache:

- *ICache enable/disable*
- *ICache operation*
- *ICache cachable control* on page 6-5
- *ICache replacement algorithm* on page 6-5
- *ICache lockdown* on page 6-5.

6.2.1 ICache enable/disable

Reset invalidates all ICache entries and disables the ICache. Setting the I bit in CP15 c1 enables the ICache. Clearing I disables it.

When the ICache and the MMU are enabled, the C bit in the relevant MMU translation table descriptor indicates whether an area of memory is *cacheable* (C). If the ICache is enabled and the MMU disabled, all instruction fetches are treated as cacheable.

When the ICache is disabled, the cache contents are ignored and all instruction fetches appear on AHB as separate nonsequential accesses. Reenabling the ICache does not change its contents. If the contents are no longer coherent with main memory, you must invalidate the ICache before enabling it (see *CP15 c7, Index and VA Cache Operations Registers* on page 4-16).

You can enable the ICache and MMU simultaneously by setting bits I and M in CP15 c1 with a single MCR instruction.

6.2.2 ICache operation

Enable the ICache as soon as possible after reset.

When the ICache is disabled, each instruction fetch results in a separate nonsequential memory access on AHB, giving very low performance to burst memory such as page mode DRAM or synchronous DRAM. When the ICache is enabled, an ICache lookup is performed for each instruction fetch regardless of the setting of the C bit in the relevant MMU translation table descriptor. If the required instruction is found in the cache, the lookup result is called a *cache hit*. If the required instruction is not found in the cache, the lookup result is called a *cache miss*.

If the instruction fetch is a cache hit and is being fetched from a cachable region of memory, then the instruction is returned from the cache to the integer unit. If the instruction fetch is a cache miss, then an 8-word cache linefill is performed, possibly replacing another entry. The entry to be replaced, the *victim*, is chosen by either random or round-robin replacement from the entries that are not locked.

If an instruction fetch is from a *noncachable* (NC) region of memory, then a single nonsequential memory access appears on the AHB. This access to the AHB is independent of the ICache being enabled.

———— **Note** ————

If a program is fetching from a noncachable region of memory, then the cache lookup results in a cache miss. The only way that it can result in a cache hit is if software has changed the value of the cachable bit in the MMU translation table descriptor without invalidating the cache contents. This is a programming error and the behavior in this case is architecturally unpredictable.

6.2.3 ICache cachable control

In the MMU translation table descriptors, the C bit defines cachable regions of memory. In CP15 c1, control register 1, the I bit enables the ICache, and the M bit enables the IMMU. Table 6-1 shows how to select cachable instructions.

Table 6-1 Selection of cachable instructions

CP15 c1 M bit	CP15 c1 I bit	MMU C bit	In debug	Memory region type
0	0	-	0	NC flat mapped
0	1	-	0	C
1	0	-	0	NC
1	1	0	0	NC
1	1	1	0	C
-	-	-	1	NC

The following sections describe the ICache behavior when accessing cachable and noncachable memory.

Cachable (C)

Reads that hit in the cache read instructions from the cache. Reads that miss in the cache cause a linefill and cannot be externally aborted. The linefill performs an AHB access.

Noncachable (NC)

Reads not cached always perform an AHB access and can be externally aborted. Cache hits never occur.

6.2.4 ICache replacement algorithm

The RR bit in CP15 c1 selects the ICache and DCache replacement algorithm. Reset selects random replacement. Setting the RR bit selects round-robin replacement.

6.2.5 ICache lockdown

Instructions can be locked into the ICache, guaranteeing an ICache hit and providing optimum and predictable execution time.

Lock instructions into the ICache by first ensuring that the code to be locked is not already in the cache. Do this by flushing either the whole ICache or specific lines. You can then use a short software routine to load the instructions into the ICache. The software routine can either be noncachable or already in the ICache, but not in an ICache line that is about to be overwritten. The instructions to be loaded must be from a memory region that is cachable.

You can perform the prefetch ICache line by writing to CP15 c9 to force the replacement counter to a specific ICache line. Then issue a prefetch ICache line operation using CP15 c7. If the prefetch is to a cachable region and misses in the ICache, the prefetch is performed. When the prefetch is complete, the replacement counter increments the pointer to the next ICache line. This operation can be repeated for multiple prefetchable ICache lines.

When all the instructions are loaded, lock them by writing to CP15 c9 to set the replacement counter base to be one higher than the number of locked cache lines.

See *DCache lockdown* on page 6-12 for a more complete explanation of cache locking.

6.3 DCache and write buffer

The 32KB DCache has 1024 lines of 32 bytes, arranged as a 64-way set-associative cache. It uses virtual addresses from the integer unit. The write buffer can hold up to eight 64-bit packets of data and four additional 64-bit packets of data in a separate castout buffer. Each data packet has an associated address packet. The write buffer can hold eight double words of regular buffered writes and an entire cache line (four double words) in a separate castout buffer. The operation of the DCache and write buffer are closely connected.

The DCache supports WT and WB memory regions, controlled by the C and B bits in each section and page descriptor within the MMU translation tables. For details see *DCache and write buffer operation* on page 6-8.

Each DCache line has:

- one valid bit, one dirty bit, and one write-back bit
- a single virtual tag address
- sixteen 32-bit data elements (eight-word line)
- a single physical address tag, used when writing modified lines back to memory.

A linefill always loads a complete eight-word line starting with the critical 64-bit data.

When a store instruction hits in the DCache, the associated dirty bit is set marking the appropriate line as modified. If the cache line is replaced due to a linefill, or if the line is the target of a DCache clean operation, the dirty bit and write-back bits are used to decide whether the line is written back to memory. The line is written back to the same physical address from which it was loaded, regardless of any changes to the MMU translation tables.

The DCache uses *allocate-on-read-miss* linefills. The RR bit in CP15 c1 selects random or round-robin replacement. Reset selects random replacement.

You can also lock data in the DCache so that it cannot be overwritten by a linefill. Lockdown operates with a granularity of $1/64^{\text{th}}$ of the cache, which is 128 words (512 bytes), with the maximum lockdown value being $63/64^{\text{th}}$ of the cache.

All data accesses are subject to MMU permission and translation checks. Data accesses that are aborted by the MMU do not cause linefills or data accesses to appear on the AHB.

The following sections describe the DCache and write buffer:

- *DCache and write buffer enable/disable* on page 6-8
- *DCache and write buffer operation* on page 6-8
- *DCache cachable and bufferable control* on page 6-9
- *DCache replacement algorithm* on page 6-11

- *Swap instructions* on page 6-11
- *DCache organization* on page 6-12
- *DCache lockdown* on page 6-12
- *Hit-Under-Miss* on page 6-13.

6.3.1 DCache and write buffer enable/disable

Reset invalidates all DCache entries, disables the DCache, and discards the contents of the write buffer.

The W bit in CP15 c1 can enable and disable the write buffer during program execution. Disabling the write buffer forces all stores (writes) to a region type of *NonCachable NonBufferable* (NCNB) regardless of the TLB region definition.

Enable the DCache by setting the C bit in CP15 c1.

The DCache must be enabled only when the MMU is enabled. This is because the MMU translation tables define the cache and write buffer configuration for each memory region.

When the DCache is disabled, the cache contents are ignored and all data accesses appear on the AHB as separate nonsequential accesses. If the cache is subsequently reenabled its contents are unchanged. Depending on the software system design, the cache might have to be cleaned after it is disabled, and invalidated before it is reenabled (see *Cache coherence* on page 6-15.)

The MMU and DCache can be enabled or disabled simultaneously with a single MCR that changes the M and C bits in CP15 c1.

6.3.2 DCache and write buffer operation

The DCache and write buffer configuration of each memory region is controlled by the C and B bits in each section and page descriptor in the MMU translation tables.

If the DCache is enabled, a DCache lookup is performed for each data access initiated by the ARM processor, regardless of the value of the C bit in the relevant MMU translation table descriptor. If the accessed virtual address matches the virtual address of an entry in the cache, the lookup result is a cache hit. If the required address does not match any entry in the cache, the lookup result is a cache miss. In this context a data access means any type of load (read), store (write), swap, or cache preload instruction.

To ensure that accesses appear on the AHB in program order, the ARM processor waits for all writes in the write buffer to complete on the AHB before starting any other AHB access. The integer unit can continue executing at full speed reading instructions and data from the caches, and writing to the DCache and write buffer while buffered writes are being written to memory over the AHB.

6.3.3 DCache cachable and bufferable control

A linefill loads eight words, starting with the critical 64-bit data word, by performing a four-beat wrapping read burst on the AHB.

A load multiple (LDM) instruction accessing NCNB or NCB regions performs a series of nonsequential read transfers on the AHB. A store multiple (STM) instruction accessing NCNB regions also performs the writes as a series of nonsequential transfers.

In the MMU translation table descriptors, the I bit and the C bit define cachable and bufferable regions of memory. In CP15 c1, control register 1, the C bit enables the DCache, and the M bit enables the DMMU. Table 6-2 shows how to select cachable and bufferable data.

Table 6-2 Selection of cachable and bufferable data

CP15 r1 M bit	CP15 r1 C bit	MMU C bit	MMU B bit	Memory region type
0	-	-	-	NCNB flat-mapped
1	0	-	-	NCNB
1	1	0	0	NCNB
1	1	0	1	NCB
1	1	1	0	WT
1	1	1	1	WB

The following sections describe the DCache and write buffer behavior for each memory region type.

NonCachable, NonBufferable (NCNB)

- Swaps are atomic operations that lock the AHB for both the read and write.
- Reads and writes are not cached, use the AHB, and can be externally aborted.
- Writes are not buffered.
- The LSU halts on reads and writes until the operation completes on the AHB.

- Cache hits should never occur.

NonCachable, Bufferable (NCB)

- Swaps to a NCB region behave like a swap to an NCNB region
- Reads are not cached, use the AHB, and can be externally aborted.
- Cache hits should never occur.
- Writes are placed in the write buffer and cannot be externally aborted.

Cachable, Write-Through (WT)

- Reads that hit in the cache read the data from the cache.
- Reads that miss in the cache cause a linefill.
- All writes are placed in the write buffer.
- Writes that hit in the cache update the cache.
- Reads and writes cannot be externally aborted.

Cachable, Write-Back (WB)

- Reads that hit in the cache read the data from the cache.
- Reads that miss in the cache cause a linefill.
- Writes that miss in the cache are placed in the write buffer.
- Writes that hit in the cache update the cache and mark the cache line as dirty.
- Cache write-backs and castouts are buffered.
- Reads and writes (write-misses and write-backs) cannot be externally aborted.

It is an operating system software error if a cache hit occurs when reading or writing a region of memory marked as NCNB or NCB. This can occur only if the operating system changes the value of the C and B bits in a page table descriptor while the cache contains data from the area of virtual memory controlled by that descriptor. The cache and memory system behavior resulting from changing the page table descriptor in this way is Unpredictable. If the operating system has to change the C and B bits of a page table descriptor, it must ensure that the caches do not contain any data controlled by that descriptor. In some circumstances, the operating system might have to clean and flush the caches to ensure this.

A read that triggers a linefill performs an eight-word burst read from the AHB and places it as a new entry in the cache, possibly replacing another line at the same location within the cache. The location that is replaced, the victim, is chosen from the nonlocked entries using either random or round-robin replacement. If the cache line being replaced is marked as dirty, indicating that it has been modified and that main memory has not

been updated to reflect the change, a cache write-back occurs. The write-back data is placed in the castout buffer at the same time that linefill data is placed in the victim line. The CPU can then continue immediately following the request issued to the DCache.

Load multiple (LDM) instructions accessing NCNB or NCB regions perform sequential bursts on the AHB. Store multiple (STM) instructions accessing NCNB regions also perform sequential bursts on the AHB.

A cache preload (PLD) instruction behaves like a load (read) single. If the region type is WT or WB, the cache preload performs a linefill. If the region type is NCNB or NCB, the cache preload stalls the memory system for one cycle and does not request anything on AHB.

The sequential burst is split into two bursts if it crosses a 1KB boundary. This is because the smallest MMU protection and mapping size is 1KB, so the memory regions on each side of the 1KB boundary might have different properties.

This means that no sequential access generated by the ARM processor crosses a 1KB boundary. You can exploit this feature to simplify memory interface design. For example, a simple page-mode DRAM controller can perform a page-mode access for each sequential access, provided the DRAM page size is 1KB or larger (see also *Cache coherence* on page 6-15).

6.3.4 DCache replacement algorithm

The DCache replacement algorithm is selected by the RR bit in CP15 c1. Random replacement is selected at reset. Setting the RR bit selects round-robin replacement.

6.3.5 Swap instructions

Swap instruction (SWP or SWPB) behavior is dependent on whether the memory region is cachable or noncachable:

- Swap instructions to cachable regions of memory are useful for implementing semaphores or other synchronization primitives in multithreaded uniprocessor software systems.
- Swap instructions to noncachable memory regions are useful for synchronization between two bus masters in a multimaster bus system. This can be two processors, or a processor and a DMA controller.

When a swap instruction accesses a cachable region of memory (WT or WB), the DCache and write buffer behavior is the same as having a load followed by a store. The AHB does not assert the HLOCK pin to swap instructions that access a cachable region and hit in the DCache. It is guaranteed that no interrupt can occur between the load and store portions of the swap.

When a swap instruction accesses a noncachable (NCB or NCNB) region of memory, the write buffer is emptied, and a single word or byte is read from the AHB. The write portion of the swap is then treated as nonbufferable, with the LSU stalled until the write is completed on the AHB. The HLOCKD pin is asserted to indicate that the read and write must be treated as an atomic operation on the bus.

Like all other data accesses, a swap to a noncachable region that hits in the cache indicates a programming error.

6.3.6 DCache organization

The DCache is organized as sixteen segments, each containing 64 lines, and each line containing eight words. The position of the line within its segment is a number from 0 to 63 which is called the index. A line in the cache can be uniquely identified by its segment and index. The index is independent of the virtual address of the line. The segment is selected by bits [8:5] of the virtual address of the line.

Bits [4:3] of the virtual address specify which 64-bit word within a cache line is accessed. Bit 2 specifies which 32-bit word in a 64-bit word is accessed. For halfword operations, bit 1 of the virtual address specifies which halfword is accessed within the word. For byte operations, bits [1:0] specify which byte within the word is accessed.

Bits [31:9] of the virtual address of the each cache line is called the tag. The virtual address tag is stored in the cache, with the eight words of data, when the line is loaded by a linefill.

Cache lookups compare bits [31:9] of the virtual address of the access with the stored tag to determine whether the access is a hit or miss. The cache is therefore said to be virtually addressed.

6.3.7 DCache lockdown

Data can be locked into the DCache causing the DCache to guarantee a hit, and providing optimum and predictable execution time.

When no data is locked in the DCache, and a linefill occurs, the replacement algorithm chooses a victim cache line to be replaced by selecting an index in the range 0 to 63. The segment is specified by bits [8:5] of the virtual address of the data access that missed.

Data is locked into the DCache by restricting the range of victim numbers produced by the replacement algorithm, so that locked down cache lines are never selected as victims. You can set the base pointer for the DCache victim generator by writing to CP15 c9. The replacement algorithm chooses a victim cache line in the range (base to 63), locking in the cache the lines with index in the range (0 to base - 1).

Data is loaded and locked into the DCache by first ensuring that the data to be locked is not already in the cache. This can be ensured by cleaning and flushing either the whole DCache or specific lines. A short software routine can then be used to load the data into the DCache.

The software routine to load the data operates by writing to CP15 c9 to force the replacement counter to a specific DCache line and then executing a load instruction to perform a cache lookup. This misses and a linefill is performed, bringing eight words of data into the cache line specified by the replacement counter, in the segment specified by bits [8:5] of the virtual address accessed by the load.

To load further lines into the cache, the software routine can loop performing one load from each line to be loaded. As each line contains eight words, each loop adds 32 (bytes) to the load address. When a linefill is acknowledged in a particular segment, the replacement increments to point to the next DCache line.

When all the data has been loaded, it is locked by writing to CP15 c9 to move the replacement counter base to be one higher than the highest index of the locked cache lines.

The software routine that loads and locks the data in the DCache can be located in a cachable region of memory providing it does not contain any loads or stores other than the loads that are used to bring the data to be locked into the DCache. The data to be loaded must be from a memory region that is cachable.

6.3.8 Hit-Under-Miss

The ARM processor supports HUM operation. Clearing the fast interrupt bit, FI, in CP15 c1 with a read-modify-write operation enables HUM operation. Reset clears FI, enabling HUM by default. Software can change the state of FI dynamically. Any pending load or store in the LSU pipeline completes before the CP15 c1 operation takes effect.

When the FI bit is set, all load misses in the data cache stall the LSU pipeline until completion of the linefill. This prevents multiple linefill requests from accumulating and so reduces the amount of activity that must complete prior to servicing an interrupt request. Setting FI also reduces the write buffer from eight entries to four entries.

When the FI bit is cleared, HUM activity occurs as described in this section. Briefly, setting FI enables load/store instructions to execute while a linefill is being serviced. If a load request misses in the DCache while a linefill for a prior request is in progress, the LSU pipeline halts. This is referred to as a *second load miss*.

There are two scenarios that can arise from the second load miss:

- The second load miss is to the cache line that is currently being filled. The data cache returns the second load miss data during an idle ARM load/store cycle following the return of the critical word from the first load miss while the linefill is still being performed. The ability to return data for the second load miss before the linefill completes is referred to as *data streaming* or *load streaming*. When data for the second load miss is returned, the LSU pipeline activity can resume.
- The second load miss is to a different cache line than the line currently being filled. The data cache completes the linefill for the first load miss before triggering any activity for handling the second miss. On completion of the first linefill, the data cache then triggers a linefill for the second load miss, promoting the second load miss to the first load miss position. Then the LSU pipeline can resume execution.

If an NCNB load, NB store, or data MMU page table walk occurs at any time during a linefill, the LSU pipeline stalls until the linefill completes. On completion of the first linefill, the transfer request that caused an NCNB load, NB store, or data MMU page table walk is enabled to advance.

During a linefill, a store can also be executed. There are also two scenarios that can arise for a store:

- The store hits in the cache line being filled. If this occurs, the store is merged or folded into the cache linefill so that coherency is maintained. That is, the data cache always has the latest copy of data. If the store region is write-through, and the write buffer is enabled, the store is placed in the write buffer. If the store region is write-back, the store is not placed in the write buffer because it has been merged with the line being filled.
- The store does not hit in the cache line being filled. If this occurs, the store is simply placed in the write buffer if the region type is write-through. If the region type is write-back and the store hits, the store updates the data cache. If the store misses, the store is placed in the write buffer.

6.4 Cache coherence

The ICache and DCache contain copies of information usually held in main memory. If these copies of memory information get out of step with each other because one is updated and the other is not updated, they are said to have become incoherent. If the DCache contains a line that has been modified by a store or swap instruction, and the main memory has not been updated, the cache line is said to be *dirty*. Clean operations force the cache to write dirty lines back to main memory.

Software is responsible for maintaining coherency between main memory, the ICache, and the DCache.

CP15 c7, Index and VA Cache Operations Registers on page 4-16 describes facilities for invalidating the entire ICache or DCache or individual ICache or DCache lines, and for cleaning the entire DCache or individual DCache lines.

To clean the entire DCache efficiently, software must loop through each cache entry using the *clean D single entry (using index)* operation or the *clean and invalidate D entry (using index)* operation. This must be performed by a two-level nested loop going through each index value for each segment (see *DCache organization* on page 6-12).

DCache, ICache, and memory coherence is generally achieved by:

1. cleaning the DCache to ensure memory is up-to-date with all changes
2. invalidating the ICache to force it to reload instructions from memory.

Software can minimize performance penalties of cleaning and invalidating caches by:

- cleaning only small portions of the DCache when only a small area of memory must be made coherent, for example, when updating an exception vector entry
- invalidating only small portions of the ICache when only a small number of instructions are modified, for example, when updating an exception vector entry
- not invalidating the ICache in situations where it is known that the modified area of memory cannot be in the cache, for example, when mapping a new page into the currently running process.

The ICache needs to be made coherent with a changed area of memory after any changes to the instructions that appear at a virtual address, and before the new instructions are executed.

Dirty data in the DCache can be pushed out to main memory by cleaning the DCache.

Cache cleaning and invalidating are necessary when:

- writing instructions to a cachable area of memory using STR or STM instructions:
 - self-modifying code

- JIT compilation
- copying code from another location
- downloading code using the EmbeddedICE JTAG debug features
- updating an exception vector entry.
- another bus master modifies a cachable area of main memory
- turning the MMU on or off
- changing the virtual-to-physical mapping in the MMU page tables
- turning the ICache or DCache on, if its contents are no longer coherent.

The DCache must be cleaned, and both caches invalidated, before the cache and write buffer configuration of an area of memory is changed by modifying the C bit or B bit in the MMU translation table descriptor. This is not necessary if the caches cannot contain any entries from the area of memory whose translation table descriptor is being modified.

Changing the process ID in CP15 c13 does not change the contents of the cache or memory, and does not affect the mapping between cache entries and physical memory locations. It only changes the mapping between addresses and cache entries. This means that changing the process ID does not lead to any coherency issues. Changing the process ID does not require cache cleaning or cache invalidation.

Reset invalidates and disables the DCache and ICache.

The pipelined design of the integer unit means that it fetches three instructions ahead of the current execution point. So, for example, before an MCR that invalidates the ICache executes, the ARM processor reads from three to five instructions following the MCR.

6.4.1 Cache cleaning when lockdown is in use

The *clean D single entry (using index)* operation only modifies the victim for that operation, not the victim pointer. The victim is set back to its previous value on the next cycle. *Clean D single entry (using VA)* and *clean and invalidate D entry (using VA)* operations do not move the victim pointer, so there is no need to reposition the victim pointer after using these operations.

6.5 Portability issues

This section describes the behavior of the ARM processor in this implementation that is architecturally Unpredictable. For portability to other ARM implementations, software must not depend on this behavior.

A read from a noncachable (NCB or NCNB) region that unexpectedly hits in the cache still reads the required data from the AHB. The contents of the cache are ignored and unchanged. This includes the read portion of a swap (SWP or SWPB) instruction.

A write to a noncachable (NCB or NCNB) region that unexpectedly hits in the cache, updates the cache and still causes an access on the AHB.

Chapter 7

Prefetch Unit

This chapter describes how the prefetch unit fetches instructions to feed to the integer core (and to coprocessors), as well as how it locates branches in the instruction stream for predicting potential changes in sequential instruction issue. It also describes the SWI functions useful for flushing the prefetch buffer. It contains the following sections:

- *About the prefetch unit* on page 7-2
- *Branch prediction activity* on page 7-3
- *Branch instruction cycle summary* on page 7-6
- *Instruction memory barriers* on page 7-8.

7.1 About the prefetch unit

The prefetch unit is responsible for fetching instructions from the memory system as required by the integer core (and coprocessors). The prefetch unit fetches instructions at up to twice the rate that the core requires them, and the prefetch buffer holds up to three instructions. The prefetch buffer enables the prefetch unit to:

- detect branch instructions ahead of the Fetch stage
- predict those branches that are likely to be taken
- remove those branches that are not likely to be taken.

The bus from the memory system to the prefetch unit is 64 bits wide. It can supply two instruction words from a doubleword-aligned address every clock cycle.

Branch prediction enables the prefetch unit to provide the branch target instruction to the Execute stage earlier than if no prediction mechanism is used. Branch prediction increases processor performance by minimizing the cycle time of branch instructions. When the prefetch unit predicts a branch as taken, it calculates the target address and fetches instructions from the new address. Depending on how full the prefetch buffer is at the time the prediction is made, the predicted branch can be reduced to two, one, or zero cycles. When the prefetch unit predicts a branch as not taken, it removes the branch from the instruction stream passed to the core. It still calculates the target address of the branch in case the prediction is incorrect. The prediction mechanism is static. It uses no history information. Conditional forward branches are predicted as not taken and conditional backward branches are predicted as taken.

The prefetch unit performs branch prediction only when the Z bit in CP15 c1 is set.

7.2 Branch prediction activity

The prefetch unit does not predict all branches. It can predict only a branch that is relative to the PC, not a branch with an absolute target address. The integer unit executes branch instructions that are left in the instruction stream passed to the core. The branch prediction logic only optimizes one branch at a time.

When the prefetch unit predicts a branch as taken, it speculatively prefetches from the target address. In speculative prefetching, all cache hits result in an instruction fetched into the prefetch buffer. Cache misses and noncachable accesses in speculative prefetching do not initiate a linefill from memory until the core has resolved the flags and the prediction is confirmed.

7.2.1 Branch folding

Depending on how many instructions are in the prefetch buffer at the time a branch is predicted, the branch may be completely removed from the instruction stream. This means:

- A branch is pulled from the instruction stream based on a prediction.
- The predicted next instruction is substituted in the place of this branch.
- No empty instruction issue slots results in the process.

Under these circumstances the branch itself takes zero cycles because it is removed altogether from the instruction stream to the core. This type of branch removal involving the direct substitution of another instruction is called *branch folding*. The condition codes of the predicted branch are folded onto the predicted next instruction, and only a single instruction is issued to the core. The condition codes of the predicted branch are called the *branch phantom*. The substituted instruction is the folded instruction.

7.2.2 Flushing the prefetch buffer

The prefetch buffer is flushed in all the following cases:

- entry into an exception processing sequence
- a load to PC
- an arithmetic manipulation of the PC
- execution of an unpredicted branch
- detection of an erroneously predicted branch.

The only change to sequential instruction fetching that does not automatically flush the prefetch buffer is the case of a predicted taken branch.

7.2.3 Branch penalty

Mispredicted branches and unpredicted taken branches have a three-cycle penalty (assuming instruction cache hit). Here *penalty* means the number of cycles in which no useful Execute stage pipeline activity can occur due to an instruction flow differing from that assumed or predicted. Table 7-1 illustrates this penalty for the case of an erroneously predicted branch. Cycles 2, 3, and 4 have nothing valid in Execute stage. The activity is similar for an unpredicted branch that is taken. Unpredicted branches that are not taken just consume their normal Execute stage and have no branch penalty.

Table 7-1 Penalty for an erroneously predicted branch

Cycle	Pipe stage	Activity
1	Execute	Branch phantom, probably with a folded instruction. Condition code evaluation results in mispredict. All instructions in earlier pipeline stages are canceled. Folded instructions are canceled.
2	Fetch	Instruction fetch from saved opposite instruction stream.
3	Issue	Correct instruction in Issue stage.
4	Decode	Correct instruction in Decode stage.
5	Execute	Correct instruction in Execute stage.

7.2.4 Optimization of branch instructions

This is a complete list of the branch optimizations performed by the branch prediction unit:

- ARM and Thumb conditional branches are predicted taken and potentially reduced to zero cycles if they branch backwards.
- ARM and Thumb conditional branches are predicted not taken and potentially reduced to zero cycles if they branch forward.
- ARM and Thumb unconditional branches are predicted taken and potentially reduced to zero cycles.
- ARM unconditional BL and BLX instructions are predicted taken and potentially reduced to one cycle.
- A Thumb BL pair (always unconditional) is predicted taken and potentially reduced to one cycle. The pair of instructions must be consecutive in memory for them to be predicted.

- A Thumb BLX pair (always unconditional) is predicted taken and potentially reduced to one cycle. The pair of instructions must be consecutive in memory for them to be predicted.

When BLs and BLXs are predicted, the instruction is changed into a link instruction and a branch instruction. The link part of the instruction is passed to the integer unit as a special MOV LR instruction. The branch part is predicted taken.

Branches are not predicted in any of the following cases:

- the Z bit in CP15 c1 is clear
- a Prefetch Abort occurred when fetching the instruction
- a breakpoint is set on the instruction address.

7.3 Branch instruction cycle summary

The number of cycles taken by the ARM10 processor to execute branch instructions depends primarily on:

- Whether or not the branch is predicted.
- Whether or not the predicted branch is correct.
- What direction the predicted branch takes, forward or backward.
- The number of instructions in the prefetch buffer ahead of the branch at the time the prediction is made. The prefetch buffer continues to issue instructions while a predicted branch target instruction is being fetched.

Table 7-2 shows the instruction cycle counts for all ARM and Thumb branches. The cycle counts are based on ICache hits, because the cycle counts of ICache misses and noncachable accesses vary widely as a function of system and implementation characteristics.

Instructions are listed here by their *ARM Architecture Reference Manual* name. Some instructions have multiple variations that distinguish unique characteristics among a common instruction, for example Thumb B(1) and Thumb B(2).

Table 7-2 ARM and Thumb branch instruction cycle counts

Instruction	Unpredicted condition code		Predicted correctly		Predicted incorrectly	
	Fail	Pass	Backward/taken	Forward/not taken	Backward/taken	Forward/not taken
ARM instructions						
B uncond	a	4	0-2	0 ^b	c	c
B cond	1	4	0-2	0 ^b	4	4
BL uncond	a	4	1-2 ^d	1-2 ^{d, e}	c	c
BL cond	2	4	e	e	e	e
BLX(1) uncond	a	4	1-2 ^d	1-2 ^{d, e}	c	c
BLX(2) uncond	a	4	f	f	f	f
BLX(2) cond	2	4	f	f	f	f
BX uncond	a	4	f	f	f	f

Table 7-2 ARM and Thumb branch instruction cycle counts (continued)

Instruction	Unpredicted condition code		Predicted correctly		Predicted incorrectly	
	Fail	Pass	Backward/taken	Forward/not taken	Backward/taken	Forward/not taken
BX cond	2	4	f	f	f	f
Thumb instructions						
B(1) cond	1	4	0-2	0 ^b	4	4
B(2) uncond	a	4	0-2	0 ^b	c	c
BL uncond	a	5 ^g	1-2 ^d	1-2 ^{d, e}	c	c
BLX(1) uncond	a	5 ^g	1-2 ^d	1-2 ^{d, e}	c	c
BLX(2) uncond	a	4	f	f	f	f
BX uncond	a	4	f	f	f	f

- a. Unconditional branches (either unconditional by instruction definition or by using cond code AL, always, cannot fail condition codes).
- b. All forward branches are only predicted when prefetch buffer contains at least 2 instructions (the branch being predicted and its preceding instruction).
- c. Unconditional branches, when predicted, can never be erroneously predicted.
- d. BL and BLX (ARM and Thumb) can never be reduced to 0 cycles by prediction because the link operation necessarily consumes a cycle.
- e. BL and BLX (ARM and Thumb) are only predicted if unconditional, in which case they are predicted taken irrespective of direction (guaranteed to be correct).
- f. BX and BLX(2) instructions, ARM and Thumb, are not pc-relative. They cannot be predicted.
- g. Thumb BL and BLX(1) instructions are encoded as two Thumb instructions. The first of these is a data processing instruction that puts an immediate into R14 then fetches from that address. This second instruction takes 4 cycles (before the next instruction is in Execute).

7.4 Instruction memory barriers

The prefetch unit performs speculative prefetching of instructions. In some circumstances it is likely that the prefetch buffer contains out-of-date instructions. In these circumstances the prefetch buffer must be flushed. An *Instruction Memory Barrier* (IMB) sequence provides a means to do this.

You can include processor-specific code in the SWI handler to implement the two IMB sequences:

IMB The IMB sequence flushes all information about all instructions.

IMBRange When only a small area of code is altered before being executed, the IMBRange sequence can efficiently and quickly flush any stored instruction information from addresses within a small range. By flushing only the required address range information, the rest of the information remains to provide improved system performance.

The IMB and IMBRange sequences are implemented as calls to specific SWI numbers.

7.4.1 Generic IMB use

Use SWI functions to provide a well-defined interface between code that is:

- independent of the ARM processor implementation on which it is running
- specific to the ARM processor implementation on which it is running.

The implementation-independent code is provided with a function that is available on all processor implementations through the SWI interface, and that can be accessed by privileged and, where appropriate, nonprivileged (User mode) code.

Using SWIs to implement the IMB instructions means that code that is written now remains compatible with future ARM processors, even if those processors implement IMB in different ways. This is achieved by changing the operating system SWI service routines for each of the IMB SWI numbers that differ from processor to processor.

7.4.2 IMB implementation

Executing the SWI instruction is sufficient to cause IMB operation. Also, both the IMB and the IMBRange sequences flush all stored information about the instruction stream.

This means that all IMB instructions can be implemented in the operating system by returning from the IMB/IMBRange service routine and that the service routines can be exactly the same. The following service routine code can be used:

```
IMB_SWI_handler
IMBRange_SWI_handler
```

```
    MOV PC, R14_svc    ; Return to the code after the SWI call
```

Note

In new code, you are strongly encouraged to use the IMBRange sequence whenever the changed area of code is small, even if there is no distinction between it and the IMB sequence. Future ARM processors might implement a faster and more efficient IMBRange sequence, and code migrated from this ARM processor can benefit when executed on future ARM processors.

7.4.3 Execution of IMB sequences

This section gives examples that show what should happen during IMB sequences. The pseudocode in the square brackets shows what should happen in the SWI routine.

Loading code from disk

Code that loads a program from a disk and then branches to the entry point of that program must use an IMB sequence after loading the program and before executing it:

```
IMB EQU 0xF00000
.
.
; code that loads program from disk
.
.
SWI IMB
    [branch to IMB service routine]
    [perform processor-specific operations to execute IMB]
    [return to code]
.
MOV PC, entry_point_of_loaded_program
.
.
```

Running BitBlt code

Compiled BitBlit routines optimize large copy operations by constructing and executing a copying loop that has been optimized for a particular operation. When writing such a routine, an IMB is required between the code that constructs the loop and the execution of the constructed loop:

```
IMBRange EQU 0xF00001
```

```

      .
      .
; code that constructs loop code
; load R0 with the start address of the constructed loop
; load R1 with the end address of the constructed loop
SWI IMBRange
    [branch to IMBRange service routine]
    [read registers R0 and R1 to set up address range parameters]
    [do processor-specific operations to execute IMBRange within address range]
    [return to code]
; start of loop code
      .
      .

```

Self-decompressing code

When writing a self-decompressing program, an IMB should be issued after the routine that decompresses the bulk of the code and before the decompressed code starts to be executed:

```

IMB EQU 0xF00000
      .
      .
; copy and decompress bulk of code
    SWI IMB
; start of decompressed code

```

Chapter 8

Bus Interface

The ARM10 processor is designed to be used within larger chip designs using the *Advanced Microcontroller Bus Architecture (AMBA)*. The ARM10 processor uses the *AMBA High-performance Bus (AHB)* as its interface to memory and peripherals.

This chapter describes the features of the bus interface not covered in the *AMBA Specification*. It contains the following sections:

- *Bus features* on page 8-2
- *AMBA AHB signals* on page 8-3
- *Arbiter signals* on page 8-6
- *AHB control signals* on page 8-7
- *Timing* on page 8-9
- *Bus interface* on page 8-10.

8.1 Bus features

The ARM10 processor uses separate AHB bus interfaces for instructions and data:

- *Instruction Bus Interface Unit (IBIU)*
- *Data Bus Interface Unit (DBIU).*

Separate bus interfaces enhances the ability to fetch and execute instructions in parallel with a data cache miss. There is no sharing of any AHB signals between the two interfaces.

The ARM10 AHB interface is always driven. When either bus master is not granted the bus, that master drives zeros onto the bus to prevent bus contention. The ARM10 processor has unidirectional inputs, outputs, and control signals.

To ensure reuse of your design with other ARM processors, including different revisions, it is strongly recommended that you use fully AMBA-compliant peripherals and interfaces early in your design cycle. The AHB timings described in this chapter are examples only, and do not provide a complete list of all possible accesses. For further details on AMBA interface and integration, see the *AMBA Specification*.

8.2 AMBA AHB signals

Table 8-1 lists the AMBA AHB signals.

Table 8-1 AMBA AHB signals

Name	Direction	Description
HADDRI[31:0]	Output	IBIU address bus.
HADDRD[31:0]	Output	DBIU address bus.
HBURSTI[2:0]	Output	IBIU burst transfer type: 000 = single transfer 010 = four-beat wrapping burst.
HBURSTD[2:0]	Output	DBIU burst transfer type: 000 = single transfer 010 = four-beat wrapping burst 011 = four-beat incrementing burst.
HCLK	Input	Clock source. This clock times all bus transfers. All signal timings are related to the rising edge of HCLK .
HCLKOUT	O	External tap from internal HCLK balanced clock tree.
HPROTI[3:0]	Output	IBIU protection control. Transfers are always opcode fetches: xxx0 = opcode fetch xxx1 = data access xx0x = user access xx1x = privileged access x0xx = not bufferable x1xx = bufferable 0xxx = not cachable 1xxx = cachable.
HPROTD[3:0]	Output	DBIU protection control. Transfers are always data accesses: xxx0 = opcode fetch xxx1 = data access xx0x = user access xx1x = privileged access x0xx = not bufferable x1xx = bufferable 0xxx = not cachable 1xxx = cachable.
HRDATAI[63:0]	Input	Read IBIU data bus. Transfers data and instructions from bus slaves to instruction side bus master in read operations.

Table 8-1 AMBA AHB signals (continued)

Name	Direction	Description
HRDATAD[63:0]	Input	Read DBIU data bus. Transfers data from bus slaves to data side bus master in read operations.
HREADYI	Input	Slave ready. HIGH means transfer finished. Can be driven LOW to extend transfer.
HREADYD	Input	Slave ready. HIGH means transfer finished. Can be driven LOW to extend transfer.
HRESETN	Input	Bus reset. This is the only active-LOW AHB signal.
HRESPI[1:0]	Input	Slave response to IBIU. Reflects status of transfer: 00 = OKAY 01 = ERROR 10 = RETRY 11 = SPLIT.
HRESPD[1:0]	Input	Slave response to DBIU. Reflects status of transfer: 00 = OKAY 01 = ERROR 10 = RETRY 11 = SPLIT.
HSIZEI[2:0]	Output	Size of IBIU transfer: 000 = byte (8 bits) 001 = halfword (16 bits) 010 = word (32 bits) 011 = doubleword (64 bits) 100 = 4 words (128 bits) 101 = 8 words (256 bits) 110 = 16 words (512 bits) 111 = 32 words (1024 bits).
HSIZED[2:0]	Output	Size of DBIU transfer: 000 = byte (8 bits) 001 = halfword (16 bits) 010 = word (32 bits) 011 = doubleword (64 bits) 100 = 4 words (128 bits) 101 = 8 words (256 bits) 110 = 16 words (512 bits) 111 = 32 words (1024 bits).

Table 8-1 AMBA AHB signals (continued)

Name	Direction	Description
HTRANSI[1:0]	Output	Selects type of IBIU transfer: 00 = IDLE 01 = BUSY (This signal is not used) 10 = NONSEQUENTIAL 11 = SEQUENTIAL.
HTRANSD[1:0]	Output	Reflects type of DBIU transfer: 00 = IDLE 01 = BUSY (This signal is not used) 10 = NONSEQUENTIAL 11 = SEQUENTIAL.
HWDATAD[63:0]	Output	DBIU write data bus. Transfers data from master to slaves in write operations.
HWRITEI	Output	IBIU transfer direction. HIGH means write transfer. LOW means read transfer.
HWRITED	Output	DBIU transfer direction. HIGH means write transfer. LOW means read transfer.

8.3 Arbiter signals

Table 8-2 lists the arbiter signals.

Table 8-2 Arbiter signals

Name	Direction	Description
HBUSREQD	Output	Request line from DBIU.
HBUSREQI	Output	Request line from IBIU.
HGRANTD	Input	AHB mastership granted to DBIU.
HGRANTI	Input	AHB mastership granted to IBIU.
HLOCKD	Output	Indicates sequence of locked DBIU transfers in SWP operations.
HLOCKI	Output	For AMBA compliance. Never asserted.

8.3.1 Arbiter interface

Figure 8-1 shows the connections between the arbiter and the BIUs.

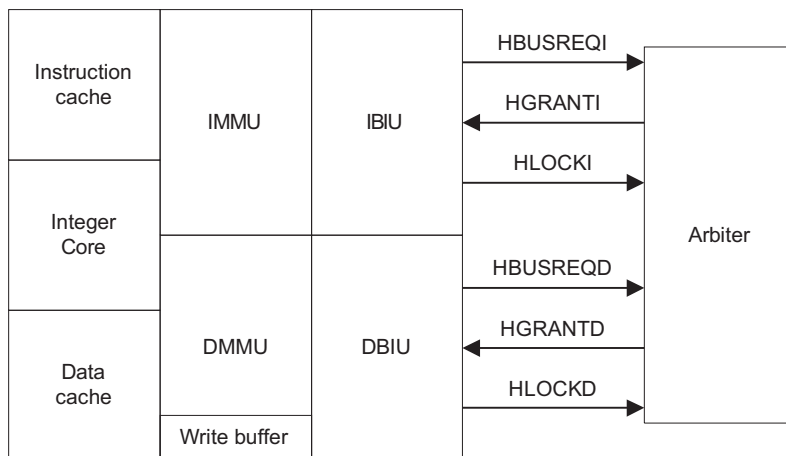


Figure 8-1 Arbiter-bus interface connections

8.4 AHB control signals

This section describes the ARM10 processor signals that control the AHB:

- **HTRANSI[1:0], HTRANSD[1:0]**
- **HSIZEI[2:0], HSIZED[2:0]**
- **HBURSTI[2:0], HBURSTD[2:0]**
- **HPROTI[3:0], HPROTD[3:0]**.

The descriptions in these sections apply to both versions of the signals listed above.

8.4.1 HTRANS[1:0]

The IBIU and DBIU use:

- 10 NONSEQ
- 11 SEQ (for linefills and bufferable STM instructions)
- 00 IDLE.

———— **Note** —————

01 BUSY is not used.

8.4.2 HSIZE[2:0]

HSIZE cannot be greater than 64 bits for a valid transfer. Table 8-3 lists typical transfer sizes.

Table 8-3 Transfer sizes

Type of transfer	Size of transfer	Comment
Linefills	64-bit	-
IBIU noncacheable reads	Depends on processor state	16-bit, 32-bit, or 64-bit
DBIU noncacheable reads	Depends on load instruction	Byte, halfword, word, or doubleword
DBIU noncacheable, nonbufferable writes	Depends on store instruction	Byte, halfword, word, or doubleword
Buffered writes	Depends on store instruction	Castouts are 64-bit

8.4.3 HBURST[2:0]

Burst lengths are shown in Table 8-4.

Table 8-4 BURST lengths

Encoding	Name	Description
000	SINGLE	Single nonsequential transfer
001	INCR	Incrementing burst of unspecified length
010	WRAP4	Four-beat wrapping burst

———— **Note** —————

In the case of a RETRY or SPLIT response during a linefill or castout, the remainder of the linefill or castout completes in nonsequential SINGLE transfers.

8.4.4 HPROT[3:0]

The values of the **HPROT** bits can be used in level 2 caches as shown in Table 8-5.

Table 8-5 Transfer attributes

Value	Meaning
HPROT0	0 = ICache linefill, core instruction fetch, or IMMU table walk 1 = DCache linefill, core load/store operation, or DMMU table walk
HPROT1	0 = User mode 1 = privileged mode
HPROT2	Depends on memory configuration: 0 = nonbufferable 1 = bufferable
HPROT3	Depends on memory configuration: 0 = noncachable 1 = cachable

8.5 Timing

The overall clocking scheme for the ARM10 processor is as follows:

- **HCLK** and **GCLK** must have coincident rising edges.
- **GCLK** can run at higher frequencies than **HCLK** if it is an integer multiple of **HCLK**.
- The integer unit, caches, MMUs, and any coprocessors run at **GCLK** speed.
- The AHB interface runs at **HCLK** speed, where $\mathbf{HCLK} = \mathbf{GCLK}/(1, 2, 3, 4, \dots)$ or $\mathbf{HCLK}:\mathbf{GCLK} = 1:N$ ($N = 1, 2, 3, 4, \dots$).

8.6 Bus interface

The bus interface is described in the following sections:

- *Topology* on page 8-11
- *Write buffer* on page 8-12.

The bus interface handles all data transfers and instruction transfers between the core clock domain and the AMBA bus clock domain. Any request from the prefetch unit or the LSU that has to go outside the ARM10 processor is handled by the bus interface in a way that is transparent to the prefetch unit and the LSU.

The following requests from the caches and MMUs drive the bus interface:

- page table walks (generated by the MMUs)
- noncachable reads
- nonbuffered writes
- linefills
- buffered writes
- CP15 write-buffer-related operations (empty write buffer and clean index).

In linefills, buffered writes are allowed to run underneath if there is room in the write buffer. Table 8-6 describes the request types.

Table 8-6 Blocking and nonblocking request types

Request type	Blocking or nonblocking
Page table walks (generated by the MMUs)	Blocking
Noncachable reads	Blocking
Nonbuffered writes	Blocking
Buffered writes	Blocking
CP15 write-buffer-related operations (empty and index clean)	Blocking
Linefills	Nonblocking

All of the request types except linefills are blocking requests. No other request can happen until the bus interface has acknowledged completion of the request. There is no priority assignment for these requests because no more than one blocking request can be asserted at any one time. It is possible for a request to the IBIU to be asserted simultaneously with a request to the DBIU. The AHB system arbiter determines priority in such cases.

For all requests, the bus interface registers the request and the associated data bits, protection bits, and address bits. The bus interface requests ownership of the AHB and, when it is granted ownership, it performs the appropriate transfer. When the transfer completes, the bus interface drops its request line and gives up ownership of the AHB.

Internal bandwidth between the bus interface and the caches and MMUs is 64 bits. Typical sizes of requests are listed in Table 8-7. (See also Table 8-3 on page 8-7.) On the AHB, **HWDATA** and **HRDATA** are 64 bits wide.

Table 8-7 Typical bus interface request sizes

Bus interface	Page walk	Linefill	Noncacheable read	Write
IBIU	32	64	16, 32, 64	-
DBIU	32	64	8, 16, 32, 64	8, 16, 32, 64

8.6.1 Topology

The bus interface consists of two completely separate blocks:

- The IBIU handles all instruction fetches and linefills.
- The DBIU performs all data loads and stores.

Both the IBIU and DBIU perform page table walks for their respective MMUs when required.

Figure 8-2 on page 8-12 shows the structure of the bus interface. The DBIU is on the left with control, read, write, and address data-path. The IBIU on the right has a read and an address data-path only because no writes ever happen on the instruction side. Both the IBIU and the DBIU have a similar layer for transferring data or instructions to and from the **HCLK** domain and further on to the rest of the AMBA system. The arrows illustrate the flow of requests and data or instructions.

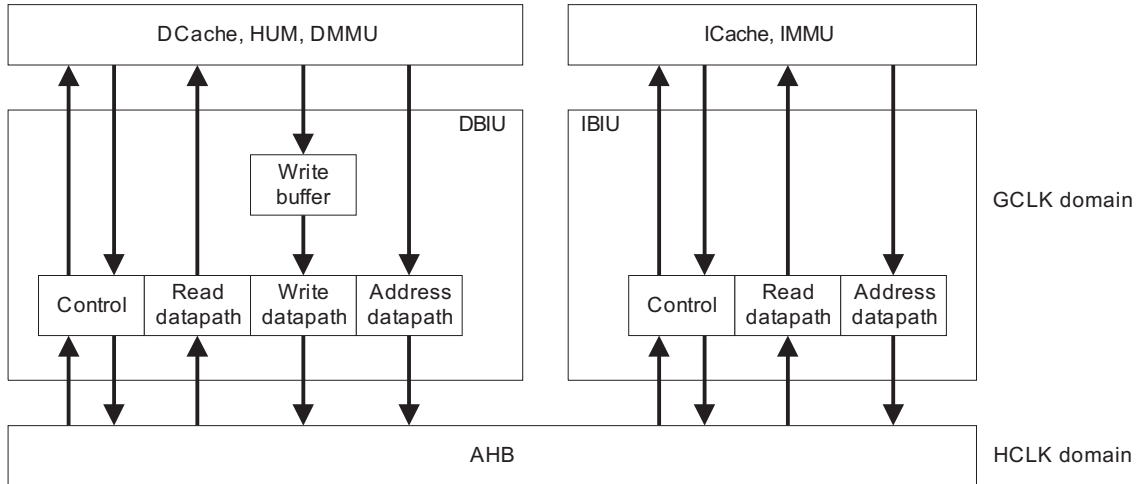


Figure 8-2 Bus interface block diagram

The DBIU and the IBIU are independent of each other. There is no efficient way of communicating between the data and the instruction side of the ARM processor, making self-modifying code difficult to accommodate.

8.6.2 Write buffer

The write buffer is the part of the DBIU used for capturing buffered writes sent from the LSU at **GCLK** speed and then writing the data back to main memory at **HCLK** speed at a later time.

The write buffer also buffers castout victims from the DCache.

The write buffer has two logical blocks:

- The circular queue write buffer
- The FIFO castout buffer.

Figure 8-3 on page 8-13 shows the structure of the circular queue write buffer and FIFO castout buffer.

Each entry in the circular queue write buffer and the FIFO castout buffer has 64 bits of data, including a 32-bit address and six protection bits. No merging of writes takes place during data insertion. Two bytes written to successive addresses take up two entries.

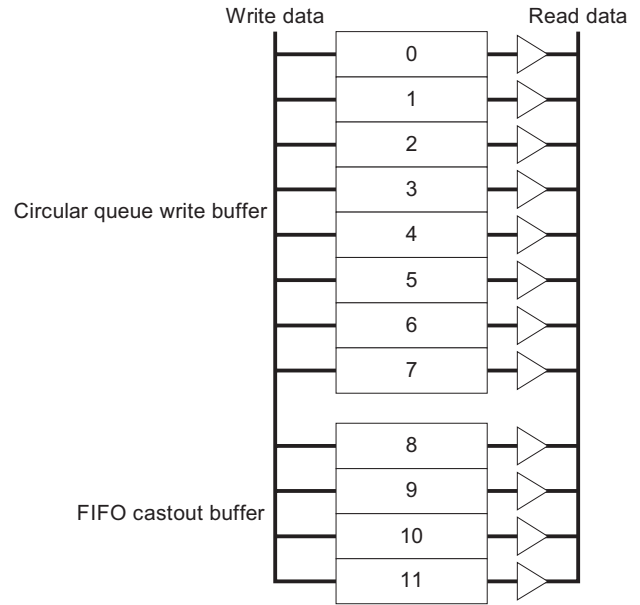


Figure 8-3 Write buffer and castout buffer

Circular queue write buffer

Two pointers are required for the write buffer:

- The *front of queue* pointer points to the next entry to write.
- The *back of queue* pointer points to the next location to read when emptying the write buffer.

The back of queue pointer always chases the front of queue pointer. The write buffer is empty when both point to the same location.

To minimize interrupt latency, the size of the write buffer can be halved. This change in size is transparent to the DBIU. It is controlled by the fast interrupt bit, FI, in control register 1 in CP15.

In the MMU translation table, three combinations of the cachable and bufferable bits, C and B, produce a buffered write, as Table 8-8 shows.

Table 8-8 Cachable and bufferable bits in buffered writes

C	B	Description
0	1	Noncachable buffered write
1	0	Write-through, considered a buffered write
1	1	Write-back, considered a buffered write

A noncachable, nonbufferable write is the only type of nonbuffered write that bypasses the write buffer.

Noncachable, nonbufferable writes are always single nonsequential writes to memory.

The DBIU empties the write buffer dynamically when either:

- it samples a blocking request and, to maintain memory coherency, must empty a number of entries prior to this request before servicing the blocking request
- it detects that the write buffer is no longer empty.

———— **Note** —————

Even with dynamic emptying, the write buffer can become full and stall the LSU. The conditions for this occurring are a combination of **HCLK:GCLK** ratio and the frequency with which buffered writes are inserted into the write buffer.

FIFO castout buffer

The castout buffer contains victims from the data cache. The castout buffer is a FIFO because the amount of data, four doublewords, is known, and the order of the data is fixed. It is always a four-beat wrapping order, where the address wraps on 32-byte boundaries. Data for the castout buffer is always inserted from location 11 to location 8. This is also the order in which data is read back out when emptying the castout buffer.

Victims from the data cache are always 64-bit aligned and take up four entries, which is exactly the size of the castout buffer.

Chapter 9

Coprocessor Interface

This chapter contains information about the coprocessor interface. It contains the following sections:

- *About the coprocessor interface* on page 9-2
- *Coprocessor interface signals* on page 9-3
- *Design considerations* on page 9-5
- *Parallel execution* on page 9-8
- *Rules for the interface* on page 9-9
- *Pipeline signal assertion* on page 9-10
- *Instruction issue* on page 9-11
- *Hold signals* on page 9-19
- *Instruction cancelation* on page 9-38
- *Bounced instructions* on page 9-45
- *Data buses* on page 9-50.

9.1 About the coprocessor interface

The coprocessor interface enables you to attach multiple coprocessors (CPs) to the ARM10 processor. To limit the number of connections required by the interface, each CP tracks the progress of instructions in the ARM10 pipeline.

To enable optimum performance from CPs, the ARM10 processor issues CP instructions as early as possible. This means that the instructions are issued speculatively, and they can be canceled later in the pipeline if, for example, an exception or branch misprediction occurs. As a result, CPs must be able to cancel instructions in late stages of the ARM10 pipeline.

Simple CPs only track the ARM10 pipeline until it is certain that a given instruction does not get canceled. At this point the CP starts to execute the instruction. More complex CPs make extensive use of the early issue of the instruction.

At certain points in the pipeline, a CP sends back signals to the ARM10 processor. These can indicate that the CP requires more time to execute or to indicate that the undefined instruction exception must be taken.

9.1.1 CP pipeline

The CP pipeline runs one cycle behind the ARM10 pipeline. This enables pipeline holds from the ARM10 processor to be registered before they are sent to the CPs. Figure 9-1 shows the ARM10 and CP pipeline stages.

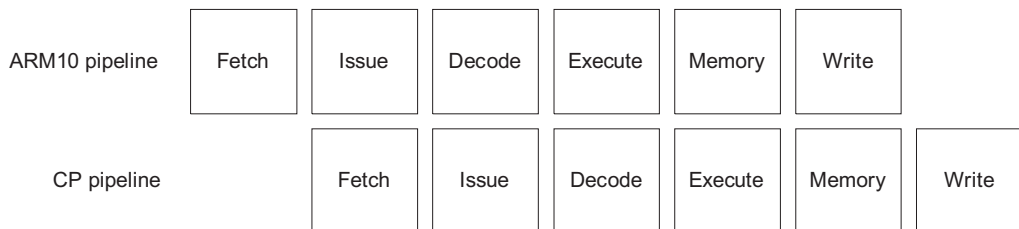


Figure 9-1 ARM10 and CP pipeline stages

9.2 Coprocessor interface signals

This section divides the CP signals according to function:

- *ARM10 instruction progression signals*
- *ARM10 instruction cancelation signals*
- *CPBOUNCEE* on page 9-4
- *Busy-waiting instruction* on page 9-4
- *CP data buses* on page 9-4
- *CP control signals* on page 9-4.

9.2.1 ARM10 instruction progression signals

The signals that indicate instruction progression are:

CPINSTRV	Valid CP instruction in ARM10 Issue stage.
CPVALIDD	Valid CP instruction in ARM10 Decode stage.
ASTOPCPD	ARM10 stalled in Decode stage in previous cycle.
CPBUSYE	ARM10 stalled in Execute stage in previous cycle.
LSHOLDCPE	ARM10 LSU stalled in Execute stage in previous cycle.
LSHOLDCPM	ARM10 LSU stalled in Memory stage in previous cycle.

9.2.2 ARM10 instruction cancelation signals

Two signals indicate ARM10 instruction cancelation:

ACANCELCP

Cancels only the instruction that was in ARM10 Execute stage in the previous cycle.

AFLUSHCP

Cancels all the instructions back from the one that was in ARM10 Execute stage in the previous cycle. **AFLUSHCP** overrides **STOP** and **VALID** signals from the ARM10 processor and causes **BUSY** signals to be deasserted in the following cycle.

9.2.3 CPBOUNCEE

The signal that indicates whether a CP can execute an instruction is:

CPBOUNCEE Takes the undefined instruction trap for the instruction that is in the ARM10 Execute stage.

9.2.4 Busy-waiting instruction

The signal that indicates whether a CP requires more time to process an instruction is:

CPBUSYE Busy-wait (stall) the ARM10 Execute stage.

———— **Note** —————

The ARM10 processor has **CPBUSYD1** and **CPBUSYD2** inputs. These are reserved for future expansion. Tie these off to a logic 0.

9.2.5 CP data buses

There are two 64-bit CP data buses:

- **STCMRCDATA** carries data from a CP to the ARM10 processor
- **LDCMRCDATA** carries data from the ARM10 processor to a CP.

9.2.6 CP control signals

CPLSLEN, **CPLSSWP**, and **CPLSDBL** are signals driven by a CP to the ARM10 processor on load/store CP instructions. They carry additional information about:

- the length of the transfer
- if upper and lower half of the data bus must be swapped before being written
- if the load/store request is for double word data.

9.3 Design considerations

This section outlines CP interface design considerations for single and multiple CPs.

9.3.1 Input and output timing

Almost all the signals on both sides of the interface must be driven straight out of registers. This is necessary because there is very little timing slack in the interface. There is very little timing slack because as few cycles as practical have been used to process a given CP instruction. This enables very high performance CPs to be built. If performance is not an issue, then timing across the interface can be greatly simplified by stalling all CP instructions in situations where timing is an issue.

9.3.2 ARM10 processor inputs and outputs

Outputs driven from the ARM10 processor go to all the CPs in the system. The inputs to the ARM10 processor from all the CPs are ANDed or ORed together before they are used. As a result, the ARM10 processor cannot tell which CP is driving its inputs. Figure 9-2 on page 9-6 shows **CPBUSYE** and **CPBOUNCEE** as examples of ARM10 coprocessor input gating. The problem of multiple CPs driving a signal at the same time is avoided, because there can only be one CP instruction in each ARM10 pipeline stage. So only one CP can own the instruction in that stage and can drive the associated signals.

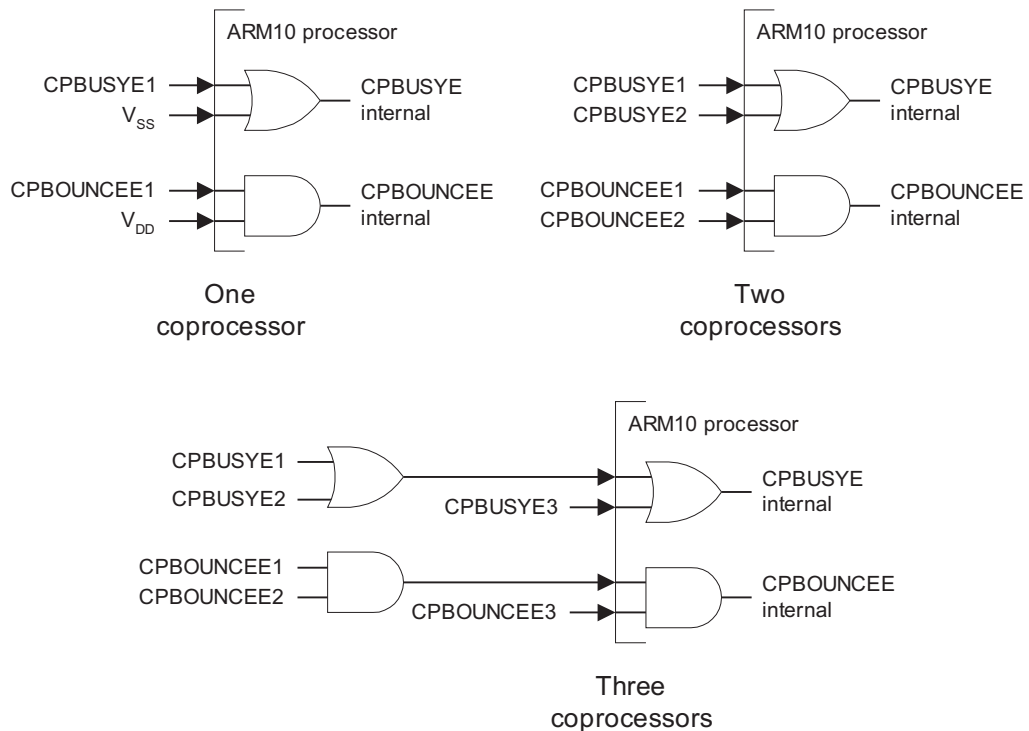


Figure 9-2 ARM10 coprocessor inputs

The ARM10 processor has control inputs for up to two external coprocessors. The two sets of inputs are differentiated by appending a *1* or a *2* to the signal name. Inputs that are not used must be tied off. By convention, single-coprocessor systems use the *1* inputs and tie off the *2* inputs. Adding more than two external coprocessors requires external gating.

Any system with more than one external coprocessor requires external gating for the **STCMRCDATA** bus inputs to the ARM10 processor. This is to avoid the necessity of routing 64-bit buses to the ARM10 processor.

9.3.3 CP input loadings

When a CP does not own the instruction associated with an ANDed signal it must drive the signal HIGH. When a CP does not own the instruction associated with an ORed signal it must drive the signal LOW. The ARM10 processor drives instruction, data, and

control outputs to all CPs, so the loading on these signals might become an issue in multiple-CP systems. Keep CP input loadings low, and buffer these signals where appropriate.

9.3.4 Combining outputs from multiple CPs

Outputs from all the CPs are ANDed or ORed together before they are used in the ARM10 processor. The AND and OR gates can be placed in the level of the design instantiating the ARM10 processor and the CPs. To aid timing for control signals, there is one level of ANDing and ORing inside the ARM10 processor. The ARM10 processor implements the ANDing and ORing necessary on the control signals of up to two external CPs. For more than two CPs, external gates must be used to OR the hold signals from the external CP into the existing inputs.

Although the ARM10 processor implements the necessary inputs for only two external CPs, this does not have to be the limiting factor in a system with three or more CPs. In such a system, the wire delays from the farthest CP probably balance the time required to AND or OR the control signal from the closer CPs. For systems with more than one CP, external gates are always required for the CP **STCMRCDATA** bus. These are not included in the ARM10 design as this would have forced the entire bus to be duplicated on the interface. Also, the freedom to place the gates anywhere in the top-level design helps with floor planning of the bus route.

9.3.5 CP ID number

The ARM10 processor issues all CP instructions to all the CPs. Each CP in the system has a unique, hardwired ID number from 0 to 15. Every CP instruction includes a CP number.

Only the CP whose ID number corresponds to the number in the CP instruction can accept the instruction. To accept an instruction, a CP must pull **CPBOUNCEE** LOW at the right time. If no CP pulls **CPBOUNCEE** LOW, then the instruction is *bounced*. That is, the ARM10 processor takes the undefined instruction trap. This enables error trapping or software emulation of a CP not present in the system.

A CP does not have to accept an instruction even if its ID corresponds to the CP number in the instruction. This is used in cases where some of the CP instructions are handled in hardware and some are handled in software.

9.4 Parallel execution

Initially, instructions progress along the ARM10 pipeline and CP pipeline in lockstep. A CP instruction moves along the ARM10 pipeline as if it were a single-cycle instruction. When the first cycle of the instruction traverses the entire length of the ARM10 pipeline, one of three things can occur:

- If the instruction is complete in the CP pipeline then it is retired in both pipelines.
- If the CP instruction is a multicycle data processing type, then the ARM10 processor and CP pipelines are decoupled. The instruction continues to iterate in the CP but is retired in the ARM10 pipeline. Once the pipelines are decoupled, the ARM10 processor cannot cancel the instruction, and the CP must complete the instruction. While the CP is working, the ARM10 processor continues to execute the following instruction stream and issues any CP instructions it hits. The CP can hold up any following CP instructions as necessary. The ARM10 processor is not explicitly signaled when the CP completes the instruction. The CP usually holds up any following instruction that is dependent on a prior instruction.
- If the CP instruction is a multicycle load or store type, then the ARM10 ALU pipeline and CP pipelines are decoupled, but the ARM10 LSU pipeline and CP pipeline remain coupled. The instruction continues to iterate in the CP and the ARM10 LSU pipeline but is retired in the ARM10 ALU pipeline. When the ARM10 ALU pipeline is decoupled, the ARM10 processor cannot cancel the instruction, and the CP must complete the instruction. While the CP and LSU are working, the ARM10 ALU pipeline continues to execute the following instruction stream and issues any CP instructions it hits. Load and store instructions stall in Decode, but data processing instructions execute if possible. Even the CP doing the load or store can run a data processing instruction in parallel if it supports this functionality. If it does not, then it must hold up the data processing instruction until the load or store instruction is complete.

Simple CPs only have to use the first of these mechanisms. They can execute multicycle instructions by holding up the ARM10 pipeline until they complete. In some systems this has a significant impact on performance.

9.5 Rules for the interface

The following rules apply to the CP pipeline and CP interface:

- No two CPs can have an instruction in the same ARM10 pipeline stage. That is, a CP instruction in a particular ARM10 pipeline stage is associated with one, and only one, CP.
- Each CP output signal is associated with one ARM10 pipeline stage. The CP that owns the instruction in that stage drives the signal.
- Outputs from the ARM10 processor must enable the CPs to track the ARM10 pipeline well enough for them to detect:
 - when to assert hold and bounce signals to ARM10 processor
 - which CP instruction that a cancel or flush signal applies to
 - when the instruction is committed and can no longer be canceled or flushed.
- A signal stalled by a hold signal becomes valid in the last cycle of the hold signal. Signals that override hold signals can be asserted at any time, and their effect must not be masked by the hold.

Internal design features of CPs might or might not conform to these rules.

9.6 Pipeline signal assertion

Table 9-1 shows where in the pipeline the coprocessor interface signals are active.

Table 9-1 Pipeline stages and active signals

	ARM10 pipeline		CP pipeline	
	Driven by ARM10	Driven by CP	Driven by ARM10	Driven by CP
CPVALIDD	Decode	-	Issue	-
CPLSLEN	-	Decode	-	Issue
CPLSSWP	-	Decode	-	Issue
CPLSDBL	-	Decode	-	Issue
CPINSTR	Issue	-	Fetch	-
CPINSTRV	Issue	-	Fetch	-
ASTOPCPD	Execute	-	Decode	-
CPBUSYE	-	Execute	-	Decode
CPLSBUSY	-	Execute	-	Decode
CPBOUNCEE	-	Execute	-	Decode
ASTOPCPE	Memory	-	Execute	-
ACANCELCP	Memory	-	Execute	-
AFLUSHCP	Memory	-	Execute	-
LSHOLDCPE	Memory	-	Execute	-
LSHOLDCPM	Write	-	Memory	-
STCMRCDATA	-	Execute	-	Decode
LDCMCRDAT A	Write	-	Memory	-

9.7 Instruction issue

CPINSTR, **CPINSTRV**, and **CPVALIDD** are the signals that control the issue of CP instructions from the ARM10 processor. These instructions go to all CPs at the same time. Only the CP that owns the instruction can drive control signals for that instruction back to the ARM10 processor.

The following sections describe these signals:

- *CPINSTR*
- *CPINSTRV* on page 9-12
- *CPVALIDD* on page 9-14
- *Example of instruction issue* on page 9-15
- *CPLSLEN, CPLSSWP, and CPLSDL* on page 9-16.

9.7.1 CPINSTR

Instructions are issued to all CPs during the ARM10 Issue stage, which is in the CP Fetch stage. The instructions are sent over a dedicated 26-bit bus, **CPINSTR**.

Usually, **CPINSTR** is only driven when there is a valid CP instruction in the ARM10 Issue stage. Occasionally, it might be driven in error because of an instruction that causes a Prefetch Abort or a branch that is incorrectly predicted. In these cases the value driven onto **CPINSTR** might decode to anything, including a CP instruction. However the instruction is still not valid because it was fetched erroneously.

CPINSTRV and **CPVALIDD** give more information about the validity of the instruction. Table 9-2 on page 9-12 shows interactions of **CPINSTR** with other signals.

The ARM10 processor drives **CPINSTR** in the ARM10 Issue stage and the CP Fetch stage.

Table 9-2 CPINSTR interactions with other signals

Signal	Interactions with CPINSTR
ASTOPCPD	Treat CPINSTR as invalid this cycle. Use its value only in the last interlocked cycle, that is, the cycle in which ASTOPCPD and all other relevant holds go LOW. The value of CPSINTR might change while ASTOPCPD is asserted if an exception or mispredicted branch occurs. Also, the prefetch unit might place a valid instruction in the Issue stage under an interlock, causing an invalid instruction on CPINSTR and CPINSTRV to change to a valid one while ASTOPCPD is asserted.
ASTOPCPE	Treat CPINSTR as invalid this cycle. Use its value only in the last interlocked cycle, that is, the cycle in which ASTOPCPE and all other relevant holds go LOW. The value of CPSINTR might change while ASTOPCPE is asserted if an exception or mispredicted branch occurs. Also, the prefetch unit might place a valid instruction in the Issue stage under an interlock, causing an invalid instruction on CPINSTR and CPINSTRV to change to a valid one while ASTOPCPE is asserted.
LSHOLDCPE	None.
CPBUSYE	Treat CPINSTR as invalid this cycle. Use its value only in the last interlocked cycle, that is, the cycle in which CPBUSYE and all other relevant holds go LOW. The value of CPSINTR might change while CPBUSYE is asserted if an exception or mispredicted branch occurs. Also, the prefetch unit might place a valid instruction in the Issue stage under an interlock, causing an invalid instruction on CPINSTR and CPINSTRV to change to a valid one while CPBUSYE is asserted.
LSHOLDCPM	None.
ACANCELCP	None.
AFLUSHCP	Invalidates instruction on CPINSTR .
CPBOUNCEE	None.

9.7.2 CPINSTRV

CPINSTR and **CPINSTRV** are the only CP interface signals that are driven in the ARM10 Issue stage. **CPINSTRV** indicates that **CPINSTR** carries an instruction worth decoding. The fact that **CPINSTRV** is asserted is not a guarantee that **CPINSTR** carries a valid CP instruction. **CPINSTRV** going LOW is a guarantee the **CPINSTR** does not carry a valid CP instruction.

CPINSTRV is a useful hint. It can be used to save power by not decoding bad instructions. To save power all bits of **CPINSTR** are also driven to 0 when **CPINSTRV** is LOW. This behavior must not be relied upon for correct function.

If **CPINSTR** carries a valid CP instruction, **CPINSTRV** does not guarantee that it will be executed. There are some cases where **CPINSTRV** is asserted for instructions that turn out to be invalid. Prefetch aborted instructions and instructions following mispredicted branches are examples of this. Not enough is known about the instruction in the ARM10 Issue stage to make **CPINSTRV** a definite indicator of a valid instruction. More is known in the ARM10 Decode stage and the signal **CPVALIDD** is used to confirm that an instruction is valid. Table 9-3 shows interactions of **CPINSTRV** with other signals.

The ARM10 processor drives **CPINSTRV** in the ARM10 Issue stage and the CP Fetch stage.

Table 9-3 CPINSTRV interactions with other signals

Signal	Interactions with CPINSTRV
ASTOPCPD	Treat CPINSTRV as invalid this cycle. Use its value only in the last interlocked cycle, that is, the cycle in which ASTOPCPD and all other relevant holds go LOW. The value of CPSINTRV might change while ASTOPCPD is asserted if an exception or mispredicted branch occurs. Also, the prefetch unit might place a valid instruction in the Issue stage under an interlock, causing an invalid instruction on CPINSTR and CPINSTRV to change to a valid one while ASTOPCPD is asserted.
ASTOPCPE	Treat CPINSTRV as invalid this cycle. Use its value only in the last interlocked cycle, that is, the cycle in which ASTOPCPE and all other relevant holds go LOW. The value of CPSINTRV might change while ASTOPCPE is asserted if an exception or mispredicted branch occurs. Also, the prefetch unit might place a valid instruction in the Issue stage under an interlock, causing an invalid instruction on CPINSTR and CPINSTRV to change to a valid one while ASTOPCPE is asserted.
LSHOLDCPE	None.
CPBUSYE	Treat CPINSTR as invalid this cycle. Use its value only in the last interlocked cycle, that is, the cycle in which CPBUSYE and all other relevant holds go LOW. The value of CPSINTRV might change while CPBUSYE is asserted if an exception or mispredicted branch occurs. Also, the prefetch unit might place a valid instruction in the Issue stage under an interlock, causing an invalid instruction on CPINSTR and CPINSTRV to change to a valid one while CPBUSYE is asserted.
LSHOLDCPM	None.

Table 9-3 CPINSTRV interactions with other signals (continued)

Signal	Interactions with CPINSTRV
ACANCELCP	None.
AFLUSHCP	Invalidates instruction.
CPBOUNCEE	None.

9.7.3 CPVALIDD

Not enough is known about the instruction in the ARM10 Issue stage to make **CPINSTRV** a definite indicator of a valid instruction. More is known in the ARM10 Decode stage, and the signal **CPVALIDD** can confirm that an instruction is valid. **CPVALIDD** goes HIGH during the ARM10 Decode stage to confirm an instruction is valid. **CPVALIDD** does not guarantee execution of the instruction, because the instruction might get canceled or flushed (see *ACANCELCP* on page 9-38 and *AFLUSHCP* on page 9-42). Table 9-4 shows interactions of **CPVALIDD** with other signals.

The ARM10 processor drives **CPVALIDD** in the ARM10 Decode stage and the CP Issue stage.

Table 9-4 CPVALIDD interactions with other signals

Signal	Interactions with CPVALIDD
ASTOPCPD	Treat CPVALIDD as invalid this cycle. Use its value only in the last interlocked cycle, that is, the cycle in which ASTOPCPD and all other relevant holds go LOW. The value of CPVALIDD might change while ASTOPCPD is asserted if an exception or mispredicted branch occurs. Also, the prefetch unit might place a valid instruction in the Issue stage under an interlock, causing an invalid instruction on CPINSTR and CPINSTRV to change to a valid one while CPVALIDD is asserted.
ASTOPCPE	Treat CPVALIDD as invalid this cycle. Use its value only in the last interlocked cycle, that is, the cycle in which ASTOPCPE and all other relevant holds go LOW. The value of CPVALIDD might change while ASTOPCPE is asserted if an exception or mispredicted branch occurs. Also, the prefetch unit might place a valid instruction in the Issue stage under an interlock, causing an invalid instruction on CPINSTR and CPINSTRV to change to a valid one while CPVALIDD is asserted.
LSHOLDCPE	None.

Table 9-4 CPVALIDD interactions with other signals (continued)

Signal	Interactions with CPVALIDD
CPBUSYE	Treat CPVALIDD as invalid this cycle. Use its value only in the last interlocked cycle, that is, the cycle in which CPBUSYE and all other relevant holds go LOW. The value of CPVALIDD might change while CPBUSYE is asserted if an exception or mispredicted branch occurs. Also, the prefetch unit might place a valid instruction in the Issue stage under an interlock, causing an invalid instruction on CPINSTR and CPINSTRV to change to a valid one while CPBUSYE is asserted.
LSHOLDCPM	None.
ACANCELCP	None.
AFLUSHCP	Invalidates instruction.
CPBOUNCEE	None.

9.7.4 Example of instruction issue

In Figure 9-3 on page 9-16, instructions 1 and 2 drive **CPINSTR**. **CPINSTRV** initially indicates that both instructions 1 and 2 are valid, but **CPVALIDD** indicates that only instruction 1 is valid. After that, instructions 3 and 4 are not valid CP instructions, so **CPINSTRV** and **CPVALIDD** are kept LOW. The numbers in the waveforms show which instruction owns the signal at that time. For example, instruction 1 owns **CPVALIDD** at edge T3. Instruction 2 owns **CPVALIDD** at edge T4. A CP registers the instruction 1 value at T3 and the instruction 2 value at T4.

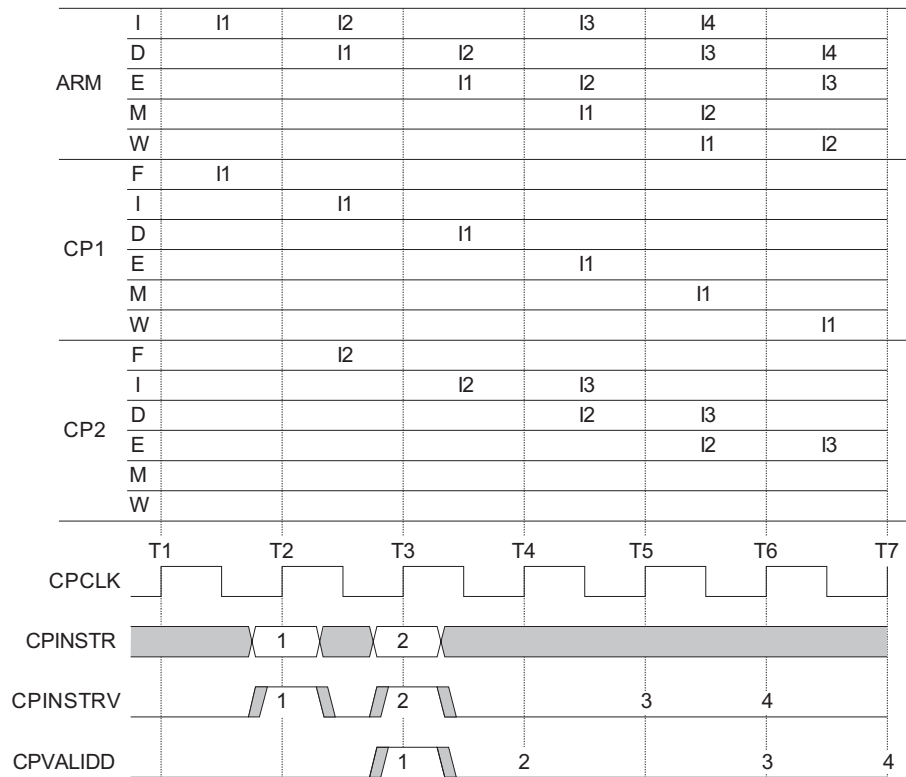


Figure 9-3 Instruction issue example

9.7.5 CPLSLEN, CPLSSWP, and CPLSDBL

A CP drives the **CPLSLEN**, **CPLSSWP**, and **CPLSDBL** signals to the ARM10 processor on load/store CP instructions. They indicate:

- the length of the transfer
- if upper and lower half of the data bus must be swapped before being written
- if the load/store request is for double-precision data.

CPLSLEN

CPLSLEN indicates the number of 32-bit data items to be transferred for the corresponding load/store CP instruction. Driving a 1 on this bus represents a single load or store data item being transferred. **CPLSLEN** must be driven with 0 if the CP is not

processing an instruction. If **ASTOPCPD** is asserted due to a hold in the ARM10 Decode stage, the **CPLSLEN** value is retained by the ARM10 processor. Table 9-5 describes the interactions of **CPLSLEN** with other signals.

The CP drives **CPLSLEN** in the CP Issue stage and the ARM10 Decode stage.

Table 9-5 CPLSLEN interactions with other signals

Signal	interactions with CPLSLEN
ASTOPCPD	CPLSLEN is registered with ASTOPCPD
ASTOPCPE	None
LSHOLDCPE	None
CPBUSYE	None
LSHOLDCPM	None
ACANCELCP	None
AFLUSHCP	Invalidates instruction
CPBOUNCEE	None

CPLSSWP

CPLSSWP indicates that the upper and lower data words on **LDCMCRDATA** and **STCMRCDATA** buses must be swapped by the ARM10 processor before being written. If **ASTOPCPD** is asserted due to a hold in the ARM10 Decode stage, the **CPLSSWP** value is retained by the ARM10 processor. Table 9-6 describes the interactions of **CPLSSWP** with other signals.

The CP drives **CPLSSWP** in the CP Issue stage and the ARM10 Decode stage.

Table 9-6 CPLSSWP interactions with other signals

Signal	Interactions with CPLSSWP
ASTOPCPD	CPLSSWP is registered with ASTOPCPD
ASTOPCPE	None
LSHOLDCPE	None
CPBUSYE	None
LSHOLDCPM	None

Table 9-6 CPLSSWP interactions with other signals (continued)

Signal	Interactions with CPLSSWP
ACANCELCP	None
AFLUSHCP	Invalidates instruction
CPBOUNCEE	None

CPLSDBL

CPLSDBL indicates that the load/store CP instruction involves a double word transfer. That is, a 64-bit quantity is being transferred. If **ASTOPCPD** is asserted due to a hold in the ARM10 Decode stage, the **CPLSDBL** value is retained by the ARM10 processor. Table 9-7 describes the interactions of **CPLSDBL** with other signals.

The CP drives **CPLSDBL** in the CP Issue stage and the ARM10 Issue stage.

Table 9-7 CPLSDBL interactions with other signals

Signal	Interactions with CPLSSWP
ASTOPCPD	CPLSDBL is registered with ASTOPCPD
ASTOPCPE	None
LSHOLDCPE	None
CPBUSYE	None
LSHOLDCPM	None
ACANCELCP	None
AFLUSHCP	Invalidates instruction
CPBOUNCEE	None

9.8 Hold signals

The following sections describe hold signals:

- *ASTOPCPD* on page 9-20
- *ASTOPCPE* on page 9-22
- *ASTOPCPE example* on page 9-23
- *LSHOLDCPE* on page 9-25
- *Example of LSHOLDCPE* on page 9-25
- *LSHOLDCPM* on page 9-27
- *CPBUSYE* on page 9-29
- *CPBUSYE example* on page 9-29
- *CPBUSYE and ASTOPCPD interaction* on page 9-30
- *ASTOPCPD with CPBUSYE* on page 9-31
- *CPBUSYE and ASTOPCPE interaction* on page 9-32
- *ASTOPCPE with CPBUSYE* on page 9-33
- *CPLSBUSY* on page 9-37.

The pipeline hold signals from the ARM10 processor keep the CP pipeline in lockstep with the ARM10 processor. Pipeline hold signals from the CPs hold up the ARM10 processor to give more time to execute an instruction. To avoid a deadlock, it is important that both sides do not factor their hold inputs back into their hold outputs. Table 9-8 summarizes the hold signals.

The hold signals are usually timing-critical. They factor huge fanout terms into pipeline holds. In high-performance systems, they must come straight out of registers in the driving block.

Table 9-8 Hold signals summary

Signal	From	To	ARM10 stage	CP stage	Comments
ASTOPCPD	ARM10	All CPs	Decode + 1	Decode	Hold CP in CP Decode because ARM10 is held in ARM10 Decode
ASTOPCPE	ARM10	All CPs	Execute + 1	Execute	Hold CP in CP Execute because ARM10 is held in ARM10 Execute
LSHOLDCPE	ARM10	All CPs	Execute + 1	Execute	Hold CP data transfers in CP Execute because LSU is held in ARM10 Execute

Table 9-8 Hold signals summary (continued)

Signal	From	To	ARM10 stage	CP stage	Comments
LSHOLDCPM	ARM10	All CPs	Memory + 1	Memory	Hold CP data transfers in CP Memory because LSU is held in ARM10 Memory
CPBUSYE	Each CP	Other CPs and ARM10	Execute	Decode	Hold ARM10 processor in ARM10 Execute
CPLSBUSY	Each CP	Other CPs	-	Decode	Holds other CPs in CP Issue

9.8.1 ASTOPCPD

ASTOPCPD indicates that the instruction in the ARM10 Decode stage did not progress into the ARM10 Execute stage in the previous cycle. It is driven out of a register following the ARM10 Decode stage. If **ASTOPCPD** is asserted, CPs must hold their Decode, Issue, and Fetch stages. The logic in these stages must keep reevaluating because **CPINSTR**, **CPINSTRV**, and **CPVALIDD** might change. Only the cycle in which **ASTOPCPD** is deasserted can be considered a valid cycle. Table 9-9 shows the interactions of **ASTOPCPD** with other signals.

The ARM10 processor drives **ASTOPCPD** in the ARM10 Execute stage and the CP Decode/CP Decode + 1 stage.

Table 9-9 ASTOPCPD interactions with other signals

Signal	Interactions with ASTOPCPD
ASTOPCPE	ASTOPCPD is usually asserted when ASTOPCPE is asserted.
LSHOLDCPE	ASTOPCPD is asserted with LSHOLDCPE when the pipelines are in lockstep. Pipelines are in lockstep unless the CP instruction has already retired from the ARM10 pipeline and is now transferring data from the LSU for a load/store multiple.
CPBUSYE	The ARM10 processor ignores CPBUSYE if ASTOPCPD is already asserted. ASTOPCPD is not asserted if a valid CPBUSYE (ASTOPCPE LOW) was received in the previous cycle.
LSHOLDCPM	None.
ACANCELCP	None.
AFLUSHCP	Flush invalidates ASTOPCPD .
CPBOUNCEE	None.

In Figure 9-4 **ASTOPCPD** is used to indicate that instruction 1 stalled in the ARM10 Decode stage for one cycle. The following values of **CPINSTR**, **CPINSTRV**, and **CPVALIDD** are invalid in all but the last cycle that was interlocked. **ASTOPCPD** is LOW as instruction 2 leaves the Decode stage indicating that it was not held up. The numbers in waveforms show which instruction owns the signal at that time.

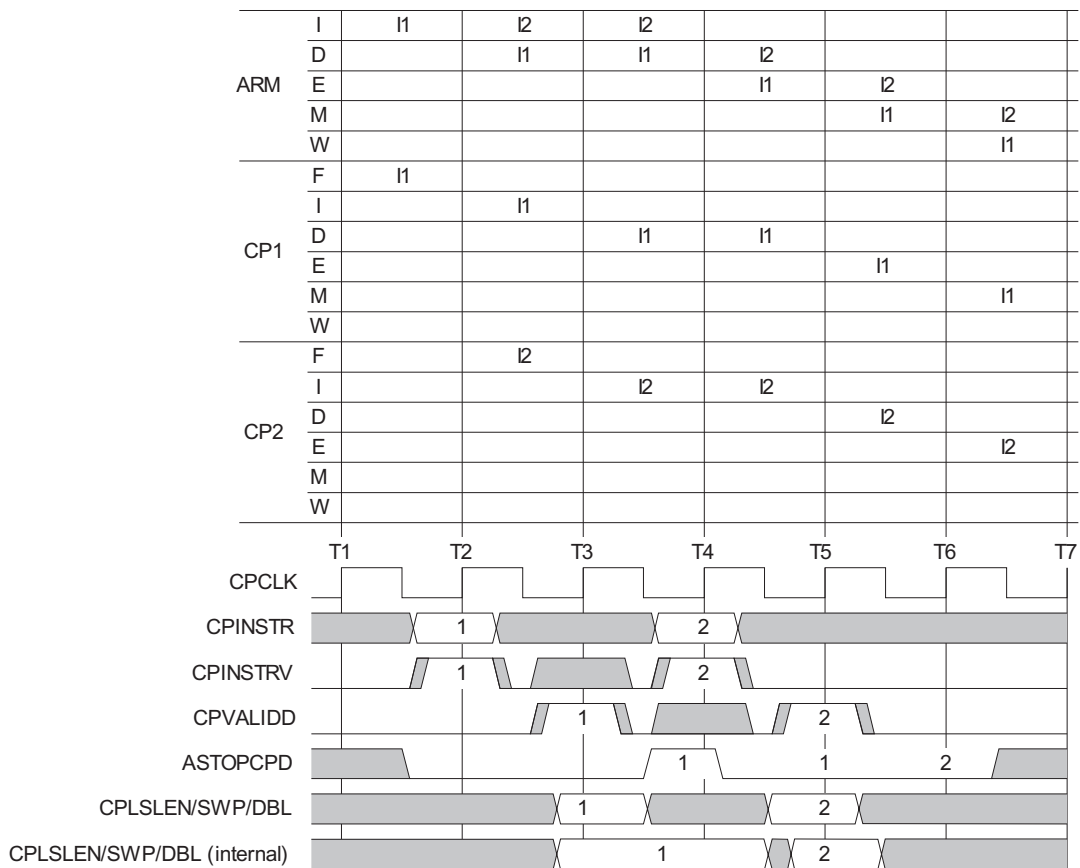


Figure 9-4 **ASTOPCPD** example

CPLSLEN, **CPLSSWP**, and **CPLSDBL** for a given instruction are driven from a CP in the cycle before **ASTOPCPD** is driven from the ARM10 processor, so the ARM10 processor must register the value of **CPLSLEN** and **CPLSSWP** and **CPLSDBL** if it is about to drive an **ASTOPCPD**.

9.8.2 ASTOPCPE

ASTOPCPE indicates that the instruction in the ARM10 Execute stage did not progress into the ARM10 Memory stage in the previous cycle. It is driven out of a register following the ARM10 Execute stage. If **ASTOPCPE** is asserted, CPs must hold their Execute, Decode, Issue, and Fetch stages. The logic in these stages must keep reevaluating as **CPINSTR**, **CPINSTRV**, and **CPVALIDD** might change. Only the cycle where **ASTOPCPE** is deasserted is a valid cycle. **AFLUSHCP** overrides **ASTOPCPE**. Table 9-10 shows the interactions of **ASTOPCPE** with other signals.

The ARM10 processor drives **ASTOPCPE** in ARM10 Execute + 1 stage and the CP Execute stage.

Table 9-10 ASTOPCPE interactions with other signals

Signal	interactions with ASTOPCPD
ASTOPCPD	None.
LSHOLDCPE	ASTOPCPE is asserted with LSHOLDCPE when the pipelines are in lockstep. Pipelines are in lockstep unless the CP has already retired from the ARM10 pipeline and is now transferring data from the LSU for a load/store multiple.
CPBUSYE	The ARM10 processor ignores CPBUSYE if ASTOPCPE is already asserted. ASTOPCPE is not asserted if CPBUSYE was asserted at the end of the previous cycle, but ASTOPCPE can be asserted when CPBUSYE deasserts. In this case, asserting ASTOPCPE continues to hold the same instruction in ARM10 Execute that was held by CPBUSYE .
LSHOLDCPM	ASTOPCPE is asserted with LSHOLDCPM when the pipelines are in lockstep. Pipelines are in lockstep unless the CP has already retired from the ARM10 pipeline and is now transferring data from the LSU for a load/store multiple.
ACANCELCP	ACANCELCP held by ASTOPCPE .
AFLUSHCP	AFLUSHCP overrides ASTOPCPE . The pipeline is flushed from Execute back.
CPBOUNCEE	CPBOUNCEE is not used until ASTOPCPE (and other relevant holds) are deasserted.

9.8.3 ASTOPCPE example

Figure 9-5 on page 9-24 shows the ARM10 processor holding instruction 1 in its Execute stage for one cycle. The numbers in the waveforms show which instruction owns the signal at that time.

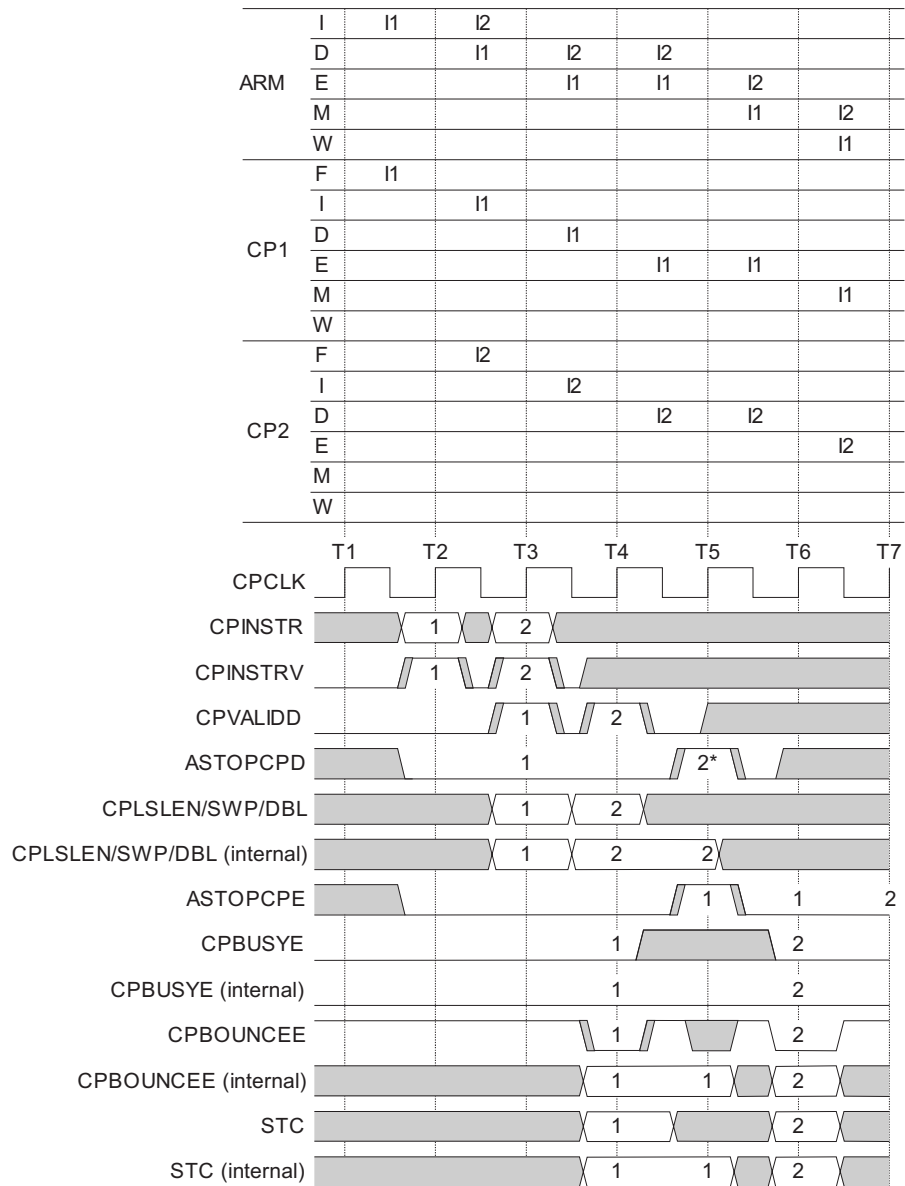


Figure 9-5 ASTOPCPE example

* **ASTOPCPD** is caused by **ASTOPCPE** and **CPBUSYE** is ignored under **ASTOPCPE**. Under an **ASTOPCPE**, **STC** is registered in the ARM10 processor.

9.8.4 LSHOLDCPE

LSHOLDCPE indicates that the load/store CP instruction in the ARM10 LSU Execute stage, did not progress in to the ARM10 LSU Memory stage in the previous cycle. It is driven out of a register following the ARM10 LSU Execute stage. If **LSHOLDCPE** is asserted, CPs must hold their Execute, Decode, Issue, and Fetch stages. If **LSHOLDCPE** is asserted, and a store is in the CP Execute stage, the **STCMRCDATA** bus value is retained by the ARM1010 processor until **LSHOLDCPE** deasserts. Table 9-11 shows the interactions of **LSHOLDCPE** with other signals.

The ARM10 processor drives **LSHOLDCPE** in the ARM10 Execute + 1 stage and the CP Execute stage.

Table 9-11 LSHOLDCPE interactions with other signals

Signal	Interactions with LSHOLDCPE
ASTOPCPD	None.
LSHOLDCPE	None.
ASTOPCPE	LSHOLDCPE is asserted with ASTOPCPE when pipelines are in lockstep. Pipelines are in lockstep unless the CP instruction has already retired from the ALU pipeline and is now transferring data to or from the LSU.
CPBUSYE	CPBUSYE indicates an Execute stage hold when the ALU and LSU pipelines are in lockstep. LSHOLDCPE indicates an LSU execute stage hold when the ALU and LSU pipelines are not in lockstep.
LSHOLDCPM	If LSHOLDCPM is asserted, LSHOLDCPE is asserted as well.
ACANCELCP	None.
AFLUSHCP	Flush invalidates LSHOLDCPE .
CPBOUNCEE	None.

9.8.5 Example of LSHOLDCPE

Figure 9-6 on page 9-26 shows the ARM10 LSU holding instruction 1 in its Execute stage for one cycle. The numbers in the waveforms show which instruction owns the signal at that time. **ASTOPCPD** is caused by **ASTOPCPE**. **CPBUSYE** is ignored under **ASTOPCPE**. Under an **LSHOLDCPE**, **STC** is registered in the ARM10 processor.

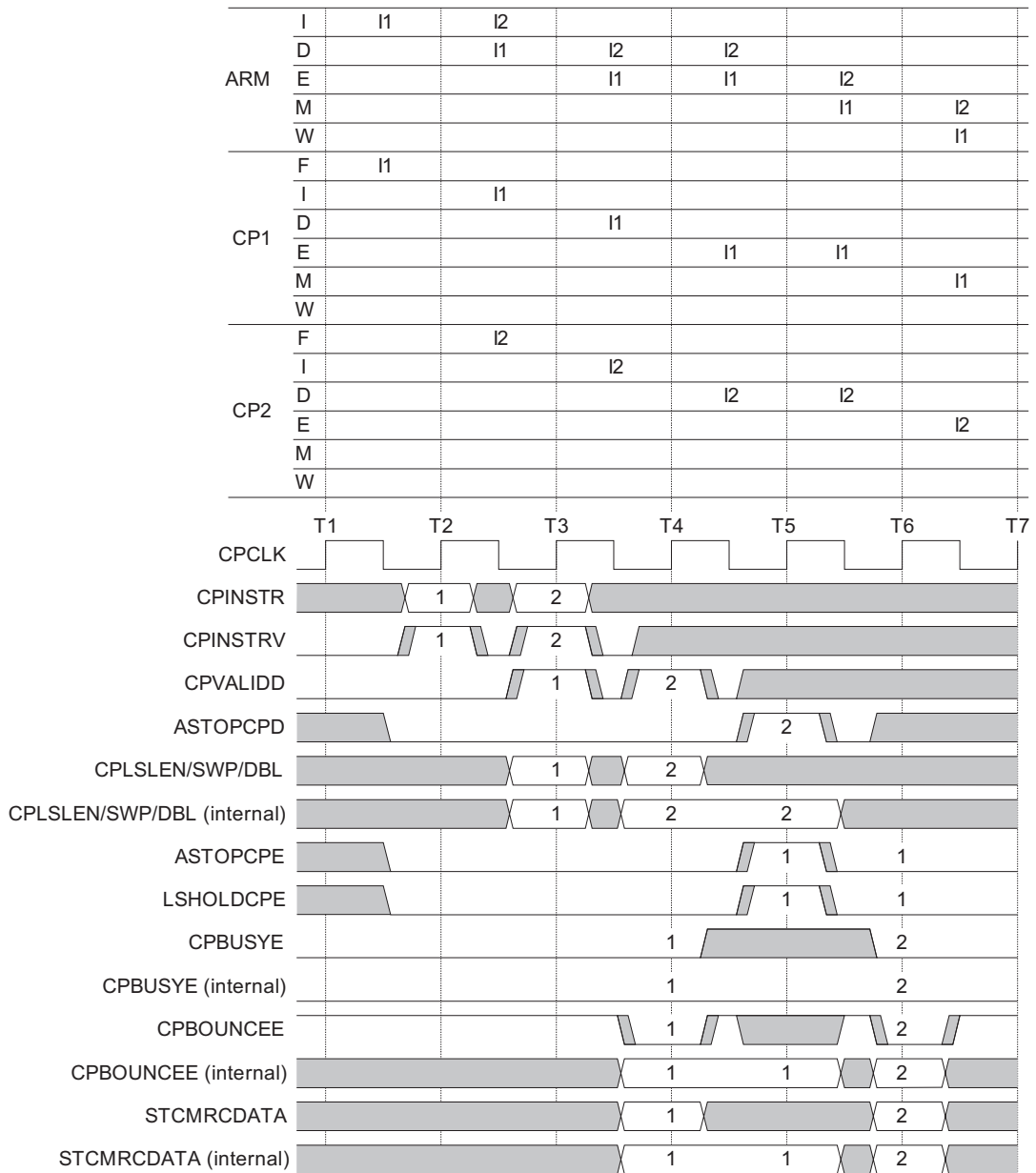


Figure 9-6 LSHOLDCPE example

9.8.6 LSHOLDCPM

LSHOLDCPM indicates that the load CP instruction in the ARM10 LSU Memory stage did not progress into the ARM10 LSU Write stage in the previous cycle or that a load cache miss occurred. It is driven out of a register following the ARM10 LSU Memory stage. If **LSHOLDCPM** is asserted, CPs must hold their Memory, Execute, Decode, Issue and Fetch stages. If **LSHOLDCPM** is asserted, and a load is in the CP Memory stage, the **LDCMCRDATA** bus value is ignored by the CP until **LSHOLDCPM** deasserts. Table 9-12 shows the interactions of **LSHOLDCPM** with other signals.

The ARM10 processor drives **LSHOLDCPM** in the ARM10 Memory + 1 stage and the CP Memory stage.

Table 9-12 LSHOLDCPM interactions with other signals

Signal	Interactions with other signals
ASTOPCPD	None
LSHOLDCPE	None
ASTOPCPE	None
CPBUSYE	None
LSHOLDCPM	If LSHOLDCPM is asserted, LSHOLDCPE is asserted as well.
ACANCELCP	None
AFLUSHCP	None
CPBOUNCEE	None

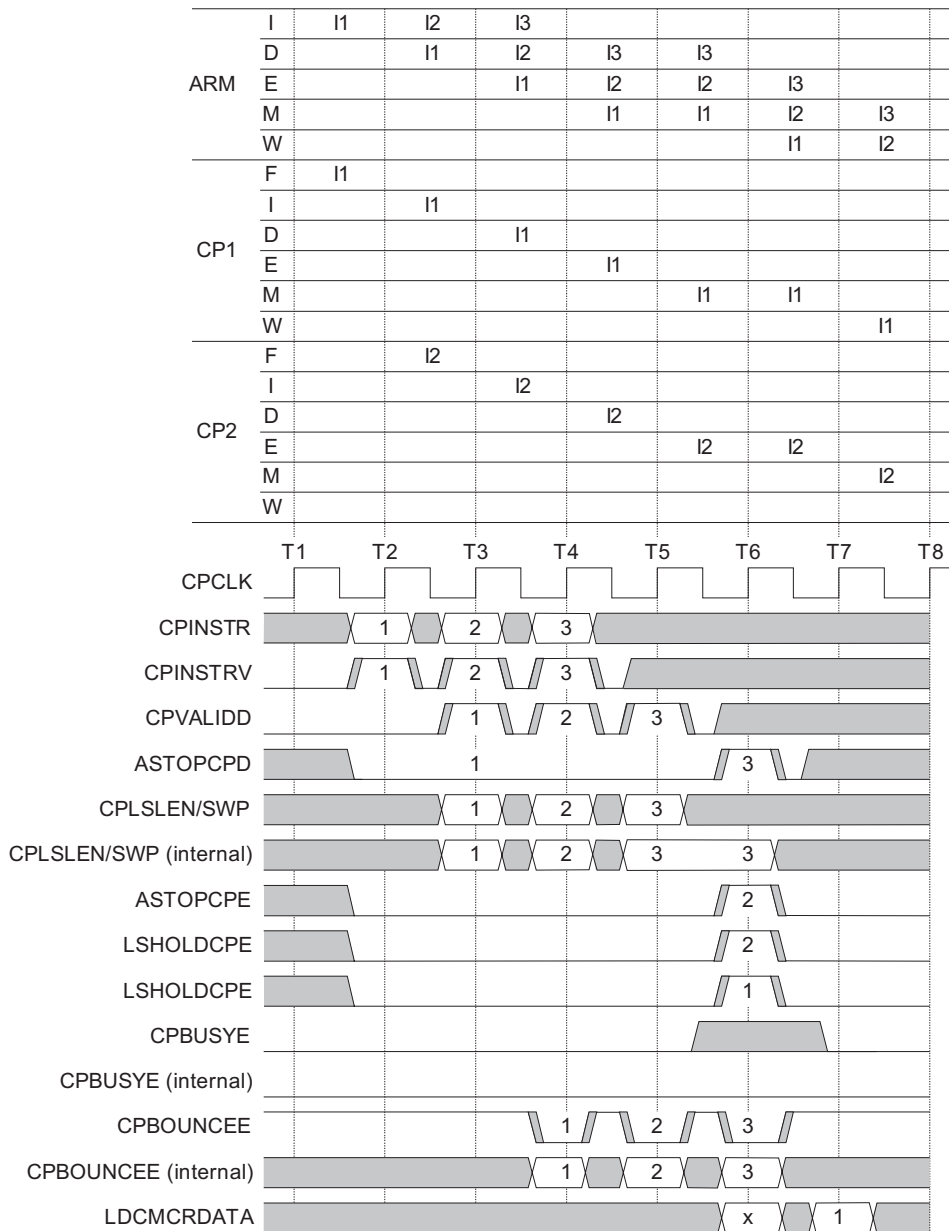


Figure 9-7 LSHOLDCPM example

9.8.7 CPBUSYE

From the ARM10 processor viewpoint, **CPBUSYE** indicates that the CP that owns the instruction in the ARM10 Execute stage wants to hold the instruction in that stage. It is asserted in the ARM10 Execute stage and must come directly out of a register. It also holds the instructions in other CP Issue stages. Table 9-13 shows the interaction of **CPBUSYE** with other signals.

The ARM10 processor drives **CPBUSYE** in the ARM10 Execute stage and the CP Decode stage.

Table 9-13 CPBUSYE interactions with other signals

Signal	interactions with CPBUSYE
ASTOPCPD	The ARM10 processor ignores CPBUSYE if ASTOPCPD is already asserted. ASTOPCPD is not asserted if a valid CPBUSYE (CPBUSY HIGH , ASTOPCPD LOW) was received in the previous cycle.
ASTOPCPE	The ARM10 processor ignores CPBUSYE if ASTOPCPE is already active. ASTOPCPE is not asserted if a valid CPBUSYE was asserted at the end of the previous cycle. ASTOPCPE is not asserted if CPBUSYE is already asserted. ASTOPCPE can be asserted in the cycle that CPBUSYE deasserts.
LSHOLDCPE	None.
LSHOLDCPM	None.
ACANCELCP	None.
AFLUSHCP	AFLUSHCP has priority over CPBUSYE .
CPBOUNCEE	CPBOUNCEE is not used until CPBUSYE (and other holds) are deasserted.

9.8.8 CPBUSYE example

In Figure 9-8 on page 9-30 instruction 1 is held in the ARM10 Execute stage by **CPBUSYE**. Numbers in waveforms show which instruction owns the signal at that time. In some CPs, instruction 1 might advance into Decode under the **CPBUSYE**. In this case instruction 1 spends two cycles in Decode rather than two in Issue. This depends on the CP implementation. For the interface this makes no difference because the interface signals still have to be driven depending upon the position of the instruction in the ARM10 pipeline.

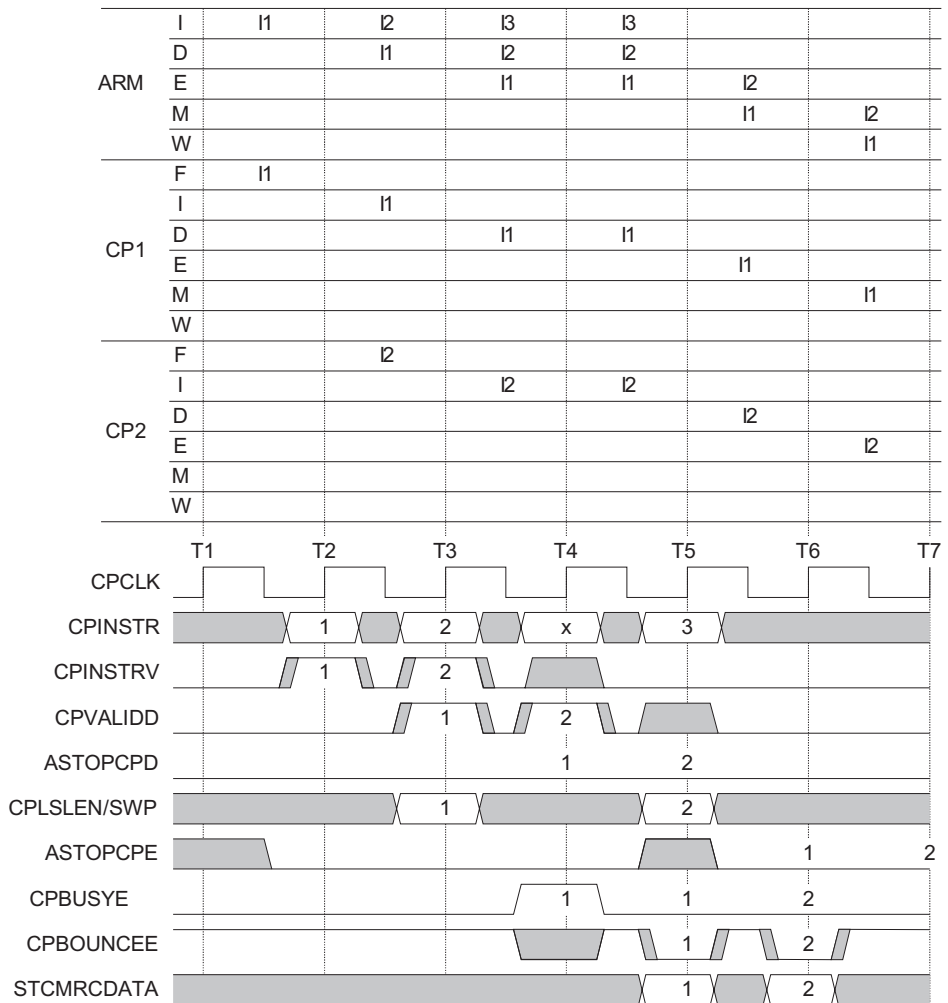


Figure 9-8 CPBUSYE example

9.8.9 CPBUSYE and ASTOPCPD interaction

There is a complex interaction between **ASTOPCPD** and **CPBUSYE**. If **ASTOPCPD** is asserted, the ARM10 processor ignores **CPBUSYE** being asserted in the same cycle, until **ASTOPCPD** deasserts. Figure 9-9 on page 9-31 shows one possible sequence of events.

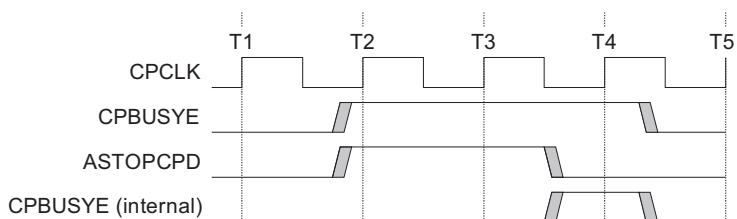


Figure 9-9 CPBUSYE ignored due to ASTOPCPD assertion

If **CPBUSYE** is asserted in the cycle before the ARM10 processor would have asserted **ASTOPCPD**, then **ASTOPCPD** is suppressed until the cycle after **CPBUSYE** deasserts. Figure 9-10 shows this sequence of events.

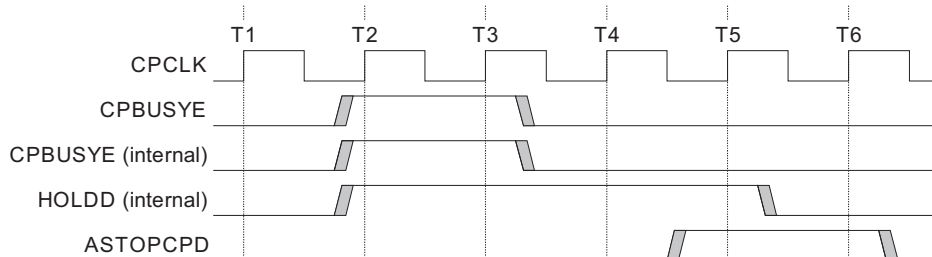


Figure 9-10 CPBUSYE asserted before ASTOPCPD

The internal hold signal **HOLDD** is usually registered to make **ASTOPCPD** in the next cycle, but this is held until **CPBUSYE** goes LOW.

9.8.10 ASTOPCPD with CPBUSYE

In Figure 9-11 on page 9-32, instruction 1 is held up by **CPBUSYE** and instruction 2 is held up by **ASTOPCPD**. An instruction in ARM10 Decode is always held up behind an instruction held by ARM10 **CPBUSYE** in Execute, unless it is flushed.

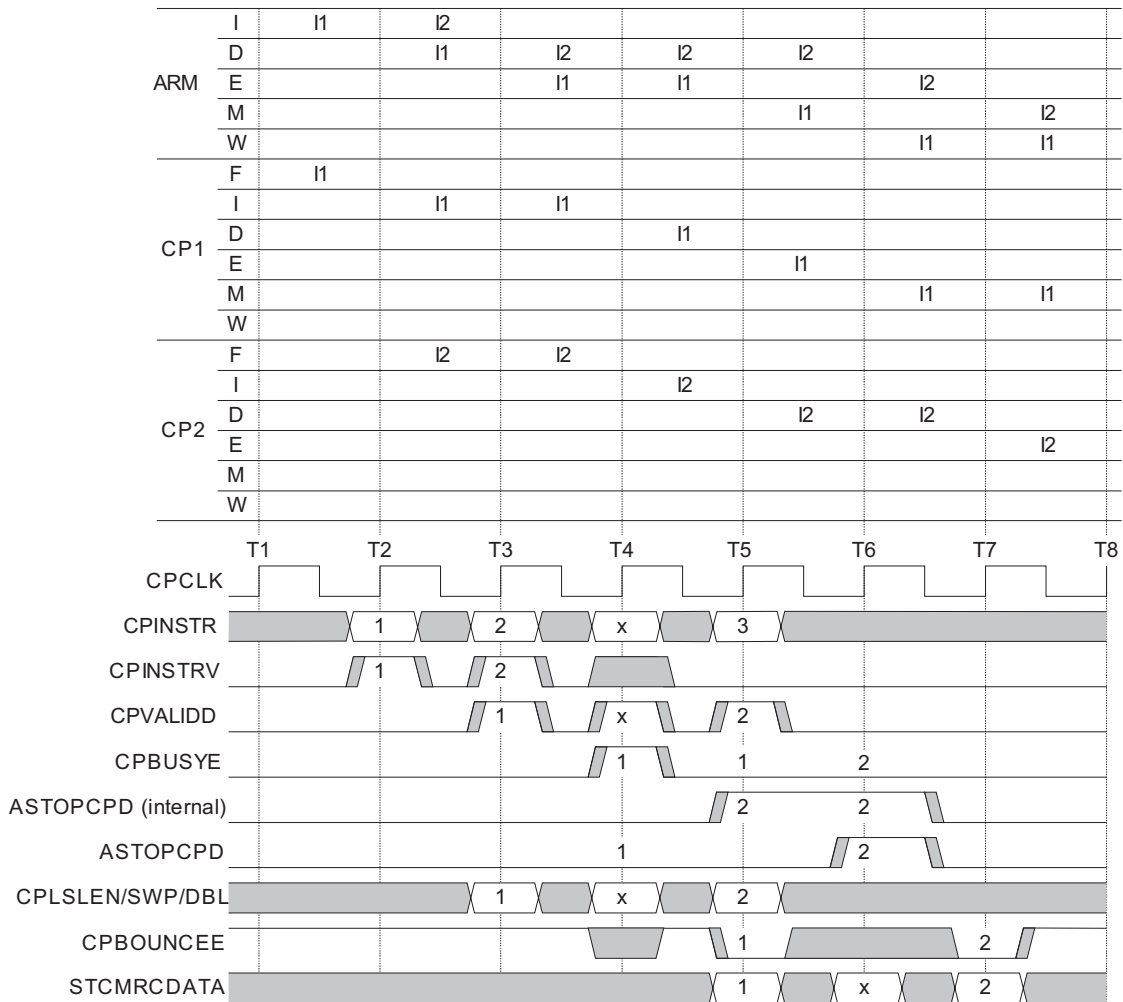


Figure 9-11 ASTOPCPD with CPBUSYE

9.8.11 CPBUSYE and ASTOPCPE interaction

There is a complex interaction between **ASTOPCPE** and **CPBUSYE**. **CPBUSYE** is asserted in the Execute stage of an instruction, **ASTOPCPE** is asserted from a register at the end of the Execute stage (E + 1). If **ASTOPCPE** is asserted in the same cycle that **CPBUSYE** is asserted then **CPBUSYE** is ignored until **ASTOPCPE** deasserts. If **CPBUSYE** was asserted in the previous cycle then **ASTOPCPE** cannot be asserted until the cycle after that in which **CPBUSYE** deasserts.

Where **ASTOPCPE** is asserted at the same time as **CPBUSYE**, the ARM10 processor ignores **CPBUSYE** until **ASTOPCPE** deasserts. In Figure 9-12, **CPBUSYE** is ignored until **ASTOPCPE** deasserts.

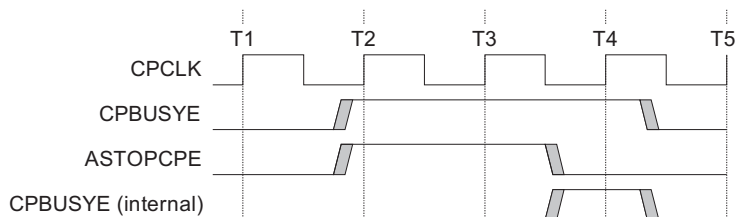


Figure 9-12 CPBUSYE ignored due to ASTOPCPE assertion

In Figure 9-13, **CPBUSYE** is asserted before **ASTOPCPE**. The ARM10 processor does not assert **ASTOPCPE** until the cycle after **CPBUSYE** deasserts. **ASTOPCPE** is holding up the same instruction, in Execute, that **CPBUSYE** held up.

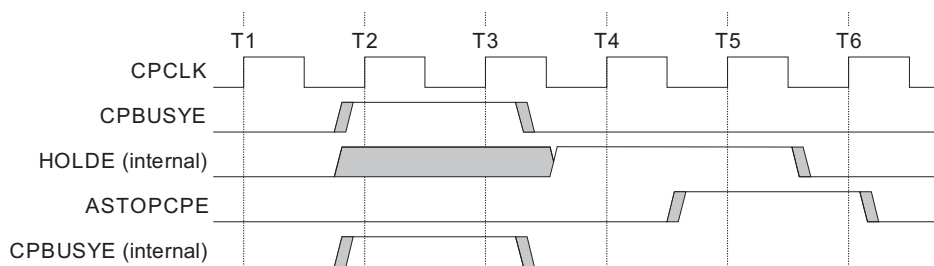


Figure 9-13 CPBUSYE asserted before ASTOPCPE

9.8.12 ASTOPCPE with CPBUSYE

In Figure 9-14 on page 9-34, instruction 2 is held up by **ASTOPCPE** and **CPBUSYE**.

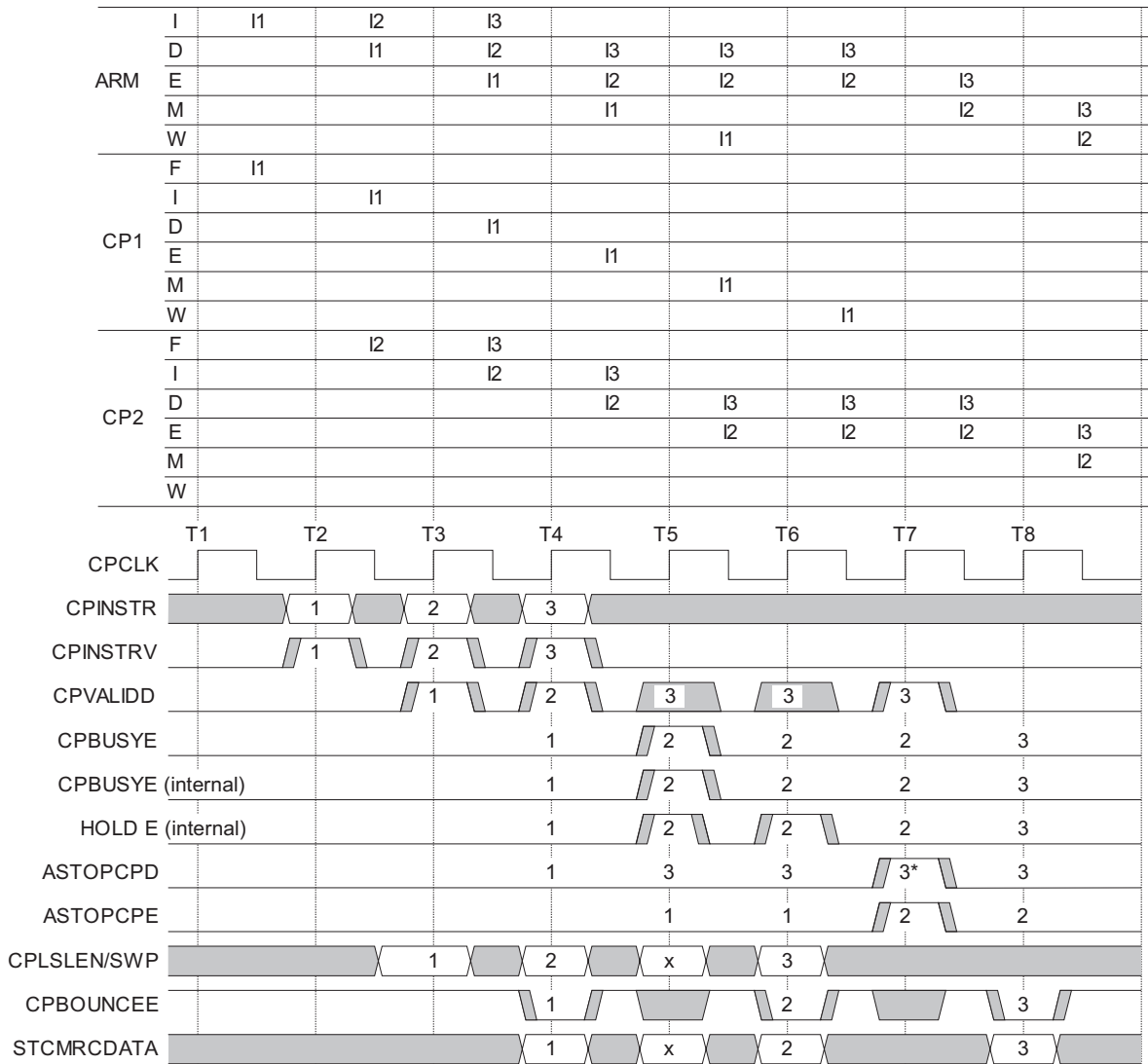


Figure 9-14 I2 held up by ASTOPCPE and CPBUSYE

*Although instruction 3 is responsible for **ASTOPCPD** at T7, instruction 2 has caused **ASTOPCPE** to be asserted and this has to be folded back into **ASTOPCPD**.

In Figure 9-15 on page 9-35, instruction 1 is held up by **ASTOPCPE** and instruction 2 is held up by **CPBUSYE**.

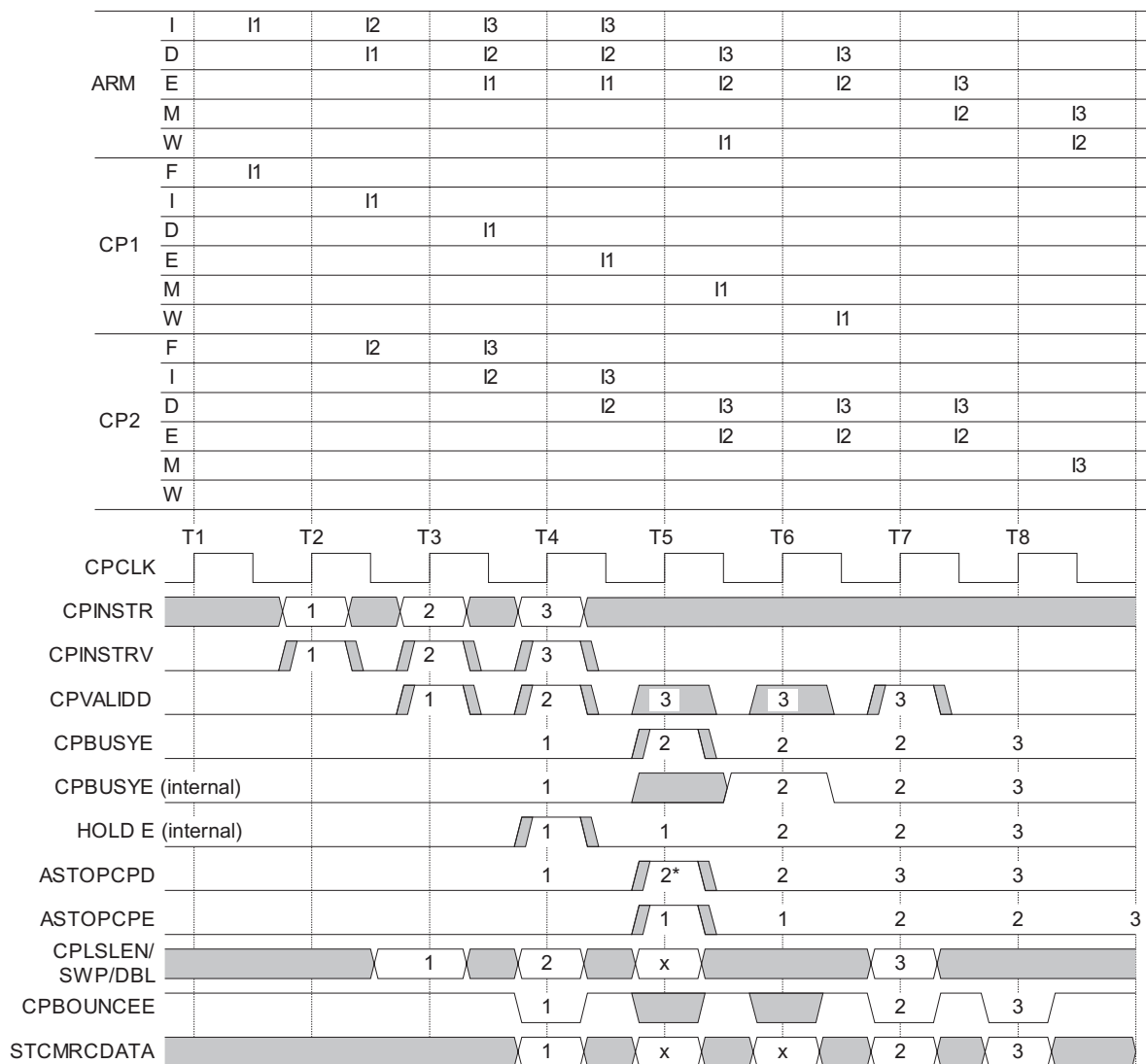


Figure 9-15 I1 held up by ASTOPCPE and I2 held up by CPBUSYE

*Although instruction 2 is responsible for driving **ASTOPCPD** at T5, instruction 1 has caused **ASTOPCPE** to be asserted and this has to be folded back into **ASTOPCPD**.

In Figure 9-16 on page 9-36, instruction 1 is held up by **CPBUSYE** and instruction 2 is held up by **ASTOPCPD**.

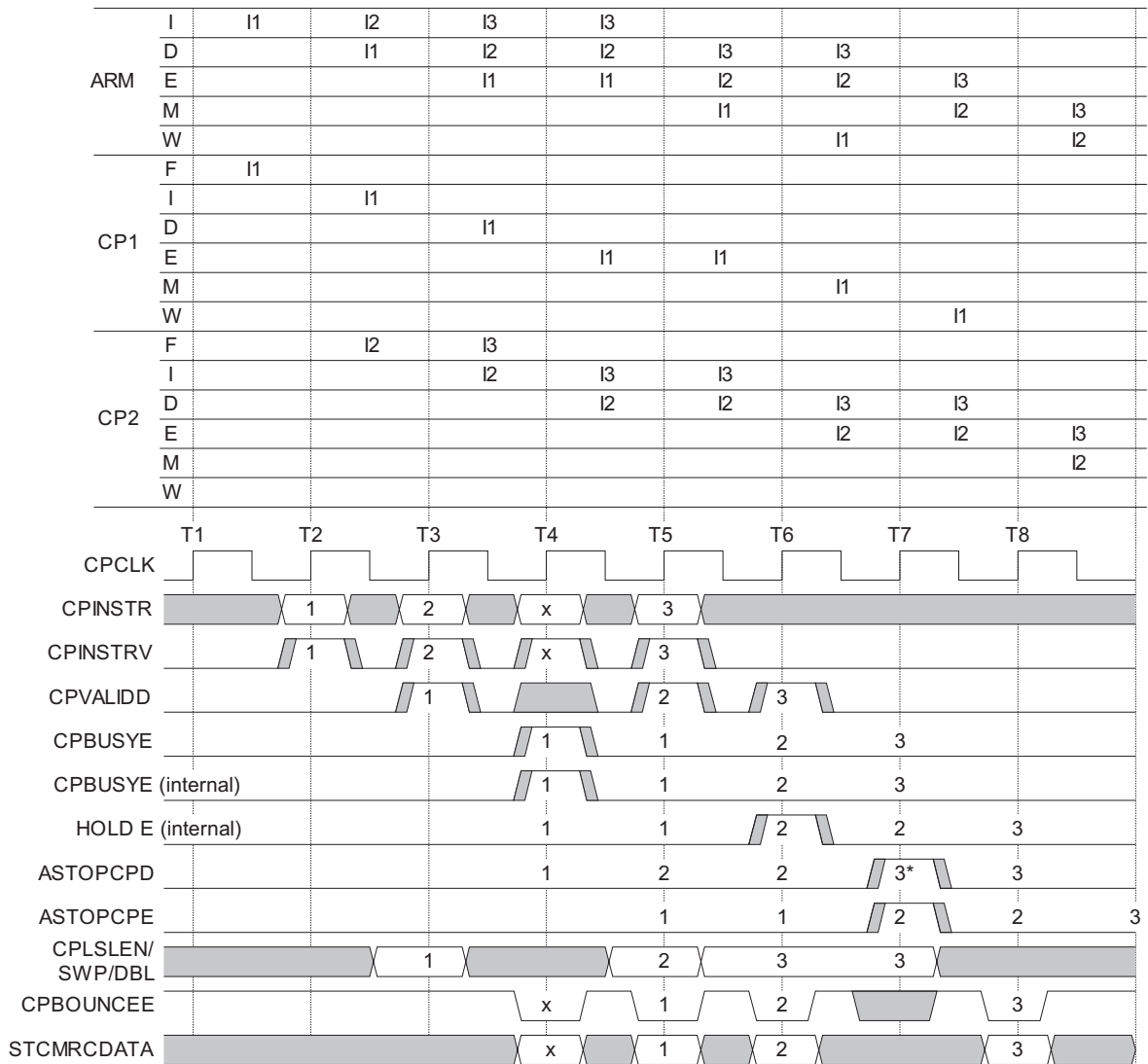


Figure 9-16 I1 held up by CPBUSYE and I2 held up by ASTOPPCPD

*In Figure 9-16 although instruction 3 is responsible for driving **ASTOPCPE** at T7, instruction 2 has caused **ASTOPCPE** to be asserted and this has to be folded back into **ASTOPPCPD**.

9.8.13 CPLSBUSY

This is driven out of a register on the CP Issue/Decode boundary (valid early in the ARM10 Execute stage). It signals to other CPs that the sender is involved in a load or store multiple data transfer and is keeping control of the **STCMRCDATA** bus. Other CPs must progress to Decode (where they are stalled by **ASTOPCPE**) but must not attempt to drive the bus until a cycle after **CPLSBUSY** deasserts.

CPLSBUSY stalls all other CPs when a long LDC is in progress. **CPLSBUSY** does not have to go to the ARM10 processor because it can only do one load/store operation at a time because they are held up in any case. **CPLSBUSY** comes out of flop and goes to other CPs. Table 9-14 shows the interactions of **CPLSBUSY** with other signals.

The CP drives **CPLSBUSY** in the CP Decode stage and the ARM10 Execute stage.

Table 9-14 CPLSBUSY interactions with other signals

Signal	Interactions with CPLSBUSY
ASTOPCPD	None
ASTOPCPE	None
LSHOLDCPE	None
LSHOLDCPM	None
ACANCELCP	None
AFLUSHCP	None
CPBOUNCEE	None

9.9 Instruction cancelation

Instruction cancelation signals are described in the following sections:

- *ACANCELCP*
- *ACANCELCP example*
- *ACANCELCP with ASTOPCPE example* on page 9-40
- *ACANCELCP with CPBUSYE example* on page 9-41
- *AFLUSHCP* on page 9-42
- *AFLUSHCP example* on page 9-42.

9.9.1 ACANCELCP

ACANCELCP indicates that the instruction that has just entered the ARM10 Memory stage must be canceled. **ACANCELCP** differs from **AFLUSHCP**. It cancels a single instruction rather than canceling all upstream instructions in the pipeline. It is driven from register following the ARM10 Execute stage. Table 9-15 shows **ACANCELCP** the interactions with other signals.

The ARM10 processor drives **ACANCELCP** in the ARM10 Memory stage and the CP Execute stage.

Table 9-15 ACANCELCP interactions with other signals

Signal	Interactions with CPBUSYE
ASTOPCPD	None
ASTOPCPE	CP ignores ACANCELCP if ASTOPCPE asserted
LSHOLDCPE	None
CPBUSYE	ACANCELCP is held in response to an active CPBUSYE
LSHOLDCPM	None
ACANCELCP	None
AFLUSHCP	AFLUSHCP has priority
CPBOUNCEE	No effect for canceled instructions

9.9.2 ACANCELCP example

ACANCELCP cancels one instruction (turns it into a NOP) but does not affect the ones around it. In this case, three instructions are issued in a row. Instruction 2 is canceled. Instructions 1 and 3 complete. The numbers in waveforms show which instruction owns

the signal at that time. The ARM10 processor ignores an indication from CP2 that I2 must bounce as the instruction is canceled. Figure 9-17 shows an example with **ACANCELCP**.

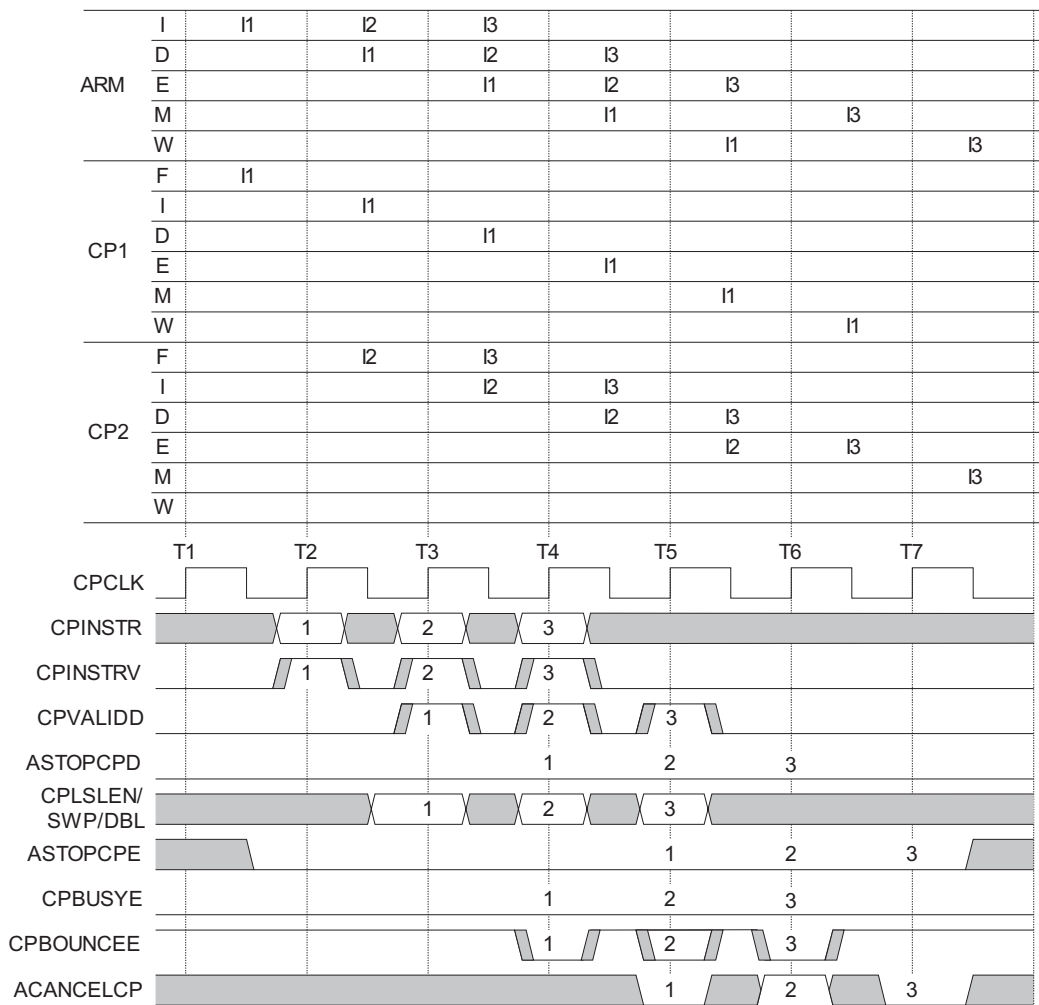


Figure 9-17 ACANCELCP example

The ARM10 processor ignores an indication from CP2 that instruction 2 must bounce because the instruction is canceled.

9.9.3 ACANCELCP with ASTOPCPE example

Instruction 1 is held up by the ARM10 processor with **ASTOPCPE**. **ACANCELCP** is valid in the last cycle that **ASTOPCPE** is asserted. Figure 9-18 shows an example of **ACANCELCP** with **ASTOPCPE**.

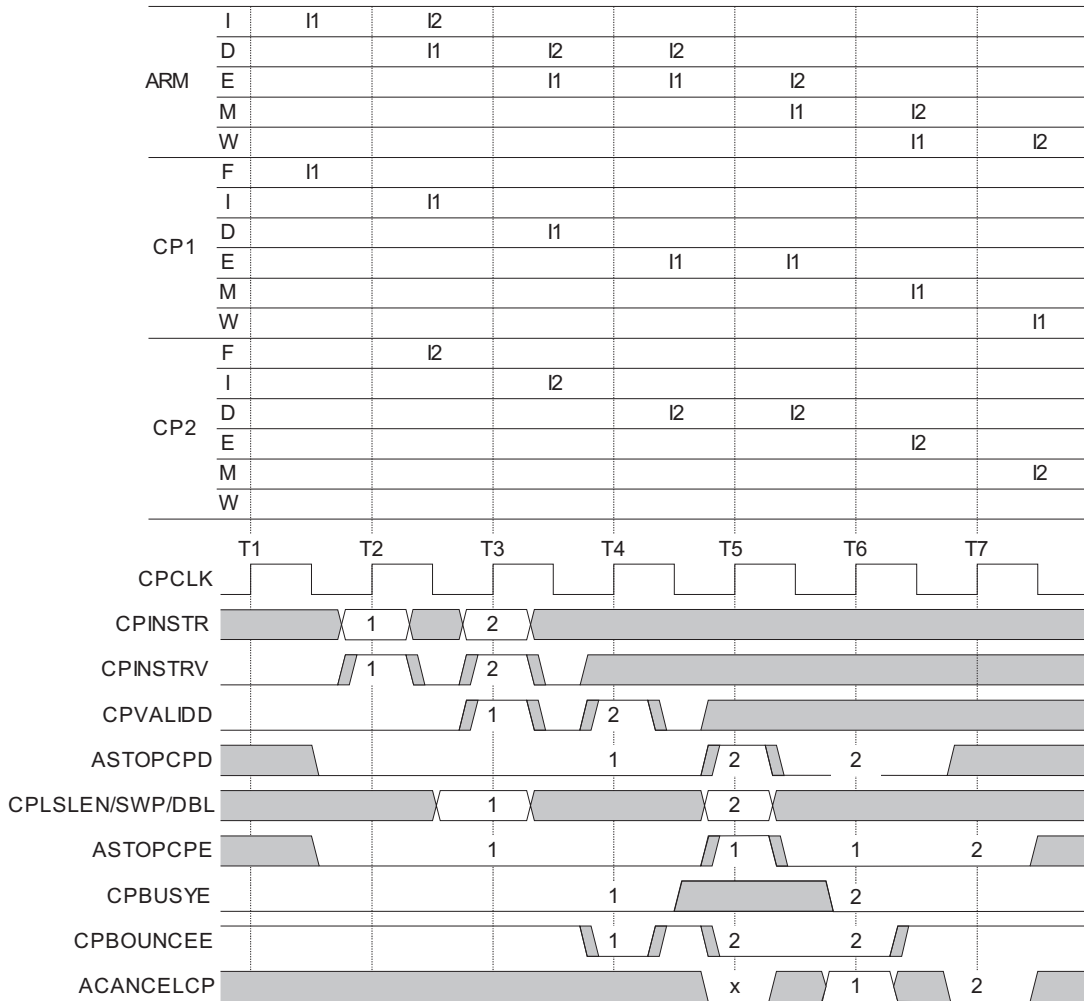


Figure 9-18 ACANCELCP with ASTOPCPE example

9.9.4 ACANCELCP with CPBUSYE example

Instruction 1 is held up by CP1 as indicated by **CPBUSYE**. **ACANCELCP** is valid in the last cycle that **CPBUSYE** is asserted.

ASTOPCPE might be asserted with **CPBUSYE**. It can then be deasserted while **CPBUSYE** is still active or might have stayed asserted when **CPBUSYE** is deasserted. When both **CPBUSYE** and **ASTOPCPE** are deasserted the pipeline must progress. Figure 9-19 shows an example of **ACANCELCP** with **CPBUSYE**.

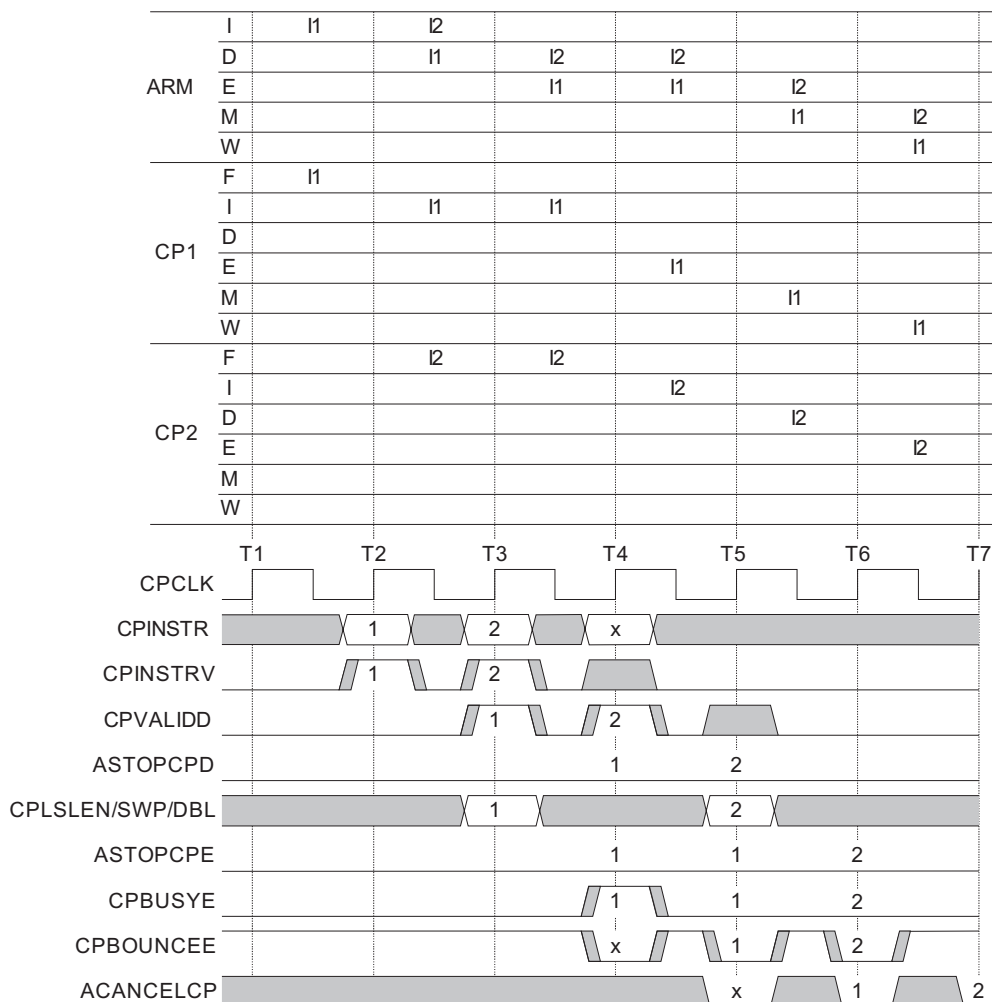


Figure 9-19 ACANCELCP with CPBUSYE example

9.9.5 AFLUSHCP

AFLUSHCP indicates that the instruction that has just entered the ARM10 Memory stage and all upstream instructions currently in the pipeline must be canceled.

AFLUSHCP differs from **ACANCELCP** because it cancels all upstream instructions in the pipeline rather than just a single instruction. It is driven from register following the ARM10 Execute stage. This means that there is no time to factor Data Aborts into the **AFLUSHCP** signal. As a result, aborted CP loads complete when a Data Abort occurs, and then be reexecuted on return from the Data Abort handler routine. It must be possible to execute any CP load more than once (before the next instruction is executed) with no noticeable effects on the CP. Table 9-16 shows the interactions of **AFLUSHCP** with other signals.

The ARM10 processor drives **AFLUSHCP** in the ARM10 Memory stage and the CP Execute stage.

Table 9-16 AFLUSHCP interactions with other signals

Signal	Interactions with CPBUSYE
ASTOPCPD	Flush overrides
ASTOPCPE	Flush overrides
LSHOLDCPE	Flush overrides
CPBUSYE	Flush overrides (deasserted in the following cycle)
LSHOLDCPM	Flush overrides
ACANCELCP	None
CPBOUNCEE	Ignored because instruction canceled by flush

AFLUSHCP supersedes the **ASTOP** and **VALID** signals from the ARM10 processor. It is used to signal that an interrupt has flushed the pipeline. As a result **CPBUSYE** must be deasserted in the following cycle to enable the interrupt to be serviced.

9.9.6 AFLUSHCP example

AFLUSHCP has to override **ASTOPCPE** and **ASTOPCPD**. Here **AFLUSHCP** is asserted for instruction 2. This might be caused by instruction 2 being bounced or a reason unrelated to the CPs, an interrupt, for example. **AFLUSHCP** has to kill the effects of instruction 2 and all following instructions currently in the pipe.

Interrupts can cause flushes at any time. So, even a valid instruction that has been busy-waited for many cycles can be flushed. When the instruction has reached the Memory stage of the ARM10 processor without **AFLUSHCP** or **ACANCELCP** being asserted it completes (with the exception of instructions that Data Abort). Figure 9-20 shows an example of this with five instructions. CP load or store instructions that cause a Data Abort are completed by the CP and rerun by the Data Abort handler. So they must be designed to be rerun with no ill effects. Figure 9-20 shows an example with **AFLUSHCP**.

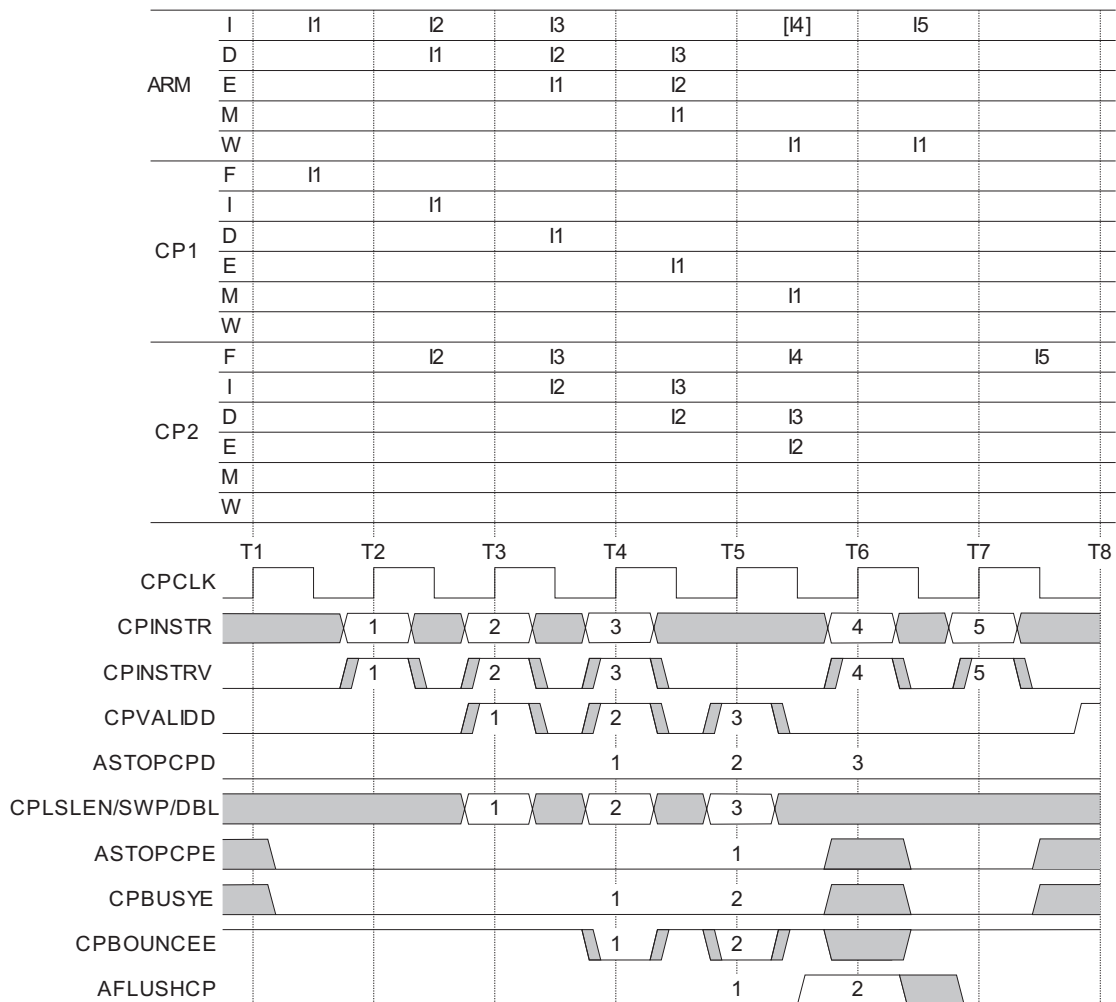


Figure 9-20 AFLUSHCP example

The ARM10 processor ignores an indication from CP2 that I2 might bounce as the instruction is canceled. Instruction 4 might be in the Issue stage. This must be flushed by **AFLUSHCP** but is also not confirmed by **CPVALIDD**. Instruction 5 is issued after the flush and is a valid instruction.

AFLUSHCP can be asserted even if hold signals such as **ACANCELCP** and/or **CPBUSYE** are asserted. In these cases, **AFLUSHCP** has the highest priority because the pipe is currently full of instructions that do not execute. This might be because of a mispredicted branch or an exception.

9.10 Bounced instructions

The following sections describe what happens when CPs cannot execute an instruction, and the undefined instruction trap must be taken:

- *CPBOUNCEE*
- *CPBOUNCEE example* on page 9-46
- *CPBOUNCEE with ASTOPCPE* on page 9-47
- *CPBOUNCEE with CPBUSYE* on page 9-48.

9.10.1 CPBOUNCEE

CPBOUNCEE is used by CPs to acknowledge ownership of CP instructions. Only a CP with an ID that matches the CPID field in the instruction can accept an instruction. If no CP accepts an instruction, the instruction is bounced to an Undefined Instruction handler, and the undefined instruction trap is taken. A CP does not have to accept all instructions with an CPID that matches its ID. This allows a mixture of hardware and software to be used to implement a CP.

The CP drives **CPBOUNCEE** out of a register at the start of the ARM10 Execute stage. When an instruction is bounced, the CP should continue to operate as if it were a NOP. If the bounced instruction passes its condition code check then the ARM10 processor indicates that the CP should flush its pipeline using **AFLUSHCP**.

The CP that owns an instruction on the **CPINSTR** bus drives LOW the **CPBOUNCEE** signal to the ARM10 processor in the CP Decode stage. If the instruction is not owned by a CP, that CP leaves **CPBOUNCEE** HIGH. The ARM10 processor ANDs all individual **CPBOUNCEE** signals internally. If **CPBOUNCEE** is HIGH across ARM10 Execute/Memory boundary, the instruction is deemed to have not been accepted by any CP, and the Undefined instruction trap is taken. A CP might bounce an instruction if the CP is unable to process that instruction or is unable to process a prior instruction and requires software support.

The ARM10 processor ignores **CPBOUNCEE** if **CPBUSYE** is asserted and registers the value of **CPBOUNCEE** at the end of the cycle that **CPBUSYE** deasserts. An active **ASTOPCPE** does not prevent the value of **CPBOUNCEE** from being registered. If a

CP is driving **CPBUSYE**, other CPs must hold **CPBOUNCEE** HIGH. The CP driving **CPBUSYE** must hold its value of **CPBOUNCEE** until the cycle after **CPBUSYE** deasserts. Table 9-17 shows the interactions of **CPBOUNCEE** with other signals.

Table 9-17 CPBOUNCEE interactions with other signals

Signal	Interactions with CPBOUNCEE
ASTOPCPD	None
ASTOPCPE	The ARM10 processor registers CPBOUNCEE even if ASTOPCPE is active
LSHOLDCPE	CPBOUNCEE is ignored until the cycle in which CPBUSYE deasserts
CPBUSYE	Flush overrides
LSHOLDCPM	None
ACANCELCP	A canceled, bounced instruction has no effect
CPBOUNCEE	Ignored as instruction canceled by flush

9.10.2 CPBOUNCEE example

CPBOUNCEE must only be considered valid in the last cycle where neither of **CPBUSYE** or **ASTOPCPE** are asserted. Usually, **AFLUSHCP** is asserted following a **CPBOUNCEE**. One case where this does not happen is when the bounced instruction is canceled at the same time using **ACANCELCP**.

Here instruction 1 completes but instruction 2 bounces and might cause an **AFLUSHCP** that cancels instruction 2 and instruction 3.

As long as one of them is HIGH at all times, **CPBUSYE** and **ASTOPCPE** can be asserted and deasserted under each other multiple times while an instruction is held in Execute. **CPBOUNCEE** is ignored until the first cycle in which both are not asserted. Figure 9-21 on page 9-47 shows an example with **CPBOUNCEE**.

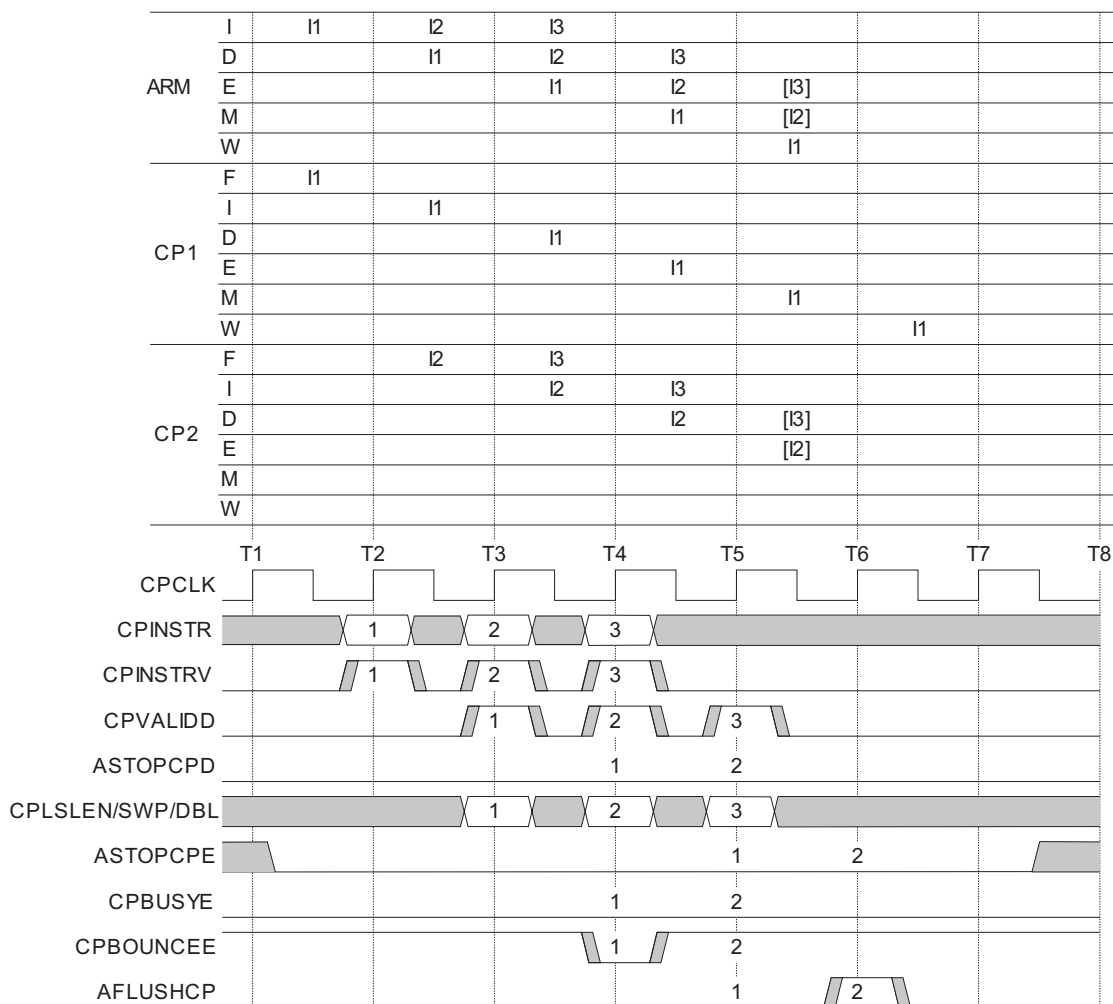


Figure 9-21 CPBOUNCEE example

The flush can occur for a number of reasons. The undefined instruction trap is a low priority exception.

9.10.3 CPBOUNCEE with ASTOPCPE

In Figure 9-22 on page 9-48 instruction 1 is held in the ARM10 Execute stage for one cycle. CPBOUNCEE is only considered valid in the cycle in which ASTOPCPE is deasserted. So, in this case, instruction 1 does not bounce and instruction 2 does.

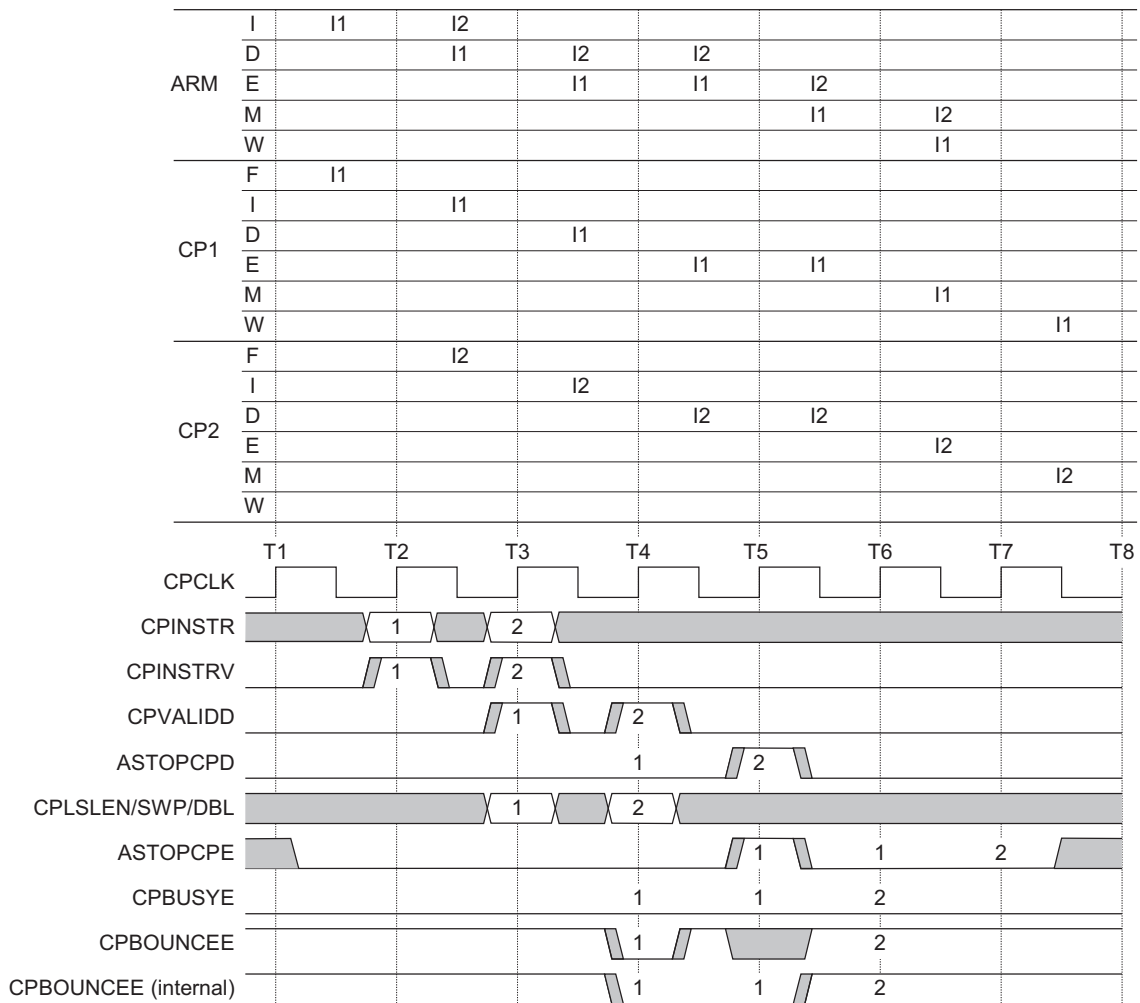


Figure 9-22 CPBOUNCEE with ASTOPCPE example

9.10.4 CPBOUNCEE with CPBUSYE

In Figure 9-23 on page 9-49 Instruction 1 is held in the ARM10 Execute stage for one cycle. **CPBOUNCEE** is only considered valid in the cycle in which **CPBUSYE** is deasserted. So, in this case, instruction 1 does not bounce and instruction 2 does.

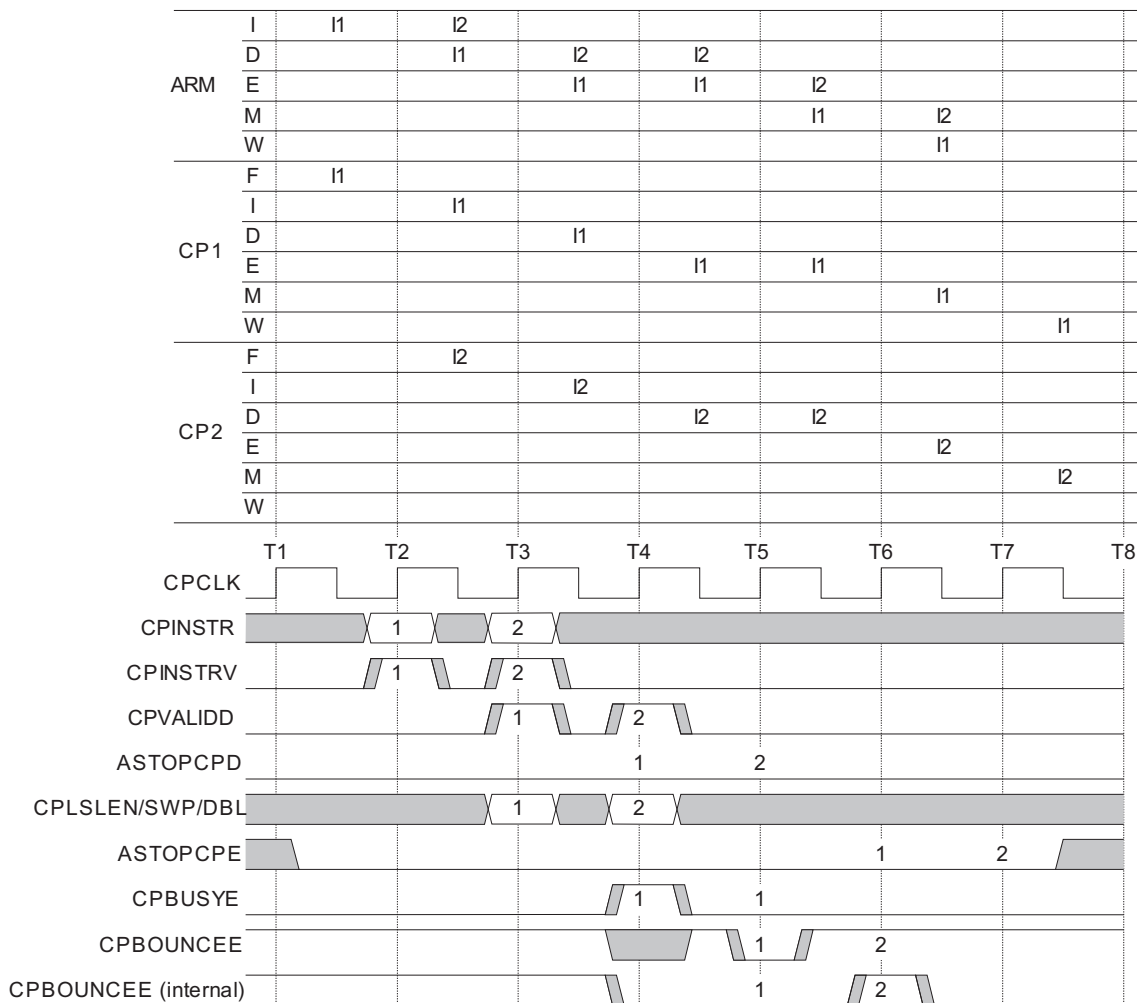


Figure 9-23 CPBOUNCEE with CPBUSYE example

9.11 Data buses

This section describes the 64-bit data buses:

- *STCMRCDATA*
- *LDCMCRDATA*.

9.11.1 STCMRCDATA

The 64-bit **STCMRCDATA** bus carries data from a CP to the ARM10 processor. For a data transfer from a CP register to an ARM10 register (MRC) the data on **STCMRCDATA** is written into a register in the ARM10 register bank. For a CP store to memory (STC), the data on **STCMRCDATA** is passed through ARM10 processor to the memory system. It is stored at an address generated by the ARM10 processor. Table 9-18 describes the interactions between **STCMRCDATA** and signals.

STCMRCDATA is driven by a CP in the ARM10 Execute stage.

Table 9-18 STCMRCDATA interactions with signals

Signal	Interactions with STCMRCDATA
ASTOPCPD	None.
ASTOPCPE	The ARM10 processor registers the value on STCMRCDATA when ASTOPCPE is asserted and the LSU pipeline and ALU pipeline are in lockstep. If the pipelines are decoupled then ASTOPCPE only affects the data processing operation that may be running under the loads or stores.
LSHOLDCPE	If the ALU and LSU pipelines are decoupled then ARM10 processor registers the value on STCMRCDATA when LSHOLDCPE is asserted.
CPBUSYE	None.
LSHOLDCPM	None.
ACANCELCP	None.
CPBOUNCEE	None.

9.11.2 LDCMCRDATA

The 64-bit **LDCMCRDATA** bus carries data from the ARM10 processor to a CP. For a data transfer from an ARM10 register to a CP register (MCR), the data on **LDCMCRDATA** is written into a register in the CP register bank. For a CP load from memory (LDC), the data on **LDCMCRDATA** is passed through the ARM10 processor from the memory system. It is loaded from an address generated by the ARM10 processor. Table 9-19 shows the interactions of **LDCMCRDATA** with other signals.

LDCMRCDATA is driven by the ARM10 processor in the ARM10 Write stage.

Table 9-19 LDCMRCDATA interactions with signals

Signal	Interactions with LDCMRCDATA
ASTOPCPD	None.
ASTOPCPE	None.
LSHOLDCPE	None.
CPBUSYE	None.
LSHOLDCPM	LSHOLDCPM indicates that the memory system did not return valid data in the previous cycle. In this case there is not valid data on LDCMRCDATA until LSHOLDCPM goes LOW.
ACANCELCP	None.
CPBOUNCEE	None.

Chapter 10

JTAG Interface

This chapter describes the JTAG interface built into the ARM10 processor. It contains the following sections:

- *JTAG interface and halt mode* on page 10-2
- *JTAG instructions* on page 10-4
- *Scan chain descriptions* on page 10-8.

10.1 JTAG interface and halt mode

JTAG-based hardware debug using halt mode provides access to the integer unit and debug logic. Access is through scan chains and the IEEE 1149.1 *Test Access Port* (TAP). The TAP state machine is illustrated in Figure 10-1.

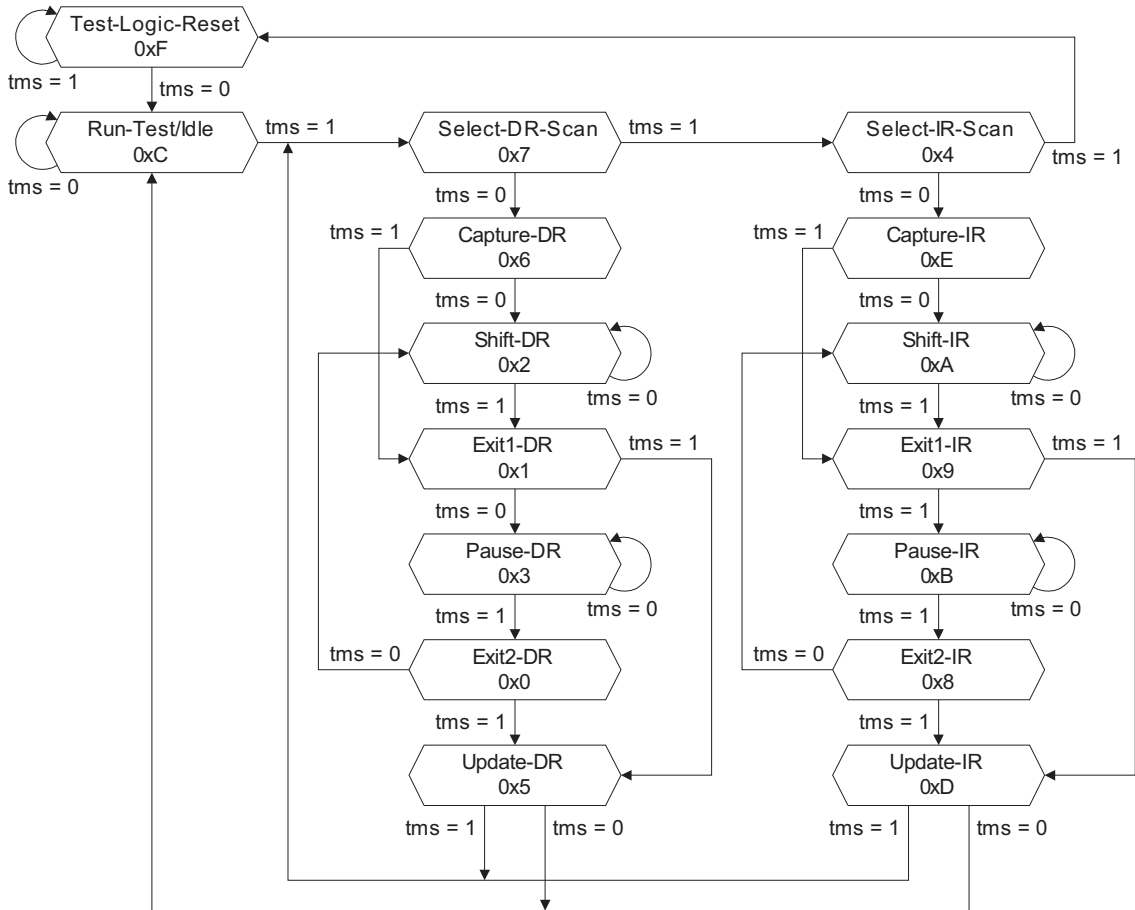


Figure 10-1 JTAG TAP state machine diagram

10.1.1 Entering debug state

Halt mode is enabled by writing a 1 to bit 30 of the *Debug Status and Control Register* (DSCR). This can only be done by external debug hardware such as Multi-ICE. When this mode is enabled, the processor halts, instead of taking an exception in software, if one of the following events occurs:

- A HALT instruction has been scanned in through the JTAG interface. The TAP controller must pass through Run-Test/Idle to issue the HALT command to the ARM10 processor.
- An exception occurs and the corresponding vector catch enable bit is set.
- A register breakpoint hits.
- A watchpoint hits.
- A BKPT instruction reaches the Execute stage of the ARM10 pipeline.
- **EDBGRQ** is asserted.

————— **Note** —————

Software debug must not be used to debug abort and FIQ handlers. Setting a vector trap on FIQ or a watchpoint or breakpoint anywhere in the vector table or handler code for FIQs or aborts can lead to the abort handler being reentered before it has saved state. The value in the abort mode link register and SPSR are overwritten and the information required to return from the handler is lost.

The core halted bit in the DSCR is set when debug state is entered. At this point, the debugger determines why the integer unit was halted and preserves the machine state. The MSR instruction can be used to change modes and gain access to all banked registers in the machine. While in debug state:

- the PC is not incremented
- external interrupts are ignored
- all instructions are read from the instruction transfer register (scan chain 4).

10.1.2 Exiting debug state

To exit from debug state, scan in the RESTART instruction through the JTAG interface. The debugger might adjust the PC before restarting, depending on the way the integer unit entered debug state. When the state machine enters the Run-Test/Idle state, normal operations resume. The delay, waiting until the state machine is in Run-Test/Idle, enables conditions to be set up in other devices in a multiprocessor system without taking immediate effect. When Run-Test/Idle state is entered, all the processors resume operation simultaneously. The core restarted bit is set when the RESTART sequence is complete. The core halted bit **DSCR0** is cleared before the core is restarted.

10.2 JTAG instructions

The JTAG interface portion of the logic implements the IEEE 1149.1 interface and supports:

- a JTAG ID register
- a bypass register
- a 4-bit instruction register.

In addition, the public instructions listed in Table 10-1 are defined.

Table 10-1 Defined public JTAG instructions

Instruction	Binary code	Hexadecimal code
EXTEST	0000	0x0
SCAN_N	0010	0x2
SAMPLE/PRELOAD	0011	0x3
RESTART	0100	0x4
CLAMP	0101	0x5
HIGHZ	0111	0x7
HALT	1000	0x8
CLAMPZ	1001	0x9
INTEST	1100	0xC
IDCODE	1110	0xE
BYPASS	1111	0xF

———— **Note** —————

All unused JTAG instructions default to the BYPASS instruction.

You can access the debug registers through either software, with MCR or MRC instructions, or through the JTAG interface port. See Chapter 11 *Debug* for details of debug registers.

To write the CP14 registers c1, c4, and c5, use the EXTEST instruction. To read CP14 registers c0, c1, and c5, use the INTEST or EXTEXT instruction.

SAMPLE/PRELOAD, CLAMP, HIGHZ, and CLAMPZ are applicable only to external scan chains and they are not supported by scan chains in the ARM10 processor. These instructions are not described in this document.

Note

The CP14 registers do not have interlocks. If the JTAG interface attempts to access a CP14 register while the ARM10 processor is writing to it, the result is Unpredictable.

10.2.1 EXTEST

EXTEST connects the selected scan chain between **TDI** and **TDO**. Loading the instruction register with the EXTEST instruction puts all the scan cells in their test mode of operation.

In the Capture-DR state, inputs to the system logic are captured by the scan cells. In the Shift-DR state, the previously captured test data is shifted out of the scan chain through **TDO**, while new test data is shifted in through the **TDI** input. Data from the boundary scan register cell is applied to the output pins in the Update-DR state. Typically, the first test stimulus to be applied using the EXTEST instruction is shifted into the boundary scan register using the SAMPLE/PRELOAD instruction.

Note

For debug, this instruction connects the selected scan chain between **TDI** and **TDO**. When the instruction register is loaded with the EXTEST instruction, the debug scan chains can be written.

Registers in CP14 that can be written by the JTAG interface, c1, c4, and c5, are written using an EXTEST instruction.

10.2.2 SCAN_N

SCAN_N connects the scan path select register between **TDI** and **TDO**. During the Capture-DR state, the fixed value 1000 is loaded into the register. During the Shift-DR state, the ID number of the desired scan path is shifted into the scan path select register. In the Update-DR state, the scan register of the selected scan chain is connected between **TDI** and **TDO**, and remains connected until a subsequent SCAN_N instruction is issued. On reset, scan chain 3 is selected by default.

10.2.3 RESTART

RESTART is used to restart the processor on exit from debug state. The scan chain path register is not affected and the processor exits debug state once the Run-Test/Idle state is entered.

10.2.4 HALT

HALT stops the integer unit and puts it into debug state. The core can only be put into debug state if debug halt mode is enabled.

10.2.5 INTEST

INTEST connects the selected scan chain between **TDI** and **TDO**. When the instruction register is loaded with the INTEST instruction, all the scan cells are placed in their test mode of operation.

This instruction enables serial testing of on-chip system logic by applying test stimuli. The test results are captured and examined by shifting out the contents of the boundary scan register. In the Capture-DR state, the value of the data applied from the integer unit logic to the output scan cells and the value of the data applied from the system logic to the input scan cells is captured.

In the Shift-DR state, the previously captured test data is shifted out of the scan chain through the **TDO** pin.

Data is typically loaded into the parallel output register stages of the boundary scan chain using the SAMPLE/PRELOAD instruction prior to its use.

———— **Note** —————

This instruction connects the selected scan chain between **TDI** and **TDO**. When the instruction register is loaded with the INTEST instruction, the debug scan chains can be read. INTEST is an optional instruction and its use is governed by the IEEE 1149.1-1990 standard and must be implemented according to those guidelines. In the Capture-DR state, the value of the data applied from the integer unit logic to the output scan cells and the value of the data applied from the system logic to the input scan cells is captured. In the Shift-DR state, the previously captured test data is shifted out of the scan chain through the **TDO** pin, while data shifted in through the **TDI** pin is ignored.

Registers c0, c1, and c5 are read with the INTEST instruction.

10.2.6 IDCODE

IDCODE connects the device identification register, or ID register, between **TDI** and **TDO**. The ID register is a 32-bit register that enables the manufacturer, part number, and version of a component to be determined through the JTAG interface. When the instruction register is loaded with the IDCODE instruction, all the scan cells are placed in their normal (System) mode of operation.

In the Capture-DR state, the device identification code is captured by the ID register. In the Shift-DR state, the previously captured device identification code is shifted out of the ID register through the **TDO** pin, while data is shifted in through the **TDI** pin into the ID register. In the Update-DR state, the ID register is unaffected.

See *TAP ID Register* on page 10-8 for details of selecting and interpreting the ID register value.

10.2.7 BYPASS

BYPASS connects a 1-bit shift register, the bypass register, between **TDI** and **TDO**. When the BYPASS instruction is loaded into the instruction register, all the scan cells are placed in their normal (System) mode of operation. This instruction has no effect on the system pins.

In the Capture-DR state, a logic 0 is captured by the bypass register. In the Shift-DR state, test data is shifted into the bypass register through **TDI** and out through **TDO** after a delay of one **TCK** cycle. The bypass register is not affected in the Update-DR state.

The first bit shifted out is a zero.

All unused JTAG instruction codes default to the BYPASS instruction.

10.3 Scan chain descriptions

This section describes the following scan chains:

- *BYPASS Register*
- *TAP ID Register*
- *Instruction Register* on page 10-9
- *Scan Chain Select Register* on page 10-9
- *Scan chain 0, Debug ID Register* on page 10-10
- *Scan chain 1, Debug Status and Control Register (DSCR)* on page 10-10
- *DSCR readable and writable bits* on page 10-13
- *Scan chain 2* on page 10-14
- *Scan chain 3* on page 10-14
- *Scan chain 4* on page 10-14
- *Scan chain 5, CP14 c5* on page 10-15
- *Scan chain 6* on page 10-15.

10.3.1 BYPASS Register

Purpose	Bypasses the device during scan testing by providing a path between TDI and TDO .
Length	1 bit
Operating mode	When the bypass instruction is the current instruction in the instruction register, serial data is transferred from TDI to TDO in the Shift-DR state with a delay of one TCK cycle. There is no parallel output from the bypass register. A logic 0 is loaded from the parallel input of the bypass register in Capture-DR state.
Order	TDI-[0]-TDO

10.3.2 TAP ID Register

Purpose	The TAP controller ID of each core type is unique. The JTAG ID of this ARM10 processor is initially 0x14A20F0F. A JTAG debugger such as Multi-ICE can easily identify the processor. The JTAG ID register is routed to the edge of the chip so that partners can create their own ID numbers by tying the pins to HIGH or LOW values. Partner-specific devices are identified by ID numbers of the form shown in Figure 10-2 on page 10-9.
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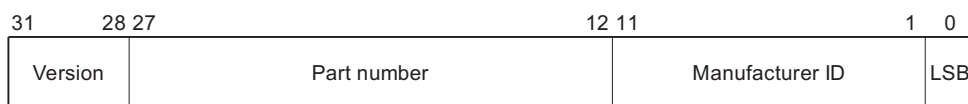


Figure 10-2 TAP ID Register

Length	32 bits
Version	Bits[31:28]
Part Number	Bits [27:12]
Manufacturer ID	Bits [11:1]
LSB	Bit 0
Operating mode	When the IDCODE instruction is current, the TAP ID register is selected as the serial path between TDI and TDO . There is no parallel output from the TAP ID register. The 32-bit ID code is loaded into the register from its parallel inputs during the Capture-DR state.
Order	TDI-[31, 30]...[1, 0]-TDO

10.3.3 Instruction Register

Purpose	Holds the current TAP instruction.
Length	4 bits
Operating mode	When in Capture-DR state, the instruction register is selected as the serial path between TDI and TDO . During the Capture-DR state, the value b0001 is loaded into this register. This is shifted out during Shift-IR, least significant bit first, while a new instruction is shifted in, least significant bit first. During the Update-IR state, the value in the instruction register becomes the current instruction. On reset, IDCODE becomes the current instruction. The value of the current instruction is reflected on the IR[3:0] output bus.
Order	TDI - 3, 2, 1, 0 - TDO

10.3.4 Scan Chain Select Register

Purpose	Holds the current active scan chain.
----------------	--------------------------------------

Length	5 bits
Operating mode	After SCAN_N has been selected as the current instruction, when in Shift-DR state, the scan chain select register is selected as the serial path between TDI and TDO . During the Capture-DR state, binary 10000 is loaded into this register. This is shifted out during Shift-DR, least significant bit first, while a new value is shifted in, least significant bit first. During the Update-DR state, the value in the register selects a scan chain to become the currently active scan chain. All further instructions such as INTEST then apply to that scan chain. The currently selected scan chain only changes when a SCAN_N instruction is executed, or a reset occurs. On reset, scan chain 3 is selected as the active scan chain. The number of the currently selected scan chain is reflected on the SCREG[4:0] output bus. The TAP controller can be used to control external scan chains in addition to those within the ARM10 processor. The external scan chain must be assigned a number and control signals must be generated for it. The number and control signals can be derived from SCREG[4:0] , IR[3:0] , TAPSM[3:0] , and TCK .
Order	TDI - 4, 3, 2, 1, 0 - TDO

10.3.5 Scan chain 0, Debug ID Register

Purpose	Debug. This scan chain is CP14 R0, the debug ID register. The debug ID register value is 0x41006201.
Length	32 bits
Order	TDI - 31, 30, 29, . . . 2, 1, 0 - TDO

10.3.6 Scan chain 1, Debug Status and Control Register (DSCR)

Purpose	Debug. This scan chain is CP14 R1, the DSCR.
Length	32 bits
Defined bits	The following bits are defined for Chain 1.
DSCR0	Core halted.
DSCR1	Core restarted.

DSCR[4:2] Method of debug entry. Table 10-2 shows the method of entry bit values.

Table 10-2 Method of debug entry bit values

DSCR[4:2]	Meaning
000	JTAG HALT instruction occurred
001	Breakpoint occurred
010	Watchpoint occurred
011	BKPT instruction occurred
100	External debug request occurred
101	Vector catch occurred
110	Data-side abort occurred
111	Instruction-side abort occurred

DSCR5 Abort occurred. This is writable only with an MCR to CP14 register c1. This bit is sticky. It is cleared with an MCR to the DSCR where this bit is written as a zero. Reset when **NTRST** = 0 or if the TAP controller is in the reset state.

DSCR6 wDTR buffer empty. This bit indicates to the core that the wDTR buffer is empty, meaning that the core can write more data into it. This is the inversion of the bit that the JTAG debugger sees if it polls the DTR by going through Capture-DR with EXTEST. The debugger must not use this bit to determine if the wDTR is empty or full because the timing between the JTAG interface signal and the core signal is different.

DSCR7 rDTR buffer full. This bit indicates to the core that the rDTR buffer is full, meaning that the debugger has written data into it. This is the inversion of the bit that the JTAG debugger sees if it polls the DTR by going through CaptureDR with INTEST. The debugger must not use this bit to determine if the rDTR is empty or full because the timing between the JTAG interface signal and the core signal is different.

DSCR[15:8] Reserved.

DSCR16	Vector catch enable, Reset. Reset when NTRST = 0 or if the TAP controller is in the Reset state.
DSCR17	Vector catch enable, undefined instruction. Reset when NTRST = 0 or if the TAP controller is in the Reset state.
DSCR18	Vector catch enable, SWI. Reset when NTRST = 0 or if the TAP controller is in the Reset state.
DSCR19	Vector catch enable, Prefetch Abort. Reset when NTRST = 0 or if the TAP controller is in the Reset state.
DSCR20	Vector catch enable, Data Abort. Reset when NTRST = 0 or if the TAP controller is in the Reset state.
DSCR21	Reserved.
DSCR22	Vector catch enable, IRQ. Reset when NTRST = 0 or if the TAP controller is in the Reset state.
DSCR23	Vector catch enable, FIQ. Reset when NTRST = 0 or if the TAP controller is in the Reset state.
DSCR[26:24]	Reserved.
DSCR27	Comms Channel Mode: 1 = comms channel activity 0 = no comms channel activity. Reset when NTRST = 0 or if the TAP controller is in the Reset state.
DSCR28	Thumb state indicator (see Table 11-5 on page 11-7). Thumb instruction: 1 = ITR contains a Thumb instruction 0 = ITR contains an ARM instruction.

- DSCR29** Execute instruction in ITR select:
 1 = instruction in ITR is sent to prefetch unit if JTAG state machine passes through Run-Test/Idle
 0 = disabled.
 Set when **NTRST** = 0 or if the TAP controller is in the Reset state.
- DSCR30** Halt/Monitor mode select:
 1 = halt mode enabled
 0 = monitor mode enabled.
 Reset when **NTRST** = 0 or if TAP controller is in Reset state.
- DSCR31** Global debug enable:
 1 = all debugging functions enabled
 0 = all debugging functions disabled (breakpoints and watchpoints).
 Reset when **NRESET** = 0 (the core Reset line).

10.3.7 DSCR readable and writable bits

The DSCR can be seen from core and from the JTAG interface. The readable and writable bits seen from the core and the JTAG debugger are summarized in Table 10-3.

Table 10-3 DSCR bits from the core

DSCR bits	View from core	View from JTAG
[1:0]	Reserved	Read-only
[4:2]	Read-only	Read-only
[5]	Reserved	Read-only
[7:6]	Read-only	Read-only
[15:8]	Reserved	Reserved
[23:16]	Read-only	Readable/writable
[26:24]	Reserved	Reserved
[30:27]	Reserved	Readable/writable
[31]	Readable/writable	Read-only

———— **Note** ————

The comms channel bits, rDTRFull and wDTREmpty, are inversions of what the debugger sees, because these bits are mirrored in the DSCR for the core, not the debugger.

Order TDI - 31, 30 . . . 1, 0 - TDO

10.3.8 Scan chain 2

Scan chain 2 is the combination of scan chain 4 and scan chain 5. Scan chain 4 is the *Instruction Transfer Register*, ITR, and scan chain 5 is the *Data Transfer Register*, DTR. The instruction complete bit, ITR0, is not included in this combination. It appears only in scan chain 4. Scan chain 2 is illustrated in Figure 10-3.

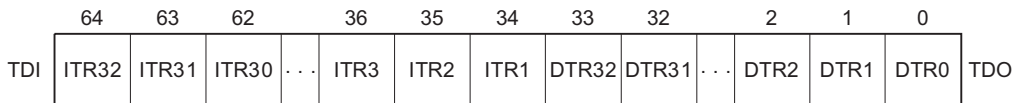


Figure 10-3 Scan chain 2

10.3.9 Scan chain 3

Purpose Can be used for external boundary scan testing. Used for interdevice testing (EXTEST) and testing the core (INTEST).

Length Undetermined

10.3.10 Scan chain 4

Purpose Debug

Length 33 bits

Purpose This scan chain is the *Instruction Transfer Register* (ITR), used to send instructions to the core through the prefetch unit. This chain consists of 32 bits of information, plus an additional bit to indicate the completion of the instruction sent to the core. Instructions scanned into the ITR are not executed unless the instruction transfer execute bit DSCR29 is asserted. Bit 0 indicates if the instruction in the ITR has completed execution.

Order TDI-[32, 31, 30]...[1, 0]-TDO

10.3.11 Scan chain 5, CP14 c5

CP14 is the *Data Transfer Register*, DTR. It consists of two separate registers, the read-only rDTR and the write-only wDTR. The two registers facilitate the creation of a bidirectional comms channel in software.

The rDTR can be loaded only through the JTAG port and is read-only by the core using an MRC instruction. The rDTR chain contains 32 bits of information plus one additional bit for the comms channel.

The wDTR can be loaded only by the core through an MCR instruction and is read-only through the JTAG port. The wDTR contains 32 bits of information plus one additional bit for the comms channel. The definition of bit 0 depends on whether the current JTAG instruction is INTEST or EXTEST. If the current instruction is EXTEST, the debugger can write to the rDTR, and bit 0 indicates if there is still valid data in the queue. If the bit is set, the debugger can write new data. When the core performs a read of the rDTR, bit 0 is automatically asserted. Conversely, if the JTAG instruction is INTEST, bit 0 indicates if there is currently valid data to read in the wDTR. If the bit is set, the JTAG interface must read the contents of the wDTR, which in turn, clears the bit. The core can then sample its own wDTR empty bit and write new data for the debugger.

The TAP controller see either rDTR or wDTR depending on the instruction only sees one register through scan chain 5, and the appropriate register is chosen depending on which instruction is used (INTEST or EXTEST).

Purpose	Debug
Length	33 bits
Order	TDI-rDTR[32]rDTR[31]...rDTR[1]rDTR[0] wDTR[32]wDTR[31]...wDTR[1]wDTR[0]-TDO

10.3.12 Scan chain 6

Purpose	ETM
Length	40 bits
Purpose	The ETM scan chain. Refer to <i>Embedded Trace Macrocell Technical Reference Manual</i> .

Chapter 11

Debug

This chapter describes the debug unit. These features assist the development of application software, operating systems, and hardware. This chapter contains the following sections:

- *About the debug unit* on page 11-2
- *Register descriptions* on page 11-5
- *Software lockout function* on page 11-15
- *Halt mode* on page 11-16
- *Monitor mode* on page 11-19
- *Values in the Link Register after aborts* on page 11-20
- *Comms channel* on page 11-21.

11.1 About the debug unit

The debug unit assists in debugging software. The debug hardware, in combination with a software debugger program, can be used to debug:

- application software
- operating systems
- ARM10-based hardware systems.

The debug unit enables you to:

- stop program execution
- examine and alter processor and coprocessor state
- examine and alter memory and input/output peripheral state
- restart the processor core.

The debug unit provides several ways to stop execution. The most common is for execution to halt when a particular memory address is accessed, either for an instruction fetch (a breakpoint), or a data access (a watchpoint). When execution has stopped, one of two modes is entered:

Halt mode All processor execution halts, and can only be restarted with hardware connected to the external JTAG interface. You can examine and alter all processor state (CPU registers), coprocessor state, memory, and input/output locations through the JTAG interface. This mode is intentionally invasive to program execution. In halt mode you can debug the processor irrespective of its internal state. Halt mode requires external hardware to control the JTAG interface. A software debugger provides the user interface to the debug hardware.

Monitor mode In monitor mode the processor stops execution of the current program and starts execution of a Debug Abort handler. The state of the processor is preserved in the same manner as all ARM exceptions (see *The ARM Architecture Reference Manual* on exceptions and exception priorities). The abort handler communicates with a debugger application to access processor and coprocessor state, and to access memory contents and input/output peripherals. Monitor mode requires a debug monitor program to interface between the debug hardware and the software debugger.

11.1.1 Halt mode and monitor mode compared

Halt mode is for nonreal-time debugging. Because of its hardware nature, you can use halt mode to debug the processor under almost all circumstances. However, real-time systems in which processor execution cannot be completely suspended are unlikely to be able to tolerate the intrusion caused by halt mode. Therefore monitor mode is provided for time-critical applications that cannot tolerate a long interruption while the processor is halted. Monitor mode relies on the processor being able to freely execute instructions to process debug requests.

11.1.2 Programming the debug unit

The debug unit is programmed using Coprocessor 14, CP14. CP14 provides:

- instruction address comparators for triggering breakpoints
- data address comparators for triggering watchpoints
- a bidirectional serial communication channel
- all other state information associated with debug.

CP14 is accessed using coprocessor instructions in both halt mode and monitor mode. BKPT instructions cause a Prefetch Abort if debug is disabled.

11.1.3 Summary of CP14 registers

All debug state is mapped into CP14 as registers. Three CP14 registers, c0, c1, and c5, can be accessed by software running on the processor. Four registers, c0, c1, c4, and c5, are accessible as scan chains from the JTAG interface. Register c4, the instruction transfer register, is accessible only as a scan chain. The remaining registers are accessible only by software operating in a privileged processor mode. Table 11-1 shows the CP14 registers and their scan chain numbers.

Table 11-1 CP14 registers and scan chain numbers

Register	Register name	Scan chain number	Software access mode	
			User	Privileged
CP14 c0	<i>Debug ID register, DIDR</i>	0	Yes	Yes
CP14 c1	<i>Debug Status and Control Register, DSCR</i>	1	Yes	Yes
CP14 c2 and c3	Reserved	-	No	No
CP14 c4	<i>Instruction Transfer Register, ITR</i>	4	No	No

Table 11-1 CP14 registers and scan chain numbers (continued)

Register	Register name	Scan chain number	Software access mode	
			User	Privileged
CP14 c5	<i>Data Transfer Register, DTR</i>	5	Yes	Yes
CP14 c6-c63	Reserved	-	No	No
CP14 c64-c69	<i>Breakpoint Address registers, BA0-BA5</i>	-	No	Yes
CP14 c70-c79	Reserved	-	No	No
CP14 c80-c85	<i>Breakpoint Control registers, BC0-BC5</i>	-	No	Yes
CP14 c86-c95	Reserved	-	No	No
CP14 c96 and c97	<i>Watchpoint Address registers, WA0 and WA1</i>	-	No	Yes
CP14 c112 and c113	<i>Watchpoint Control registers, WC0 and WC1</i>	-	No	Yes
CP14 c114 and c127	Reserved	-	No	No

The register file has space reserved for up to 16 breakpoints and 16 watchpoints. A particular implementation can have any number from 2 to 16. The processor has six instruction-side breakpoints and two data-side watchpoints.

There are two requirements to enable debugging:

- An enable bit in the debug status and control register enables debug functionality through software. Reset clears the enable bit, disabling all debug functionality. The processor ignores external debug requests, and BKPT instructions cause Prefetch Aborts. In this mode, an operating system can quickly enable and disable debugging on individual tasks as part of the task-switching sequence.
- The **DBGEN** pin allows the debug features of the processor to be disabled entirely.

The **DBGEN** pin must be tied HIGH to enable the debug functionality of the core. **DBGEN** must be tied LOW only when debugging is never required.

The CRm and opcode2 fields are used to encode the debug register number, where the register number is {opcode2, CRm}.

11.2 Register descriptions

This section describes the CP14 registers:

- *CP14 c0, Debug ID Register*
- *CP14 c1, Debug Status and Control Register* on page 11-6
- *CP14 c2, c3* on page 11-8
- *CP14 c4, Instruction Transfer Register* on page 11-8
- *CP14 c5, Data Transfer Register* on page 11-8
- *CP14 c6-c63* on page 11-9
- *CP14 c64-c69, Breakpoint Address Registers* on page 11-9
- *CP14 c70-c79* on page 11-10
- *CP14 c80-c85, Breakpoint Control Registers* on page 11-10
- *CP14 c86-c95* on page 11-12
- *CP14 c96 and c97, Watchpoint Address Registers* on page 11-12
- *CP14 c112 and c113, Watchpoint Control Registers* on page 11-13
- *CP14 c114-c127* on page 11-14

11.2.1 CP14 c0, Debug ID Register

The Debug ID Register, DIDR, is a read-only register and contains 0x41006201. It is software accessible in both user and privileged modes and also as a scan chain from the JTAG interface. Table 11-2 shows the instructions for reading DIDR.

Table 11-2 Debug ID Register instructions

Instruction	Description
MRC p14,0,Rd,c0,c0,0	Copies contents of debug ID register into Rd

Figure 11-1 shows the DIDR bit fields.

31	24 23	20 19	16 15	12 11	8 7	4 3	0
Designer code 0100 0001	SBZ 0000	Architecture 0000	Breakpoints 0110	Watchpoints 0010	SBZ 0000	Revision 0001	

Figure 11-1 Debug ID Register

Table 11-3 describes the DIDR bit fields.

Table 11-3 Encoding of the Debug ID Register

Bits	Meaning
[31:24]	Designer code
[23:20]	Should Be Zero
[19:16]	Debug architecture version
[15:12]	Number of implemented register breakpoints
[11:8]	Number of implemented watchpoints
[7:4]	Should Be Zero
[3:0]	Revision number

11.2.2 CP14 c1, Debug Status and Control Register

The Debug Status and Control Register, DSCR, is a read/write register. It is software accessible in both user and privileged modes and also as a scan chain from the JTAG interface.

Table 11-4 shows the instructions for accessing DSCR.

Table 11-4 Debug Status and Control Register instructions

Instruction	Description
MRC p14, 0, Rd, c0, c1, 0	Copies contents of debug status and control register into Rd
MCR p14, 0, Rd, c0, c1, 0	Copies contents of Rd into debug status and control register

Figure 11-2 shows the DSCR bit fields.

31	30	29	28	27	24	23	22	21	20	19	18	17	16	15	8	7	6	5	4	2	1	0
GE	H	E	T	Reserved	CF	C1	R	CD	CP	CS	CU	CR	Reserved	RF	WE	R	MOE	Reserved				

Figure 11-2 Debug Status and Control Register

Table 11-5 describes the DSCR bit fields.

Table 11-5 Encoding of Debug Status and Control Register

Bits	Definition
[31]	GE, global debug enable bit. Reset clears GE. 1 = All debugging functions enabled 0 = All debugging functions disabled.
[30]	H, halt mode bit. Reset clears H. 1 = halt mode 0 = monitor mode.
[29]	E, execute bit. 1 = execute instruction in ITR when in JTAG Run-Test/Idle state 0 = do not execute instruction in ITR when in JTAG Run-Test/Idle state.
[28]	T, Thumb instruction bit: 1 = ITR contains a Thumb instruction 0 = ITR contains an ARM instruction.
[27:24]	Reserved.
DSCR[23:22] and DSCR[20:16] are used to catch ARM exceptions. The effect of setting one of these bits is the same as setting a register breakpoint on the address of the exception vector.	
[23]	CF, vector catch FIQ bit, read-only.
[22]	CI, vector catch IRQ bit, read-only.
[21]	Reserved.
[20]	CD, vector catch Data Abort bit, read-only.
[19]	CP, vector catch Prefetch Abort bit, read-only.
[18]	CS, vector catch Software Interrupt bit, read-only.
[17]	CU, vector catch Undefined Instruction bit, read-only.
[16]	CR, vector catch reset bit, read-only.
[15:8]	Reserved.
[7]	RF, rDTR buffer full bit, read-only: 1 = new data placed in the rDTR through the JTAG interface that can be read with a MRC or STC instruction 0 = no new data placed in the rDTR through the JTAG interface.

Table 11-5 Encoding of Debug Status and Control Register (continued)

Bits	Definition
[6]	WE, wDTR buffer empty bit, read-only: 1 = the wDTR buffer is ready to have data written to it 0 = data has not been read through the JTAG interface.
[5]	-
[4:2]	MOE, method of entry bits, read-only: 000 = JTAG halt instruction 001 = breakpoint hit 010 = watchpoint hit 011 = breakpoint instruction requested 100 = external debug requested asserted 101 = vector catch occurred 110 = data-side abort occurred 111 = instruction-side abort occurred.
[1:0]	-

11.2.3 CP14 c2, c3

CP14 registers c2 and c3 are reserved.

11.2.4 CP14 c4, Instruction Transfer Register

The Instruction Transfer Register, ITR, is a read/write register and is accessible only as a scan chain from the JTAG interface.

11.2.5 CP14 c5, Data Transfer Register

The Data Transfer Register, DTR, is a read/write register. It is software accessible in both user and privileged modes and also as a scan chain from the JTAG interface. Table 11-6 shows the instructions for accessing DTR.

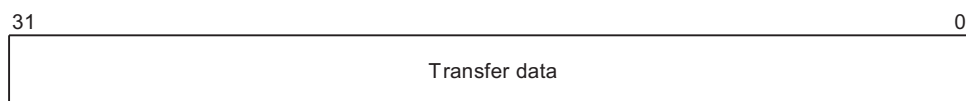
Table 11-6 Data Transfer Register instructions

Instruction	Description
MRC p14, 0, Rd, c0, c5, 0	Copies contents of DTR into Rd

Table 11-6 Data Transfer Register instructions (continued)

Instruction	Description
MCR p14,0,Rd,c0,c5,0	Copies contents of Rd into DTR
LDC p14,c5,<addressing mode>	Loads value accessed in memory into DTR
STC p14,c5,<addressing mode>	Stores contents of DTR to memory

Figure 11-3 shows the DTR bit field.

**Figure 11-3 Data Transfer Register**

———— **Note** ————

Physically, the DTR is two separate registers, the rDTR for reading and the wDTR for writing.

11.2.6 CP14 c6-c63

CP14 registers c6-c63 are reserved.

11.2.7 CP14 c64-c69, Breakpoint Address Registers

The Breakpoint Address Registers, BA0-5, are read/write registers and are accessible only by software operating in privileged mode. Table 11-7 shows the instructions for accessing BA0-5.

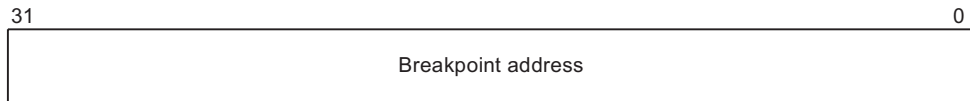
Table 11-7 Breakpoint Address Register instructions

Register	Instruction	Description
CP14 c64, BA0	MRC p14,0,Rd,c0,c0,4	Copies contents of BA0 into Rd
	MCR p14,0,Rd,c0,c0,4	Copies contents of Rd into BA0
CP14 c65, BA1	MRC p14,0,Rd,c0,c1,4	Copies contents of BA1 into Rd
	MCR p14,0,Rd,c0,c1,4	Copies contents of Rd into BA1

Table 11-7 Breakpoint Address Register instructions (continued)

Register	Instruction	Description
CP14 c66, BA2	MRC p14, 0, Rd, c0, c2, 4	Copies contents of BA2 into Rd
	MCR p14, 0, Rd, c0, c2, 4	Copies contents of Rd into BA2
CP14 c67, BA3	MRC p14, 0, Rd, c0, c3, 4	Copies contents of BA3 into Rd
	MCR p14, 0, Rd, c0, c3, 4	Copies contents of Rd into BA3
CP14 c68, BA4	MRC p14, 0, Rd, c0, c4, 4	Copies contents of BA4 into Rd
	MCR p14, 0, Rd, c0, c4, 4	Copies contents of Rd into BA4
CP14 c69, BA5	MRC p14, 0, Rd, c0, c5, 4	Copies contents of BA5 into Rd
	MCR p14, 0, Rd, c0, c5, 4	Copies contents of Rd into BA5

Figure 11-4 shows the BA0-5 bit field.

**Figure 11-4 Breakpoint Address Registers**

11.2.8 CP14 c70-c79

CP14 registers c70-c79 are reserved.

11.2.9 CP14 c80-c85, Breakpoint Control Registers

The Breakpoint Control Registers, BC0-5, are read/write registers and are accessible only by software operating in privileged mode. Table 11-8 shows the instructions for accessing BC0-5.

Table 11-8 Breakpoint Control Register instructions

Register	Instruction	Description
CP14 c80, BC0	MRC p14, 0, Rd, c0, c0, 5	Copies contents of BC0 into Rd
	MCR p14, 0, Rd, c0, c0, 5	Copies contents of Rd into BC0

Table 11-8 Breakpoint Control Register instructions (continued)

Register	Instruction	Description
CP14 c81, BC1	MRC p14, 0, Rd, c0, c1, 5	Copies contents of BC1 into Rd
	MCR p14, 0, Rd, c0, c1, 5	Copies contents of Rd into BC1
CP14 c82, BC2	MRC p14, 0, Rd, c0, c2, 5	Copies contents of BC2 into Rd
	MCR p14, 0, Rd, c0, c2, 5	Copies contents of Rd into BC2
CP14 c83, BC3	MRC p14, 0, Rd, c0, c3, 5	Copies contents of BC3 into Rd
	MCR p14, 0, Rd, c0, c3, 5	Copies contents of Rd into BC3
CP14 c84, BC4	MRC p14, 0, Rd, c0, c4, 5	Copies contents of BC4 into Rd
	MCR p14, 0, Rd, c0, c4, 5	Copies contents of Rd into BC4
CP14 c85, BC5	MRC p14, 0, Rd, c0, c5, 5	Copies contents of BC5 into Rd
	MCR p14, 0, Rd, c0, c5, 5	Copies contents of Rd into BC5

Figure 11-5 shows the BC0-5 bit fields.

**Figure 11-5 Breakpoint Control Registers**

Table 11-9 describes the BC0-5 bit fields.

Table 11-9 Encoding of Breakpoint Control Registers

Bit	Name	Definition
[31:5]	-	Should Be Zero.
[4:3]	IT	Instruction type bit: 00 = reserved 10 = ARM instruction 01 = Thumb instruction 11 = either.
[2:1]	SA	Supervisor access bit: 00 = reserved 10 = privileged 01 = user 11 = either.
[0]	E	Enable bit. Reset clears E: 0 = register disabled 1 = register enabled.

11.2.10 CP14 c86-c95

CP14 registers c86-c95 are reserved.

11.2.11 CP14 c96 and c97, Watchpoint Address Registers

The Watchpoint Address Registers, WA0 and WA1, are read/write registers and are accessible only by software operating in privileged mode. Table 11-10 shows the instructions for accessing WA0 and WA1.

Table 11-10 Watchpoint Address Register instructions

Register	Instruction	Description
CP14 c96, WA0	MRC p14, 0, Rd, c0, c0, 6	Copies contents of WA0 into Rd
	MCR p14, 0, Rd, c0, c0, 6	Copies contents of Rd into WA0
CP14 c97, WA1	MRC p14, 0, Rd, c0, c1, 6	Copies contents of WA1 into Rd
	MCR p14, 0, Rd, c0, c1, 6	Copies contents of Rd into WA1

Figure 11-6 on page 11-13 shows the watchpoint address bit field.

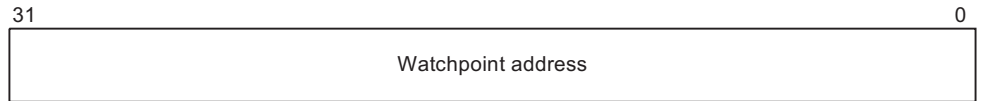


Figure 11-6 Watchpoint Address Registers

11.2.12 CP14 c112 and c113, Watchpoint Control Registers

The Watchpoint Control Registers, WC0 and WC1, are read/write registers and are accessible only by software operating in privileged mode. Table 11-11 shows the instructions for accessing WC0 and WC1.

Table 11-11 Watchpoint Control Register instructions

Register	Instruction	Description
CP14, c112, WC0	MRC p14, 0, Rd, c0, c0, 7	Copies contents of WC0 into Rd
	MCR p14, 0, Rd, c0, c0, 7	Copies contents of Rd into WC0 control
CP14, c113, WC1	MRC p14, 0, Rd, c0, c1, 7	Copies contents of WC1 into Rd
	MCR p14, 0, Rd, c0, c1, 7	Copies contents of Rd into WC1

Figure 11-7 shows the WC0 and WC1 bit fields.



Figure 11-7 Watchpoint Control Registers

Table 11-12 describes the WC0 and WC1 bit fields.

Table 11-12 Encoding of Watchpoint Control Registers

Bits	Definition
[31:11]	Should Be Zero.
[10:9]	DA[1:0] address mask. Bit 10: 1 = exclude DA1 in comparison 0 = include DA1 in comparison. Bit 9: 1 = exclude DA0 in comparison 0 = include DA0 in comparison.
[8]	Should Be Zero.
[7:5]	Byte/halfword/word/any size: 000 = reserved 001 = byte 010 = halfword 011 = byte or halfword 100 = word 101 = word or byte 110 = word or halfword 111 = any size.
[4:3]	Load/store/either: 00 = reserved 10 = load 01 = store 11 = either.
[2:1]	Supervisor: 00 = reserved 10 = privileged 01 = user 11 = either.
[0]	Enable, clear on a system reset 0 = register disabled 1 = register enabled.

11.2.13 CP14 c114-c127

CP14 registers c114-c127 are reserved.

11.3 Software lockout function

When the JTAG debugger is attached to an evaluation board or test system, it indicates its presence by setting the halt/monitor mode bit in the DSCR. When breakpoint and watchpoint registers have been configured, software cannot alter them if the halt/monitor mode bit remains **HIGH** because the debugger retains control. In this mode, software can still write to the comms channel register.

11.4 Halt mode

Halt mode is for debugging the processor using external hardware connected to the JTAG interface. The external hardware provides an interface to a JTAG debugger application. Halt mode can be selected only by setting bit 30 the H bit (bit 30) of the DSCR, which is only writable through the JTAG interface.

In halt mode the processor stops executing instructions if one of the following events occurs:

- an instruction is fetched from a breakpointed memory location
- a data fetch (load or store) occurs from a watchpointed data location
- a breakpoint instruction is executed
- the external **EDBGRQ** signal is asserted
- a HALT instruction has been scanned into the JTAG instruction register
- an exception occurs and the corresponding vector catch bit is set.

When the processor is halted, it is controlled by sending instructions to the integer unit through the JTAG port. Any valid instruction sequence can be scanned into the processor, and the effect of the instruction on the integer unit is as if the instruction is executed under normal operations. Some specific exceptions are described *Sending instructions to the integer unit* and *Using DSCR29 for fast data uploads and downloads* on page 11-17. Also accessible through the JTAG interface is a register to transfer data between CP14 and the JTAG debugger.

The integer unit is restarted by executing a JTAG RESTART instruction.

11.4.1 Sending instructions to the integer unit

Two registers in CP14 are used to communicate with the processor:

- the *Instruction Transfer Register*, ITR
- the *Data Transfer Register*, DTR.

The ITR is used to insert an instruction into the processor pipeline. While in debug state, most of the processor time is spent waiting for a valid instruction in the ITR. Undefined instructions fed to the integer unit through the debugger are *Unpredictable*. Instructions that cause exceptions cause *Unpredictable* behavior. In halt mode, the PC is not incremented as instructions are executed. However, branches and instructions that modify the PC directly update the PC value.

11.4.2 Using DSCR29 for fast data uploads and downloads

DSCR29 enables instructions to be repeatedly issued to the integer unit. When this bit is set, each time the JTAG TAP controller enters the Run-Test/Idle state, the instruction currently residing in the ITR is sent to the prefetch unit for execution. If this bit is clear, no instruction is passed to the prefetch unit. The instruction in the JTAG instruction register must be either INTEST or EXTEST.

The execute feature enables fast uploads and downloads of data. For example, a download sequence might consist of:

1. Scan chain 2, the combination of scan chains 4 and 5, is selected in the ScanNReg, then the JTAG instruction is set to EXTEST for writing.
2. An integer unit write instruction (an STC) and data are loaded into the ITR and DTR, respectively.
3. When the TAP controller passes through the Run-Test/Idle state, the instruction in the ITR is executed by the processor.
4. The scan chain can be switched to the DTR only (chain 5) and polled until the status bit in wDTR0 indicates the completion of the instruction.

More data can then be loaded into DTR and the instruction reexecuted by passing through Run-Test/Idle. The STC instruction must specify base address write-back so that the addresses are automatically updated.

A similar mechanism can increase the performance of upload:

1. First, the JTAG instruction is changed to EXTEST.
2. Using chain 2, a read instruction such as LDC can be scanned into the ITR.
3. The JTAG instruction is switched to INTEST for reading.
4. The scan chain can then be switched to the DTR and polled until the instruction completes. By passing through the Run-Test/Idle state on the way to Shift-DR (for polling), the instruction in the ITR is issued to the integer unit.

Repeat this process until the last word is read.

11.4.3 Accessing processor state

Reading the contents of the integer unit register file requires individual moves from an ARM10 register to CP14 register c5 using MRC and MCR instructions. The data is then scanned out of the DTR.

Byte and halfword transfers are performed by transferring both the address and data into ARM10 registers and then executing the appropriate ARM instructions.

Transfers to and from coprocessors are performed by moving data through an ARM10 register. For this reason all coprocessors must have all data accessible using MRC and MCR (otherwise a data buffer in writable memory must be used).

11.5 Monitor mode

Monitor mode is useful in real-time systems when the integer unit cannot be halted to collect information. Engine controllers and servo mechanisms in hard drive controllers that cannot stop the code without physically damaging the components are examples.

For situations that can only tolerate a small intrusion into the instruction stream, monitor mode is ideal. Using this technique, code can be suspended with an exception long enough to save off state information and important variables. The code continues when the exception handler is finished. The MOE bits in the DSCR can be read to determine what caused the exception.

11.5.1 Entering and exiting monitor mode

Monitor mode is the default mode on Reset. Only an external debugger can change the mode bit in the DSCR. When monitor mode is enabled, the processor takes an exception, rather than halting, if one of the following events occurs:

- a register breakpoint is hit
- a watchpoint is hit
- a breakpoint instruction reaches the Execute stage of the ARM10 pipeline
- an exception is taken and the corresponding vector trap bit is set.

The global debug enable bit in the DSCR must be set or no action is taken. Exiting the exception handler must be done in the normal fashion, for example, restoring the PC to (R14 – 0x4) for prefetch exceptions or moving R14 into the PC for BKPT instructions because they are skipped.

Watchpoints cause Data Abort exceptions. Register breakpoints cause Prefetch Abort exceptions.

11.5.2 Reading and writing Breakpoint and Watchpoint Registers

When in monitor mode, all breakpoint and watchpoint registers can be read and written with MRC and MCR instructions from a privileged processing mode.

11.6 Values in the Link Register after aborts

After an exception, r14, the link register, holds an address for exception processing. This address is used to return after the exception is processed and to address the faulted instruction. BKPT can also generate a Prefetch Abort exception. Prefetch Aborts and Data Aborts might not want to rerun the faulted instruction. BKPT exceptions might or might not want to rerun the instruction at the address of the breakpoint instruction.

Table 11-13 shows the values in the link register after exceptions.

Table 11-13 Values in the Link Register after exceptions

Faulted instruction type	Value left in link register		Address of faulted instruction		Address of following instruction	
	ARM	Thumb	ARM	Thumb	ARM	Thumb
Prefetch Abort	PC + 4	PC + 4	R14 – 4	R14 – 4	R14	R14 – 2
BKPT Used in software debug	PC + 4	PC + 4	R14 – 4	R14 – 4	R14	R14 – 2
Register breakpoint Used in software debug	PC + 4	PC + 4	R14 – 4	R14 – 4	R14	R14 – 2
Data Abort	PC + 8	PC + 8	R14 – 8	R14 – 8	R14 – 4	R14 – 6

For watchpoints, the watchpointed instruction is completed, and the link register points to the instruction at which execution should restart after the handler has finished. The restart address might be several instructions after the faulted instruction.

Table 11-14 shows the values left in the link register and the address of the instruction at which execution must restart.

Table 11-14 Value in the Link Register after a watchpoint

Faulted instruction type	Value left in link register		Address of restart instruction
	ARM	Thumb	ARM
Watchpoint Used in software debug	PC + 8	PC + 8	R14 – 8

11.7 Comms channel

The comms channel is implemented using the two physically separate DTRs and a full/empty bit pair to augment each register, creating a bidirectional data port. One register can be read from the JTAG interface and is written from the ARM10 processor. The other register is written from the JTAG interface and read by the processor. The full/empty bit pair for each register is automatically updated by the debug unit hardware, and is accessible to both the JTAG interface and to software running on the processor.

When the debugger performs comms channel activities, it indicates this to the hardware by setting DSCR27 in scan chain 1. This forces the least significant bit of the wDTR to indicate the state of the comms channel registers.

To read data from the wDTR, the debugger loads the INTEST instruction into the JTAG instruction register and then scans out the contents of the wDTR register. If the LSB of the 33-bit packet of data is HIGH, the data is valid. The bit is then cleared by this read. If the bit is a 0, meaning that the core has not written any data for the debugger, the external hardware can poll the DSCR to see if the core halted.

To write data into the rDTR, the debugger scans the EXTEST instruction into the JTAG instruction register and then scans data into the rDTR. When the debugger goes to write more data, it polls the LSB of the register until the LSB is HIGH. If the LSB is LOW, indicating the rDTR is still full and the core has not read the old data, then the new data shifted in is not loaded into the rDTR.

Because halt mode and monitor mode are mutually exclusive, the transfer registers are not used for any other purpose in monitor mode.

Figure 11-8 illustrates the output from the comms channel.

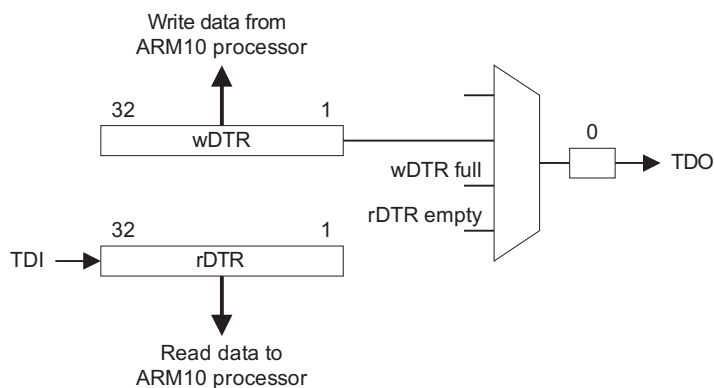


Figure 11-8 Comms channel output

Chapter 12

Instruction Cycle Summary and Interlocks

This chapter gives the instruction cycle counts and examples of interlock timing. This chapter contains the following sections:

- *Cycle timing considerations* on page 12-2
- *Instruction cycle counts* on page 12-3
- *Interlocks* on page 12-22.

12.1 Cycle timing considerations

Complex instruction dependencies make it impossible to describe briefly the exact behavior of all instructions in all circumstances. The tables in this chapter are accurate in most cases but must never be used instead of running code on a cycle-accurate model of the ARM10 processor.

Two performance-enhancing architectural features make it particularly difficult to count the number of cycles an instruction takes:

- branch prediction
- the independent *Load/Store Unit* (LSU).

12.1.1 Branch prediction

With branch prediction enabled, it is impossible to look at a branch in isolation and tell how many cycles it takes. The cycle count depends on where the branch is in memory and what the processor was doing beforehand.

If instruction accesses are hitting in the ICache, then the prefetch buffer is likely to be full. This means the prefetch unit has plenty of time to predict branches and fetch from their targets. In this case, correctly predicted branches look like they take no cycles at all. They are *folded*.

If the prefetch unit was recently flushed, or is fetching from external memory, its buffer can be empty or only partially full. In these cases, the branch predictor does not always have time to completely remove a branch, and it can take one or more cycles before the following instruction is issued. This is described in more detail in *Branch instructions* on page 12-8.

12.1.2 Load/store unit

The independent LSU can process a load or store multiple instruction while data processing operations are executed in the ALU pipeline. However, there are a number of scenarios in which the pipeline is forced to stop and wait for the LSU to complete. The cycle in which the LSU completes a load or store multiple instruction depends on several things:

- how many accesses hit in the cache and TLB
- the 64-bit alignment of the initial access
- the proximity of accesses to a 1K protection region boundary.

This is described in more detail in *Load multiple and store multiple instructions* on page 12-13.

12.2 Instruction cycle counts

Unless stated otherwise, cycle counts and result latencies described here are best case numbers. They assume:

- no outstanding data dependencies between an instruction and a previous instruction
- the instruction does not encounter any resource conflicts
- all data accesses hit in the DCache and do not cross protection region boundaries
- all instruction accesses hit in the ICache.

The tables in this section show the number of cycles an instruction takes to execute and the number of cycles after which the result of the instruction is available to a following instruction. These numbers differ because after an instruction has left the Execute stage of the pipeline, a second instruction can start to execute, even when the first instruction has not produced its final result. This is only the case when the second instruction is not dependent on the result from the first.

Note

Instructions that change the PC cause the pipeline to be flushed and restarted with a fetch of a new instruction. By the time the new instruction executes, it is likely that any dependencies on previous instructions have been cleared.

Three figures are given for each instruction:

Condition pass cycles

This is the number of cycles taken if the instruction passes its condition code check, that is, the number of cycles between this instruction starting to execute and the next instruction starting to execute. This is usually the same as the number of iterations the instruction makes in the Execute stage of the ALU pipeline.

Note

A load or store multiple instruction is a single-cycle operation in the ALU pipeline but iterates in the LSU pipeline until completed.

If an instruction changes the instruction stream, then the *condition pass cycles* indicates the number of cycles before the new PC is available plus the number of cycles it takes to refill the pipe to the point where a new instruction enters Execute in the next cycle.

Condition fail cycles

This is the number of cycles taken if the instruction fails its condition code check, that is, the number of cycles between this instruction entering the Execute stage of the pipeline and failing its condition code check and the next instruction entering the Execute stage.

Result cycles

This is the number of cycles it takes for the instruction to produce its result. It is the number of cycles that must be taken up by the current instruction and following independent instructions before a dependent instruction can be run without interlocking. It can be larger than condition pass cycles in cases where an instruction produces a result later than the Execute stage of the pipeline.

If condition pass cycles is greater than result cycles for an instruction, then the result is always available to a following instruction.

See *Interlocks* on page 12-22 for details of result forwarding paths and the pipeline stages in which instructions have to read registers.

Instructions that change mode by writing the control section of the CPSR are highlighted in some of the tables because they have to wait for the LSU pipe to empty. This is noted in the tables because it makes a significant difference to the execution time if there are any outstanding load misses. Exceptions also change mode, causing a delay while the LSU pipe empties.

The instructions are described in the following sections:

- *Data processing instructions* on page 12-5
- *Multiply instructions* on page 12-7
- *Branch instructions* on page 12-8
- *MRS and MSR instructions* on page 12-9
- *SWI instruction* on page 12-9
- *Load and store instructions* on page 12-9
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- *Thumb multiply instructions* on page 12-19
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- *Thumb load instructions and store instructions* on page 12-20
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12.2.1 Data processing instructions

The simple data processing instructions are:

AND, EOR, SUB, RSB, ADD,

ADC, SBC, RSC, CMN, ORR,

ORR, MOV, BIC, MVN, TST,

TEQ, CMP, QADD, QDADD, QSUB, QDSUB, CLZ

Table 12-1 shows the addressing mode 1 subcategories of data processing instructions.

Table 12-1 Subcategories of data processing instructions

Subcategory	Format	Example
Immediate	OP Rd, Rn, #imm	ADD R1, R2, #1
Register	OP Rd, Rn, Rm	AND R1, R2, R3
Immediate shifted register	OP Rd, Rn, Rm LSL #imm	AND R1, R2, R3 LSL #1
Register shifted register	OP Rd, Rn, Rm LSL Rs	AND R1, R2, R3 LSL R4

Note

A simple unshifted move to the PC (R15) is a special case that operates faster than most data processing operations with the PC as their destination. This enables fast execution of MOV PC, LR, and other simple jumps.

Table 12-2 shows examples of data processing cycle counts. In the table, any of the simple data processing operations can be substituted for AND.

Table 12-2 Cycle counts of data processing instructions

Example instruction	Notes	Change mode	Pass	Fail	Result available
AND Rd, Rn, #imm	-	No	1	1	1
AND Rd, Rn, Rm	-	No	1	1	1
AND Rd, Rn, Rm LSL #imm	-	No	1	1	1
AND Rd, Rn, Rm LSL Rs	-	No	2	2	2
ANDS Rd, Rn, #imm	Set flags	No	1	1	1

Table 12-2 Cycle counts of data processing instructions (continued)

Example instruction		Notes	Change mode	Pass	Fail	Result available
ANDS	Rd, Rn, Rm	Set flags	No	1	1	1
ANDS	Rd, Rn, Rm LSL #imm	Set flags	No	1	1	1
ANDS	Rd, Rn, Rm LSL Rs	Set flags	No	2	2	2
AND	PC, Rn, #imm	To PC	No	1 + 4	1	N/A
AND	PC, Rn, Rm	To PC	No	1 + 4	1	N/A
AND	PC, Rn, Rm LSL #imm	To PC	No	1 + 4	1	N/A
AND	PC, Rn, Rm LSL Rst	To PC	No	2 + 4	2	N/A
ANDS	PC, Rn, #imm	To PC, restore CPSR	Yes	1 + 4	1	N/A
ANDS	PC, Rn, Rm	To PC, restore CPSR	Yes	1 + 4	1	N/A
ANDS	PC, Rn, Rm LSL #imm	To PC, restore CPSR	Yes	1 + 4	1	N/A
ANDS	PC, Rn, Rm LSL Rs	To PC, restore CPSR	Yes	2 + 4	2	N/A
MOV	PC, Rn	Zero shift MOV to PC	No	1 + 3	1	N/A
CLZ	Rd, Rm	-	No	1	1	1
QADD	Rd, Rm, Rn	Sets Q flag	No	1	1	2
QSUB	Rd, Rm, Rn	Sets Q flag	No	1	1	2
QDADD	Rd, Rm, Rn	Sets Q flag	No	1	1	2
QDSUB	Rd, Rm, Rn	Sets Q flag	No	1	1	2

Most data processing instructions take one cycle to execute, after which their result is available for use. The exceptions are instructions that involve register-controlled shifts, saturating instructions, and instructions that write to the PC.

A simple MOV from a register, with no shift that writes the PC requires three extra cycles to refill the pipeline. More complex operations that write to the PC take four extra cycles to refill the pipeline.

12.2.2 Multiply instructions

Table 12-3 shows the cycle counts of multiply instructions. For long multiplies, the least significant word of the result is always the first available. The most significant word is available in the following cycle. This is why there are two cycle counts for instructions whose results extend over one word.

Table 12-3 Cycle counts of multiply instructions

Instruction	Notes	Pass	Fail	Rd (Lo/Hi)	Flags
SMUL<x><y> Rd, Rm, Rs	$16 \times 16 \rightarrow 32$	1	1	2	-
SMLA<x><y> Rd, Rm, Rs, Rn	$16 \times 16 + 32 \rightarrow 32$	2	2	2	-
SMLAL<x><y> RdLo, RdHi, Rm, Rs	$16 \times 16 + 64 \rightarrow 64$	2	2	2/3	-
SMULW<x> Rd, Rm, Rs	$32 \times 16 \rightarrow 32$, upper 32 bits	1	1	2	-
SMLAW<x> Rd, Rm, Rs, Rn	$32 \times 16 + 32 \rightarrow 32$, upper 32 bits	2	2	2	-
MUL Rd, Rm, Rs	$32 \times 32 \rightarrow 32$	2	2	3	-
MULS Rd, Rm, Rs	$32 \times 32 \rightarrow 32$, set flags	4	2	3	4
MLA Rd, Rm, Rs, Rn	$32 \times 32 + 32 \rightarrow 32$	2	2	3	-
MLAS Rd, Rm, Rs, Rn	$32 \times 32 + 32 \rightarrow 32$, set flags	4	2	3	4
UMULL RdLo, RdHi, Rm, Rs	$32 \times 32 \rightarrow 64$, unsigned	3	2	3/4	-
UMULLS RdLo, RdHi, Rm, Rs	$32 \times 32 \rightarrow 64$, unsigned, set flags	5	2	3/4	5
UMLAL RdLo, RdHi, Rm, Rs	$32 \times 32 + 64 \rightarrow 64$, unsigned	3	2	3/4	-
UMLALS RdLo, RdHi, Rm, Rs	$32 \times 32 + 64 \rightarrow 64$, unsigned, set flags	5	2	3/4	5
SMULL RdLo, RdHi, Rm, Rs	$32 \times 32 \rightarrow 64$, signed	3	2	3/4	-
SMULLS RdLo, RdHi, Rm, Rs	$32 \times 32 \rightarrow 64$, signed, set flags	5	2	3/4	5
SMLAL RdLo, RdHi, Rm, Rs	$32 \times 32 + 64 \rightarrow 64$, signed	3	2	3/4	-
SMLALS RdLo, RdHi, Rm, Rs	$32 \times 32 + 64 \rightarrow 64$, signed, set flags	5	2	3/4	5

If the number of pass cycles is greater than the number of result cycles, then the result cycles dominate. Multiplies that set the flags other than Q have to sit in Execute stage for several cycles, because the the ALU must calculate the new flags. Sometimes it might be possible to use a multiply that does not set the flags, followed by a compare of the result that does set the flags. This is appropriate where a useful instruction can be inserted between the multiply and the compare.

12.2.3 Branch instructions

This section describes the following instructions:

B, BL, BX, and BLX.

When branch prediction is enabled, unconditional and conditional backward branches are predicted taken, and conditional forward branches are predicted not taken. See *Branch instruction cycle summary* on page 7-6 for more detail.

Table 12-4 Cycle counts of branch instructions

Instruction	Unpredicted			Predicted	
	Pass	Fail	Predictable	Correctly	Incorrectly
B<address>	4	1	Yes	0 to 2 ^a	4
BL <address>	4	2	Yes	1 to 2	4
BX Rm	4	2	No	-	-
BLX Rm	4	2	No	-	-
BLX <Imm24>	4	2	Yes	1 to 2	4

a. Assuming all accesses hit in the I cache. When the prefetch unit has had time to fold a branch it appears to take 0 cycle. When the prefetch unit has been recently been flushed and is empty it takes 2 cycles to obtain the instruction at the branch target.

12.2.4 MRS and MSR instructions

MSR instructions that write just the flags run quickly. MSRs that change mode take more cycles and have to wait for the LSU pipeline to be empty before they start to execute. Table 12-5 shows the cycle counts for MRS and MSR instructions.

Table 12-5 Cycle counts of MRS and MSR instructions

Example instruction	Notes	Change mode	Pass	Fail
MRS Rd, CPSR	-	No	1	1
MRS Rd, SPSR	-	No	1	1
MSR_f CPSR, Rn	Only flags	No	1	1
MSR_f CPSR, #<cns>	Only flags	No	1	1
MSR CPSR, Rn	Not only flags	Yes	4	1
MSR CPSR, #<cns>	Not only flags	Yes	4	1
MSR SPSR, Rn	-	No	3	2
MSR SPSR, #<cns>	-	No	3	2

12.2.5 SWI instruction

This section describes the SWI instruction:

A SWI instruction takes four cycles, or two cycles if it fails its condition code check. This is true for the ARM and Thumb SWI instructions.

12.2.6 Load and store instructions

This section describes the following instructions:

LDR, LDRD, LDRB, LDRBT, LDRH, LDRSB, LDRSH, LDRT,

STM, STR, STRD, STRB, STRBT, STRH, STRT.

Loads and stores all take one cycle to execute unless they use a scaled register offset, in which case they take two. Loaded data is available for use after one more cycle.

Loads to the PC take six cycles unless they use a scaled register offset, when they take seven. The behavior of load and store multiple instructions is best assessed using a cycle-accurate model of the ARM10 processor.

Table 12-6 shows the cycle counts of the load instructions.

Table 12-6 Cycle counts of load instructions

Example instruction	Pass	Fail	Base write-back result	Load data
LDR PC, [Rn], #<cns>	6	2	1	-
LDR PC, [Rn, #<cns>]	6	2	-	-
LDR PC, [Rn, #<cns>]!	6	2	1	-
LDR PC, [Rn], Rm, <shf><cns>	7	2	2	-
LDR PC, [Rn, Rm]	6	2	-	-
LDR PC, [Rn, Rm]!	6	2	1	-
LDR PC, [Rn, Rm, <shf><cns>]	7	2	-	-
LDR PC, [Rn, Rm, <shf><cns>]!	7	2	2	-
LDR Rd, [Rn], #<cns>	1	1	1	2
LDRT Rd, [Rn], #<cns>	1	1	1	2
LDRB Rd, [Rn], #<cns>	1	1	1	2
LDRBT Rd, [Rn], #<cns>	1	1	1	2
LDR Rd, [Rn, #<cns>]	1	1	-	2
LDR Rd, [Rn, #<cns>]!	1	1	1	2
LDRB Rd, [Rn, #<cns>]	1	1	-	2
LDRB Rd, [Rn, #<cns>]!	1	1	1	2
LDR Rd, [Rn], Rm, <shf><cns>	2	2	1	3
LDRT Rd, [Rn], Rm, <shf><cns>	2	2	1	3
LDRB Rd, [Rn], Rm, <shf><cns>	2	2	1	3
LDRBT Rd, [Rn], Rm, <shf><cns>	2	2	1	3
LDR Rd, [Rn, Rm]	1	1	-	2
LDR Rd, [Rn, Rm]!	1	1	1	2
LDR Rd, [Rn, Rm, <shf><cns>]	1	1	-	3
LDR Rd, [Rn, Rm, <shf><cns>]!	1	1	1	3

Table 12-6 Cycle counts of load instructions (continued)

Example instruction	Pass	Fail	Base write-back result	Load data
LDRB Rd, [Rn, Rm]	1	1	-	2
LDRB Rd, [Rn, Rm]!	1	1	1	2
LDRB Rd, [Rn, Rm, <shf><cns>]	1	1	-	3
LDRB Rd, [Rn, Rm, <shf><cns>]!	1	1	1	3
LDRD Rd, [Rn], Rm	1	1	1	2
LDRD Rd, [Rn], #<cns>	1	1	1	2
LDRD Rd, [Rn, Rm]	1	1	-	2
LDRD Rd, [Rn, Rm]!	1	1	1	2
LDRD Rd, [Rn, #<cns>]	1	1	-	2
LDRD Rd, [Rn, #<cns>]!	1	1	1	2
LDRSB Rd, [Rn], Rm	1	1	1	2
LDRSB Rd, [Rn], #<cns>	1	1	1	2
LDRSB Rd, [Rn, Rm]	1	1	-	2
LDRSB Rd, [Rn, Rm]!	1	1	1	2
LDRSB Rd, [Rn, #<cns>]	1	1	-	2
LDRSB Rd, [Rn, #<cns>]!	1	1	1	2
LDRH Rd, [Rn], Rm	1	1	1	2
LDRH Rd, [Rn], #<cns>	1	1	1	2
LDRH Rd, [Rn, Rm]	1	1	-	2
LDRH Rd, [Rn, Rm]!	1	1	1	2
LDRH Rd, [Rn, #<cnt>]	1	1	-	2
LDRH Rd, [Rn, #<cnt>]!	1	1	1	2
LDRSH Rd, [Rn], Rm	1	1	1	2
LDRSH Rd, [Rn], #<cns>	1	1	1	2
LDRSH Rd, [Rn, Rm]	1	1	-	2

Table 12-6 Cycle counts of load instructions (continued)

Example instruction	Pass	Fail	Base write-back result	Load data
LDRSH Rd, [Rn, Rm]!	1	1	1	2
LDRSH Rd, [Rn, #<cns>]	1	1	-	2
LDRSH Rd, [Rn, #<cns>]!	1	1	1	2

Table 12-7 shows the cycle counts of the store instructions.

Table 12-7 Cycle counts of store instructions

Example instruction	Pass	Fail	Base write-back result
STR Rd, [Rn], #<cns>	1	1	1
STRT Rd, [Rn], #<cns>	1	1	1
STRB Rd, [Rn], #<cns>	1	1	1
STRBT Rd, [Rn], #<cns>	1	1	1
STR Rd, [Rn, #<cns>]	1	1	-
STR Rd, [Rn, #<cns>]!	1	1	1
STRB Rd, [Rn, #<cns>]	1	1	-
STRB Rd, [Rn, #<cns>]!	1	1	1
STR Rd, [Rn], Rm, <shf><cns>	1	1	1
STRT Rd, [Rn], Rm, <shf><cns>	1	1	1
STRB Rd, [Rn], Rm, <shf><cns>	1	1	1
STRBT Rd, [Rn], Rm, <shf><cns>	1	1	1
STR Rd, [Rn, Rm]	1	1	-
STR Rd, [Rn, Rm, <shf><cns>]	2	2	-
STR Rd, [Rn, Rm]!	1	1	1
STR Rd, [Rn, Rm, <shf><cns>]!	2	2	1
STRB Rd, [Rn, Rm]	1	1	-
STRB Rd, [Rn, Rm, <shf><cns>]	2	2	-

Table 12-7 Cycle counts of store instructions (continued)

Example instruction	Pass	Fail	Base write-back result
STRB Rd, [Rn, Rm]!	1	1	1
STRB Rd, [Rn, Rm, <shf><cns>]!	2	2	1
STRH Rd, [Rn], Rm	1	1	1
STRH Rd, [Rn], #<cns>	1	1	1
STRH Rd, [Rn, Rm]	1	1	-
STRH Rd, [Rn, Rm]!	1	1	1
STRH Rd, [Rn, #<cnt>]	1	1	-
STRH Rd, [Rn, #<cnt>]!	1	1	1
STRD Rd, [Rn], Rm	1	1	1
STRD Rd, [Rn], #<cns>	1	1	1
STRD Rd, [Rn, Rm]	1	1	-
STRD Rd, [Rn, Rm]!	1	1	1
STRD Rd, [Rn, #<cns>]	1	1	-
STRD Rd, [Rn, #<cns>]!	1	1	1

12.2.7 Load multiple and store multiple instructions

A simple LDM takes one cycle in Execute after which it operates independently in the load/store pipeline. Following instructions can then execute in the integer pipeline. If a dependent instruction is reached, then the integer pipeline stops until the LDM has loaded the required data or has completed. Dependent instructions are those that require data that has not yet been loaded or those that must be executed in the LSU. Instructions that must be executed in the LSU include all instructions that write to the PC, except branch instructions. An LDM loads two registers per cycle. If the initial access is not to a 64-bit aligned address, an extra cycle is required because only a single register can be loaded in the first cycle.

If an LDM loads the PC, it is loaded from the last access, and five more cycles are required to refill the pipeline. Instructions are not allowed to run under an LDM that changes the processor mode or T bit, or if access is to a noncachable, nonbufferable region of memory.

A simple STM operates the same as an LDM, except that instructions following an STM are held up if they try to write to a register that has not yet been stored. Table 12-8 shows the cycle counts of simple store instructions where L is the number of cycles it takes to load the part of the list before the PC. For example, if the list of registers is {R1, R2, R3, PC}, L is 1 or 2 depending on whether the address to load R1 from is aligned to 64 bits. If it is aligned, R1 and R2 are loaded in one cycle. If not, then it takes one cycle to load R1 and a second cycle to load R2 and R3.

Table 12-8 Cycle counts of load multiple and store multiple instructions

Example instruction	Change mode	Pass	Fail	Write-back	First data
STM Rn, <...>	No	1	1	-	-
STM Rn!, <...>	No	1	1	1	-
STM Rn, <...>^	No	1	1	-	-
STM Rn!, <...>^	No	1	1	1	-
LDM Rn, <...noPC>	No	1	1	-	2
LDM Rn!, <...noPC>	No	1	1	1	2
LDM Rn, <...noPC>^	No	1	1	-	2
LDM Rn!, <...noPC>^	No	1	1	1	2
LDM Rn, <...PC>	No	L + 6	2	-	2
LDM Rn!, <...PC>	No	L + 6	2	1	2
LDM Rn, <...PC>^	Yes	L + 6	2	-	2
LDM Rn!, <...PC>^	Yes	L + 6	2	1	2

12.2.8 Preload instructions

Table 12-9 shows the cycle counts of preload instructions. See *Cachable, Write-Back (WB)* on page 6-10 for more information on this instruction.

Table 12-9 Cycle counts of preload instructions

Instruction	Cycles
PLD [Rn, #-<cns>]	1
PLD [Rn, #<cns>]	1
PLD [Rn, -Rm]	1
PLD [Rn, -Rm, <shf><cns>]	2
PLD [Rn, Rm]	1
PLD [Rn, Rm, <shf><cns>]	2

12.2.9 Coprocessor instructions

This section describes the following instructions:

CDP, LDC, MCR, MCRR, MRC, MRRC, STC.

Table 12-10 shows the cycle counts of the coprocessor instructions. The maximum number of cycles taken by one of these instructions depends on the coprocessor involved. Cycles shown are the minimum cycle count for a tightly coupled coprocessor such as the VFP10 (Rev 1) coprocessor. Other coprocessors may have greater minimum cycle count.

Table 12-10 Cycle counts of coprocessor instructions

Example instruction	Pass	Fail	W/B	Data	Flags
CDP <copr>, <op1>, CRd, CRn, CRm, <op2>	1	1	-	-	-
MCR <copr>, <op1>, Rd, CRn, CRm, <op2>	1	1	-	-	-
MCRR <copr>, <op>, [Rd], [Rn], <CRm>	1	1	-	-	-
MRC <copr>, <op1>, Rd, CRn, CRm, <op2>	1	1	-	2	-
MRC <copr>, <op1>, PC, CRn, CRm, <op2>	2	2	-	2	2
MRRC <copr>, <op>, [Rd], [Rn], <CRm>	1	1	-	2	-

Table 12-10 Cycle counts of coprocessor instructions (continued)

Example instruction	Pass	Fail	W/B	Data	Flags
STC <copr>, CRd, [Rn], {option}	1	1	1	-	-
STC <copr>, CRd, [Rn], #<cns>!	1	1	1	-	-
STCL <copr>, CRd, [Rn], {option}	1	1	1	-	-
STCL <copr>, CRd, [Rn], #<cns>!	1	1	1	-	-
STC <copr>, CRd, [Rn], #<cns>]	1	1	-	-	-
STC <copr>, CRd, [Rn], #<cns>]!	1	1	1	-	-
STCL <copr>, CRd, [Rn], #<cns>]	1	1	-	-	-
STCL <copr>, CRd, [Rn], #<cns>]!	1	1	1	-	-
LDC <copr>, CRd, [Rn], {option}	1	1	1	2	-
LDC <copr>, CRd, [Rn], #<cns>!	1	1	1	2	-
LDCL <copr>, CRd, [Rn], {option}	1	1	1	L + 2	-
LDCL <copr>, CRd, [Rn], #<cns>!	1	1	1	L + 2	-
LDC <copr>, CRd, [Rn], #<cns>]	1	1	-	2	-
LDC <copr>, CRd, [Rn], #<cns>]!	1	1	1	2	-
LDCL <copr>, CRd, [Rn], #<cns>]	1	1	-	L + 2	-
LDCL <copr>, CRd, [Rn], #<cns>]!	1	1	1	L + 2	-

12.2.10 Semaphore instructions

This section describes the following instructions:

SWP and SWPB.

A swap takes two cycles, but before it can be executed, all outstanding loads and stores are completed. Table 12-11 shows the cycle counts of swap instructions.

Table 12-11 Cycle counts of swap instructions

Example instruction	Pass	Fail	Result available
SWP Rd, Rm, [Rn]	2	2	2
SWPB Rd, Rm, [Rn]	2	2	2

12.2.11 Thumb data processing instructions

Thumb data processing instructions behave in a way similar to ARM instructions. Table 12-12 shows the cycle counts of Thumb data processing instructions.

Table 12-12 Cycle counts of Thumb data processing instructions

Example instruction	Number of cycles	Result available
LSL Rd, Rm, #sh_imm5	1	1
LSR Rd, Rm, #sh_imm5	1	1
ASR Rd, Rm, #sh_imm5	1	1
ADD Rd, Rn, Rm	1	1
SUB Rd, Rn, Rm	1	1
ADD Rd, Rn, #imm3	1	1
SUB Rd, Rn, #imm3	1	1
MOV Rd, #imm8	1	1
CMP Rd, #imm8	1	1
ADD Rd, #imm8	1	1
SUB Rd, #imm8	1	1
AND Rd, Rm	1	1
EOR Rd, Rm	1	1
LSL Rd, Rs	2	2
LSR Rd, Rs	2	2

Table 12-12 Cycle counts of Thumb data processing instructions (continued)

Example instruction	Number of cycles	Result available
ASR Rd, Rs	2	2
ADC Rd, Rm	1	1
SBC Rd, Rm	1	1
ROR Rd, Rs	2	2
TST Rn, Rm	1	1
NEG Rd, Rm	1	1
CMP Rd, Rm	1	1
CMN Rd, Rm	1	1
ORR Rd, Rm	1	1
BIC Rd, Rm	1	1
MVN Rd, Rm	1	1
ADD Rd, Hm	1	1
ADD Hd, Rm	1	1
ADD Hd, Hm	1	1
CMP Rd, Hm	1	1
CMP Hd, Rm	1	1
CMP Hd, Hm	1	1
MOV Rd, Hm	1	1
MOV Hd, Rm	1	1
MOV Hd, Hm	1	1
ADD Rd, PC, #imm	1	1
ADD Rd, SP, #imm	1	1
ADD SP, #imm	1	1
SUB SP, #imm	1	1
ADD PC, Rm	5	-

Table 12-12 Cycle counts of Thumb data processing instructions (continued)

Example instruction	Number of cycles	Result available
ADD PC, Hm	5	-
MOV PC, Rm	5	-
MOV PC, Hm	5	-

12.2.12 Thumb multiply instructions

The Thumb multiply instruction behaves in a way similar to the ARM MULS instruction. Table 12-13 shows the cycle count of the Thumb multiply instruction.

Table 12-13 Cycle count of the Thumb multiply instruction

Example instruction	Notes	Number of cycles	Result	
			Rd	Flags
MUL Rd, Rm	$32 \times 32 + 32 \rightarrow 32$, set flags	4	3	4

12.2.13 Thumb branch instructions

Thumb BL and BLX to an immediate value are encoded as two Thumb instructions. The first instruction is a data processing instruction that puts an immediate value into R14. This takes one cycle. The second instruction adds an immediate value to R14 and fetches from that address. This takes four cycles before the next instruction is in Execute. Table 12-14 shows the cycle counts of Thumb branch instructions.

Table 12-14 Cycle counts of Thumb branch instructions

Instruction	Unpredicted			Predicted	
	Pass	Fail	Predictable	Correctly	Incorrectly
B<address>	4	1	Yes	0 to 2 ^a	4
BL <address>	1 + 4	1	Yes	1 to 2	4
BX Rm	4	1	No	-	-
BLX Rm	1 + 4	1	No	-	-
BLX <Imm>	1 + 4	1	Yes	1 to 2	4

- a. Assuming all accesses hit in the I cache. When the prefetch unit has had time to fold a branch it appears to take 0 cycle. When the prefetch unit has been recently flushed and is empty it takes 2 cycles to obtain the instruction at the branch target (See Chapter 7 *Prefetch Unit*).

12.2.14 Thumb SWI instruction

This section describes the SWI instruction:

An SWI instruction takes four cycles, or two cycles if it fails its condition code check. This is true for both the ARM and Thumb SWI instruction.

12.2.15 Thumb load instructions and store instructions

Thumb load/store instructions behave in a way similar to ARM load/store instructions. Table 12-15 shows the cycle counts of Thumb store instructions.

Table 12-15 Cycle counts of Thumb store instruction

Example instruction	Number of cycles	Data
STR Rd, [Rn, Rm]	1	-
STRB Rd, [Rn, Rm]	1	-
STRH Rd, [Rn, Rm]	1	-
STR Rd, [Rb, #imm5]	1	-
STRB Rd, [Rb, #imm5]	1	-
STRH Rd, [Rn, #imm5]	1	-
STR Rd, [SP, #imm8]	1	-

Table 12-16 shows the cycle counts of Thumb load instructions.

Table 12-16 Cycle counts of Thumb load instructions

Example instruction	Number of cycles	Data
LDR Rd, [Rn, Rm]	1	2
LDRB Rd, [Rn, Rm]	1	2
LDRSB Rd, [Rn, Rm]	1	2
LDRH Rd, [Rn, Rm]	1	2

Table 12-16 Cycle counts of Thumb load instructions (continued)

Example instruction	Number of cycles	Data
LDRSH Rd, [Rn, Rm]	1	2
LDR Rd, [Rb, #imm5]	1	2
LDRB Rd, [Rb, #imm5]	1	2
LDRH Rd, [Rn, #imm5]	1	2
LDR Rd, [SP, #imm8]	1	2

12.2.16 Thumb load multiple and store multiple instructions

Thumb load/store multiple instructions behave in the same way as ARM load/store multiple instructions. Table 12-17 shows the cycle counts of Thumb load/store multiple instructions.

Table 12-17 Cycle counts of Thumb load/store multiple instructions

Example instruction	Number of cycles	W/B	First data
PUSH {rlist}	1	-	-
PUSH {rlist, LR}	1	-	-
STMIA Rn!, {rlist}	1	1	-
POP {rlist}	1	-	2
POP {rlist, PC}	L + 6	-	2
LDMIA Rn!, {rlist}	1	1	2

L is the number of cycles it takes to load the part of the list before the PC. For example, for {R1, R2, R3, PC} L is 1 or 2 depending on whether the address to load R1 from is aligned to 64 bits. If it is aligned, R1 and R2 is loaded in one cycle. If not, then it takes one cycle to load R1 and a second cycle to load R2 and R3.

12.3 Interlocks

In almost all cases, the integer core uses forwarding to resolve data dependencies between instructions. For the remaining cases, hardware-imposed interlocks (pipeline stalls) are used to ensure the correct operation of an instruction.

The most common causes of data dependency interlocks are instructions that have a source register that is loaded from memory by the previous instruction. The previous instruction might be an LDR, in which case this data is usually available after a one-cycle interlock. In the case of an LDM, the interlock lasts until the register is loaded. The data processing instruction gets as far as Decode before it interlocks. It interlocks in Decode because this is where it reads its source registers.

Pipeline interlocks are also used to resolve hardware dependencies in the pipeline. Some common examples of hardware dependencies are:

- a new load waiting for the LSU to finish an existing LDM or STM
- a load that misses when the *Hit-Under-Miss* (HUM) slot is already occupied
- a new multiply waiting for a previous multiply to free up the first stage of the multiplier.

The integer core generates most interlocks as late as possible. For instance, a multiply accumulate instruction can start before the accumulate operands are available and stops only when the values are required. This gives the maximum time possible for previous instructions to generate the required data and minimizes occurrences of interlocks.

The integer core implements forwarding paths to enable almost any result to be used as soon as it is calculated. The forwarding paths are shown in Figure 12-1 on page 12-23.

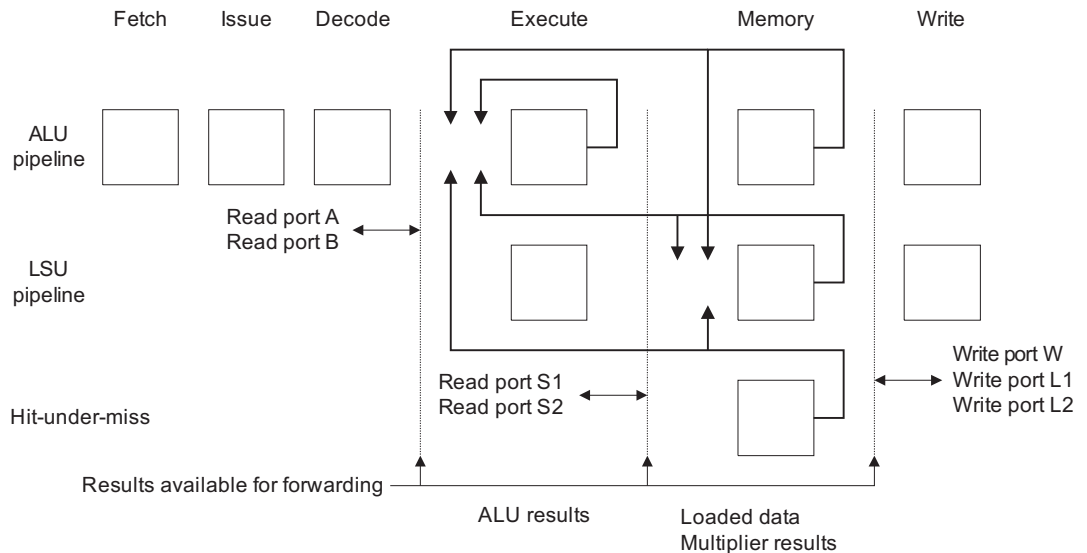


Figure 12-1 Pipeline forwarding paths

The register bank has four read ports:

- Port A
- Port B
- Port S1
- Port S2.

In the second phase of the Decode stage, the integer unit reads port A and port B. Ports A and B are for operands for ALU and multiply instructions and registers to generate addresses for loads, stores, and unpredicted branches.

In the second phase of the Execute stage, the integer unit reads port S1 and port S2. Ports S1 and S2 are for store data for STRs and STMs and for transfers to coprocessors.

The register bank has three write ports:

- Port W
- Port L1
- Port L2.

The integer unit writes to port W, port L1, and port L2 in the first phase of the Write stage. Port W is for writing results from the ALU pipeline. The results include ALU operations, multiplies, and base register write-backs for loads and stores. Ports L1 and L2 are for writing loaded data for LDRs and LDMs and for transfers from coprocessors.

Writes take place in the first phase, so the values are in the registers ready for reads to take place in the second phase. This means there is no need for forwarding paths from Write to earlier stages.

The Execute-to-Execute forwarding paths are used to forward ALU results to following ALU operations.

The Memory-to-Memory forwarding paths are used to forward loaded data to following stores.

The Memory-to-Execute forwarding paths are used to forward one-cycle-old ALU results, freshly loaded data, or multiply results to following ALU operations.

12.3.1 Examples of interlocking and forwarding

Example 12-1 and Example 12-2 illustrate interlocking and forwarding.

Example 12-1 is the simplest case of forwarding. The ADD is dependent on the MOV as the MOV writes R0 and the ADD reads it. The write of 1 into register r0 does not happen until the Write stage of the pipeline, but the correct value for r0, a 1, is forwarded to the ADD at the start of the Execute stage by the Execute-to-Execute forwarding path. This enables the ADD to run with no interlocks.

Example 12-1

```
MOV R0, #1
ADD R1, R0, #1
```

In Example 12-2, the ADD is dependent on the MOV, and there is a single-cycle SUB between them. The write of 1 to r0 has not happened when the ADD is reading its source registers because the MOV is in the Memory stage when the ADD is in the Decode stage. The correct value for r0, a 1, is forwarded to the Execute stage by the ALU pipeline Memory-to-Execute forwarding path. This enables the ADD to run with no interlock.

Example 12-2

```
MOV R0, #1
SUB R1, R2, #2
ADD R2, R0, #1
```

In Example 12-3, the data loaded into r0 is only available at the end of the Memory stage of the LDR, so the ADD interlocks in the Decode stage for one cycle after which the data is available for forwarding to the Execute stage.

Example 12-3

```
LDR R0, [R1, R2]
ADD R3, R0, #1
```

In Example 12-4, the STR data depends on the data loaded by the LDR but there is no interlock because the data is available in time to be forwarded to the Memory stage of the STR.

Example 12-4

```
LDR R0, [R1, R2]
STR R0, [R3, R4]
```

In Example 12-5, the STR address depends on the loaded data from the LDR. In this case there is an interlock for a cycle because the registers used to generate addresses are required in the Execute stage, and r0 is not available until the data is loaded at the end of the Memory stage.

Example 12-5

```
LDR R0, [R1, R2]
STR R3, [R0, R4]
```

In Example 12-6, the source register for the MOV depends on the LDR base write-back to R1. There is no interlock because the write-back value is calculated in the ALU pipeline in the Execute stage and is immediately available for forwarding to the Execute stage of the following instruction.

Example 12-6

```
LDR R0, [R1, R2]!
MOV R3, R1
```

In Example 12-7, there are no data dependencies between the loads. If the first LDR misses in the cache and the HUM slot is empty, then it is assigned to the HUM slot. The second LDR runs underneath it. If the second LDR also misses in the cache, the pipeline interlocks until a load is completed.

Example 12-7

```
LDR R0, [R1, R2]
LDR R3, [R4, R5]
```

In Example 12-8, both loads run without interlocking if they both hit in the cache. If the first LDR misses, the second LDR is held up in the Execute stage to prevent the possibility of having instructions that write to the same register in both the LSU pipeline Memory stage and the HUM buffer.

Example 12-8

```
LDR R0, [R1, R2]
LDR R0, [R1, R2]
```

In Example 12-9, there are no data dependencies between the instructions. There are no interlocks even if the LDR misses, because the data processing instructions can run underneath a miss.

Example 12-9

```
LDR R0, [R1, R2]
ADD R3, R4, R5
SUB R6, R7, R8
```

In Example 12-10, the ADD depends on the LDR. If the LDR hits in the cache, R0 is loaded in time for the ADD to read it without an interlock. If the LDR misses and the data is not returned for a few cycles, then the MOV instructions run underneath the LDR. The ADD interlocks in the Decode stage and waits for loaded the data to be available for forwarding.

Example 12-10

```
LDR R0, [R1, R2]
```

```
MOV R3, R4
MOV R5, R6
ADD R7, R0, R8
```

In Example 12-11, the ADD depends on the LDR. They both write the same register. If the LDR hits in the cache there are no interlocks. If the LDR misses and the data is not returned for a few cycles, then the moves run underneath the LDR. The ADD only gets as far as Memory where it interlocks until R0 has first been written by the LDR.

Example 12-11

```
LDR R0, [R1, R2]
MOV R3, R4
MOV R5, R6
ADD R0, R7, R8
```

In Example 12-12, the LDMIA tries to load R1 first. (Depending upon 64-bit address alignment, R2 might be loaded at the same time as R1.) The MOV is dependent on the LDMIA so it is held up for at least one cycle until the data for R1 is available for forwarding. If the load to R1 (or R1 and R2) misses, then the LDMIA continues until it completes or a second miss occurs. The MOV is always held up until the data loaded to R1 is available.

Example 12-12

```
LDMIA R0, {R1-R7}
MOV R8, R1
```

In Example 12-13 the STR depends on the LDMIA load data. If the LDMIA hits on its first access the data is available to the STR, but the STR cannot run in any case because the LDMIA is occupying the LSU. When the LDMIA is finished, the STR runs. The LDMIA can have up to one miss and still leave the LSU pipeline. The STR then runs under the LDMIA load miss that is in the HUM slot. Clearly there is one case when the STR is still not run, when the LDMIA miss was the load to R1.

Example 12-13

```
LDMIA R0, {R1-R7}
STR R1, [R8, R9]
```

In Example 12-14 there is a data dependency between the LDMIA load data and the MOV source register. Register r7 is the last register to be loaded by the LDMIA so the MOV is held up for a long time.

Example 12-14

```
LDMIA R0, {R1-R7}
MOV R8, R7
```

In Example 12-15 there is a data dependency between the LDMIA load to r5 and the destination register of the MOV. The MOV is held up in the Memory stage of the ALU pipe until the LDMIA has written to r5. In this case there are two different instructions in the Memory stage of the LSU pipe and the ALU pipe both of which write to the same register. This is resolved by always allowing the LSU pipe to write its results first because it always contains the first of the two instructions in program order.

Example 12-15

```
LDMIA R0, {R1-R7}
MOV R5, #1
```

In Example 12-16 there is a data dependency between the LDMIA store of r5 and the destination register of the MOV. The MOV is held up in memory until the STMIA has read r5. The MOV is then allowed to overwrite r5.

Example 12-16

```
STMIA R0, {R1-R7}
MOV R5, #1
```

In Example 12-17 there are no data dependencies between the load multiple instructions. If a single load (one register or two 64-bit aligned registers) from the first LDMIA misses then it is assigned to the HUM slot. The second LDMIA then starts. There are no interlocks if the second LDMIA does not miss until after the miss for the first LDM is resolved.

Example 12-17

```
LDMIA R0, {R1-R7}
```

LDMIA R8, {R9-R13}

Chapter 13

Design for Test

This chapter describes the *Design For Test* (DFT) features of the ARM10 processor and describes how best to integrate the DFT features into a *System on a Chip* (SoC). This chapter contains the following sections:

- *Test modes and ports* on page 13-2
- *Scan chain configuration* on page 13-6
- *Clocks and clock gating* on page 13-8
- *Wrapper cells* on page 13-11
- *Reset* on page 13-17
- *Memories* on page 13-18
- *Memory BIST waveforms* on page 13-27
- *Cache upload/download, manufacturing test* on page 13-34
- *Test signal value tables* on page 13-40.

13.1 Test modes and ports

This section describes the test modes and test ports:

- *ATPG modes*
- *Test ports* on page 13-3
- *Test pinout requirements* on page 13-5.

13.1.1 ATPG modes

A1020WMUXINSEL and **A1020WMUXOUTSEL** configure the wrapper for internal test mode, external test mode, or functional mode.

Writing to **A1020WMUXINSEL** and **A1020WMUXOUTSEL** selects the test mode as shown in Table 13-1.

Table 13-1 ATPG mode selection

Mode	A1020WMUXINSEL	A1020WMUXOUTSEL
Internal test mode	1	0
External test mode	0	1
Functional mode	0	0

Internal test mode

In internal test mode, all input wrapper cells are inward-facing to control core inputs and observe all outputs during test.

Serial core test mode is an internal test mode configuration in which all of the scan chains are connected serially with the wrapper chain attached last. The last cell in the wrapper chain is a lockup latch so that this output can be connected to another clock domain and retain *safe shift* properties. That is, values can be shifted from one scan cell to the next with no risk of error due to clock skew. In this mode, the wrapper clock must be in phase with **GCLK**. Capture cycles cannot occur safely if there are delay differences between the clock domains. **UDLTEST** must be 0 during serial core test mode. The **SCORETEST** signal enables serial core test mode.

External test mode

In external test mode, all input wrapper cells observe external logic and all output wrapper cells control external logic.

13.1.2 Test ports

The test ports in Table 13-2 must be instantiated as specified for ARM10 testing to operate correctly.

Table 13-2 Test port signals

Port name	I/O	Type	Description
A1020DFTCKEN	I	Static	Enables internal core clocks.
A1020SCANEN	I	Dynamic	Scan enable for all internal domains.
A1020SCANMODE	I	Static	Puts device in scan mode.
A1020SCANOUT[23:0]	O	Dynamic	Scan output ports, cache download outputs, memory BIST outputs.
SCANIN[23:0]	I	Dynamic	Scan input ports, cache upload inputs, memory BIST inputs.
A1020DFTRESET	I	Dynamic	Provides direct control over asynchronous reset in scan mode.
A1020TEST	I	Static	Enables cache upload or download mode and BIST test modes.
A1020TESTCFG[2:0]	I	Dynamic	Choose cache upload, download, or BIST test mode.
HRESETN	I	Dynamic	Hard reset.
SFRESETN	I	Dynamic	Soft reset.
TDI	I	Dynamic	JTAG scan-in.
TMS	I	Dynamic	JTAG test mode select.
TCK	I	Dynamic	JTAG test clock.
NTRST	I	Dynamic	JTAG test reset.
TDO	O	Dynamic	JTAG scan-out. Two-state signal externally controlled by ARM10 TDOEN output.

Caution

Because JTAG access occurs with wrappers disabled, JTAG accesses during a cache upload pattern requires additional pin constraints. Lack of constraints on these input pins may result in pattern failure.

A workaround is to constrain the signals during cache upload test as shown in Table 13-3 on page 13-4.

Table 13-3 Cache upload signal constraints

Signal	Connection	Note
CPBUSYD1	0	
CPBUSYE1	0	If from the VFP10, CPBUSY signals can be constrained to the correct state by asserting VFP10SAFE . Tie unused CPBUSY signals to ground.
CPBUSYD2	0	
CPBUSYE2	0	
PMRXACK	0	Acknowledge signals from the power manager are anticipated to be zero after hard reset. Tie unused power manager acknowledge signals to ground.
PMTXACK	0	
FIFOFULL	0	If from the ETM10, FIFOFULL can be constrained to correct state by asserting ETM10SAFE . Tie unused FIFOFULL to ground.

Table 13-4 lists wrapper test signals. There are 24 scan-in and 24 scan-out ports. However, even in 12 or 6 scan chain configuration, a minimum of 16 scan inputs and 16 scan outputs must be ported out to accommodate memory *Built-In Self-Test* (BIST) and cache upload/download modes.

Table 13-4 Test port wrapper signals

Port name	I/O	Type	Description
A1020DFTWCKEN	I	Static	Enables wrapper clock A1020WCLK to dedicated test cells.
A1020RSTSAFE	I	Static	Enables reset of portion of core while testing external logic.
A1020SAFE	I	Static	Forces safe values onto core outputs. Used during ARM10 test.
A1020WCLK	I	Dynamic	Wrapper clock for dedicated wrapper cells.
A1020WMUXINSEL	I	Static	Puts dedicated wrapper cells in internal test mode, external test mode, or functional mode.
A1020WMUXOUTSEL	I	Static	
A1020WSCANEN	I	Dynamic	Scan enable for all wrapper cells.
A1020WSCANOUT[2:0]	O	Dynamic	Output ports for wrapper scan chains.
SCANMUX12	I	Static	Gives access to 12 separate internal scan chains and three wrapper chains. Clearing both SCANMUX12 and SCANMUX6 gives 24 separate internal scan chains and three wrapper chains.

Table 13-4 Test port wrapper signals (continued)

Port name	I/O	Type	Description
SCANMUX6	I	Static	Gives access to six separate internal scan chains and one wrapper chain.
SCORETEST	I	Static	Concatenates all internal and wrapper scan chains.
UDLTEST	I	Static	Enables only shared wrapper cells. Must be asserted during 3-wrapper chain mode.
WSCANIN[2:0]	I	Dynamic	Input ports for wrapper scan chains.

A test control module can be created to control the states of these signals. Table 13-20 on page 13-40.

SCORETEST, **SCANMUX6**, and **SCANMUX12** port states depend on how many scan chains are required during test. When dynamic test signals are connected at chip level, they must make single-cycle timing to the first flip-flop encountered. All signals in Table 13-2 on page 13-3 and Table 13-4 on page 13-4 except **A1020DFTCKEN** must be disabled in functional mode. In functional mode, **A1020DFTCKEN** must be enabled. **UDLTEST** must be LOW for serial core test mode and 6-chain mode. **UDLTEST** must be HIGH for 12-chain mode and 24-chain mode.

13.1.3 Test pinout requirements

Simple and safe implementation of the test pinout in your design requires porting all of the signals listed in Table 13-2 on page 13-3 to your external pinout. Carefully designing a DFT control block can reduce the pin count of the test interface by controlling the static test signals through mode selection. See *Test signal value tables* on page 13-40 for reference tables.

13.2 Scan chain configuration

The ARM10 processor is a partial scan design. Scan chains in the core can be configured as follows:

- 3 wrapper scan chains in 24 core scan chain mode
- 3 wrapper scan chains in 12 core scan chain mode
- 1 wrapper scan chain in 6 core scan chain mode.

Table 13-5 shows how tying **SCANMUX6** and **SCANMUX12** HIGH or LOW selects the scan chain configuration.

Table 13-5 Scan chain configurations

Configuration	SCANMUX12 value	SCANMUX6 value	UDLTES	Maximum chain length
			T value	
24 internal scan chains and 3 wrapper chains	0	0	1	377
12 internal scan chains and 3 wrapper chains	1	0	1	697
6 internal scan chains and 1 wrapper chain	0	1	0	1300
Restricted	1	1	-	-
3 wrapper scan chains	-	-	1	320
1 wrapper scan chain	-	-	0	830
All chains concatenated, serial core test mode	0	0	0	8947

Table 13-6 shows the wrapper scan chain concatenations.

Table 13-6 Wrapper scan chain concatenations

Mode	Scan chains concatenated	Scan-in	Scan-out
SCANMUX12	23, 11	SCANIN11	A1020SCANOUT11
SCANMUX12	22, 10	SCANIN10	A1020SCANOUT10
SCANMUX12	21, 9	SCANIN9	A1020SCANOUT9
SCANMUX12	20, 8	SCANIN8	A1020SCANOUT8
SCANMUX12	19, 7	SCANIN7	A1020SCANOUT7

Table 13-6 Wrapper scan chain concatenations (continued)

Mode	Scan chains concatenated	Scan-in	Scan-out
SCANMUX12	18, 6	SCANIN6	A1020SCANOUT6
SCANMUX12	17, 5	SCANIN5	A1020SCANOUT5
SCANMUX12	16, 4	SCANIN4	A1020SCANOUT4
SCANMUX12	15, 3	SCANIN3	A1020SCANOUT3
SCANMUX12	14, 2	SCANIN2	A1020SCANOUT2
SCANMUX12	13, 1	SCANIN1	A1020SCANOUT1
SCANMUX12	12, 0	SCANIN0	A1020SCANOUT0
SCANMUX6	23, 11, 17, 5	SCANIN5	A1020SCANOUT5
SCANMUX6	22, 10, 16, 4	SCANIN4	A1020SCANOUT4
SCANMUX6	21, 9, 15, 3	SCANIN3	A1020SCANOUT3
SCANMUX6	20, 8, 14, 2	SCANIN2	A1020SCANOUT2
SCANMUX6	19, 7, 13, 1	SCANIN1	A1020SCANOUT1
SCANMUX6	18, 6, 12, 0	SCANIN0	A1020SCANOUT0

13.3 Clocks and clock gating

There are three clock domains in the core and one clock for the dedicated cells in the wrapper:

GCLK	Is the largest clock domain within the core.
HCLK	Is delay-matched with GCLK . HCLK also drives some shared wrapper cells.
TCK	is not synchronized with any other clock domain. It must have separate clock control during the capture cycle.
A1020WCLK	Is the wrapper clock. Its timing is not perfectly delay-matched with any of the other clocks, so take care to prevent hold time failures during test. In production scan mode, A1020WCLK must be $180^\circ \pm 8\%$ out of phase with GCLK . In serial core test mode A1020WCLK must be in phase with GCLK .

Table 13-7 shows the scan chains and the related clock domains.

Table 13-7 Scan chain clocks

Chain name	Scan-in	Scan-out	Maximum chain length	Clock domain
Chain 23	SCANIN[23]	A1020SCANOUT[23]	305	GCLK/TCK
Chain 22	SCANIN[22]	A1020SCANOUT[22]	336	GCLK
Chains 20-0	SCANIN[20:0]	A1020SCANOUT[20:0]	377	GCLK
Chain 21	SCANIN[21]	A1020SCANOUT[21]	332	GCLK/HCLK
Wrappers 2 and 1	WSCANIN[2:1]	A1020WSCANOUT[2:1]	257	A1020WCLK
Wrapper 0	WSCANIN[0]	A1020WSCANOUT[0]	320	HCLK

13.3.1 Scan mode clocking

The ARM10 processor patterns are created with **GCLK** and **HCLK** pin equivalenced. They always have the same activity. You can drive both of these clocks from one clock source. In other words, the test patterns expect **GCLK** and **HCLK** to arrive coincidentally. The timing from the input clock pin or pins must be delay-matched to the **GCLK** and **HCLK** port as shown in Figure 13-1 on page 13-9.

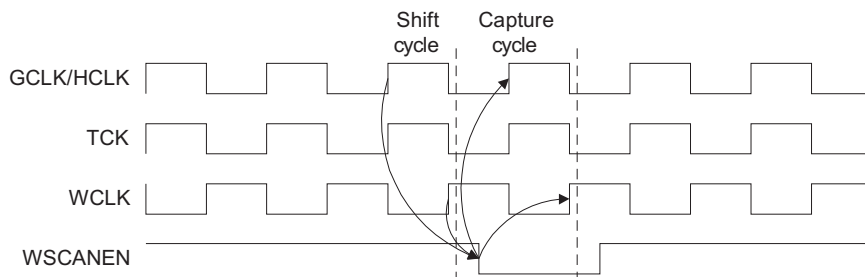


Figure 13-1 Production scan mode clocking

A1020WCLK is 180° out of phase with **GCLK** during production scan mode and any wrapper mode as shown in Figure 13-1. This is to prevent hold timing errors because **GCLK** and **A1020WCLK** are not perfectly delay-matched within the core.

A1020WCLK can be created by inverting **GCLK**, but the timing of these two signals to the ports of the ARM10 processor must be closely matched. **TCK** is not delay-matched with any other clock. During the capture cycle, **TCK** is never toggled at the same time as any other clock on the ARM10 processor. There are lock-up latches in the scan chains wherever they cross clock domains to allow safe shift. The timing to the **TCK** port should be $\frac{1}{4}$ of **GCLK**.

———— **Note** ————

Due to the mixture of shared and dedicated wrapper cells in the wrapper scan chain, **A1020WSCANEN** is the scan enable for both the **HCLK** and **A1020WCLK** domains. To prevent setup or hold time issues for either clock edge, position the edges of **A1020WSCANEN** carefully during use of the wrapper.

13.3.2 Clocking in serial core test mode

During serial core test mode, all scan enables must remain asserted. All clocks are coincident as shown in Figure 13-2 on page 13-10. The scan chains in the ARM10 processor are concatenated into one scan chain with the wrapper scan chain attached last. There is a lock-up latch on the end of the wrapper scan chain. There are also lock-up latches wherever two scan chains from different clock domains are connected.

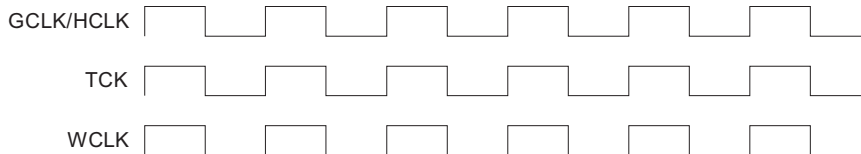


Figure 13-2 Clocking in serial core test mode

13.3.3 Clock gating

A1020DFTCKEN and **A1020DFTWCKEN** are the clock gating signals that gate **GCLK** and **A1020WCLK** respectively. While these signals are enabled, **HCLK** is not gated. In functional mode, **A1020DFTCKEN** must be enabled and **A1020DFTWCKEN** must be disabled. **A1020DFTCKEN** must be enabled whenever **GCLK** is used. **A1020DFTCKEN** can be disabled when **GCLK** is not used. **A1020DFTWCKEN** must be enabled when **A1020WCLK** is used. **A1020DFTWCKEN** must be disabled when **A1020WCLK** is not used.

13.4 Wrapper cells

This section describes the different kinds of wrapper cells:

- *Dedicated input and output wrapper cells*
- *Reset dedicated wrapper cell* on page 13-12
- *Direct control of reset* on page 13-14
- *Shared wrapper cell* on page 13-14.

13.4.1 Dedicated input and output wrapper cells

Figure 13-3 on page 13-12 shows a dedicated input wrapper cell and a dedicated output wrapper cell.

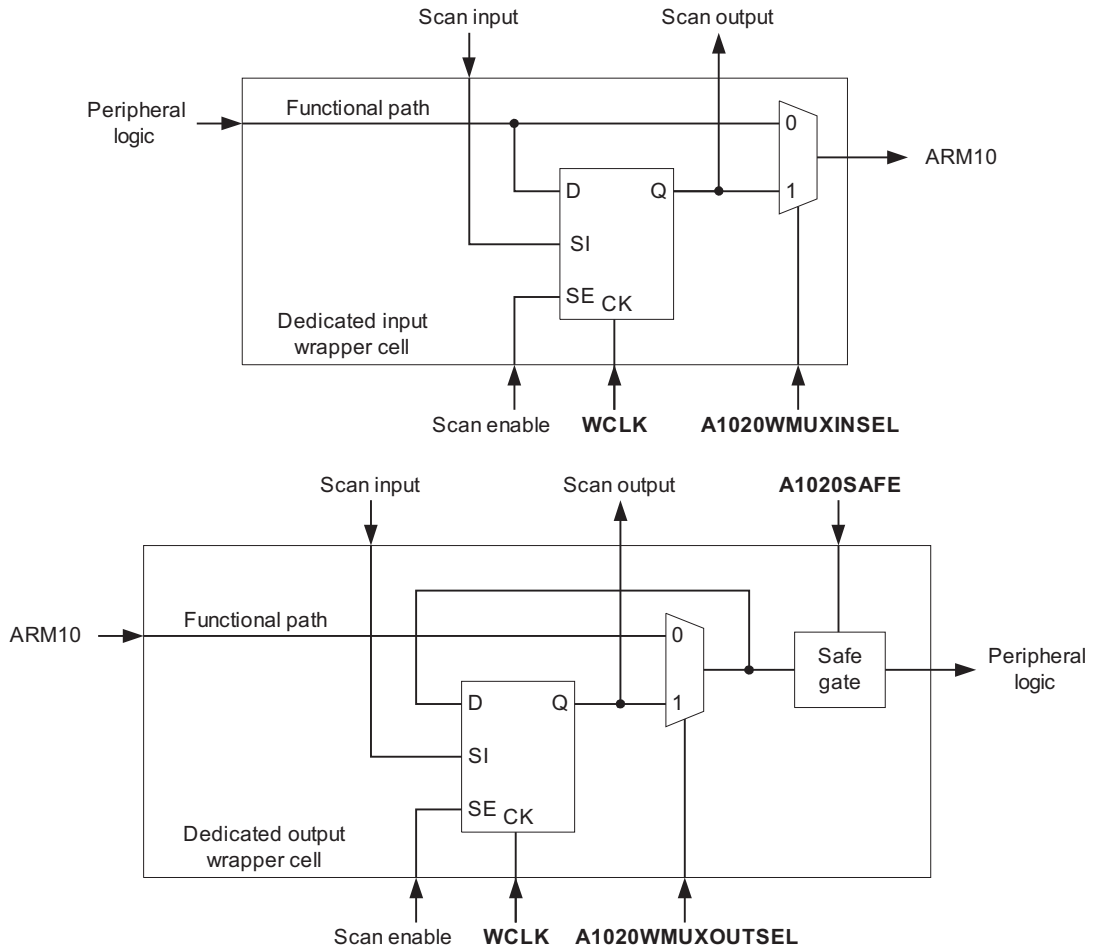


Figure 13-3 Dedicated input and output wrapper cells

13.4.2 Reset dedicated wrapper cell

There is a third type of wrapper cell designed for asynchronous reset inputs. Figure 13-4 on page 13-13 shows the elements of the reset dedicated wrapper cell.

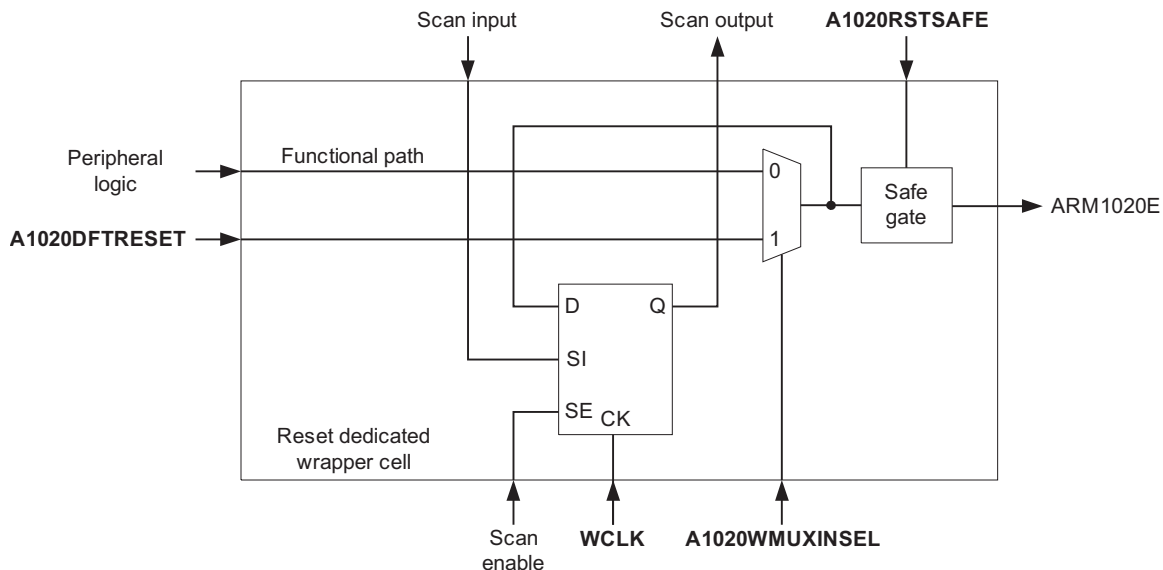


Figure 13-4 Reset dedicated wrapper cell

During external test mode, the safe gate on the reset wrapper cells enables the reset of the core to reduce power and to keep the core safe. In addition, all asynchronous resets are directly controllable during scan mode.

The ARM10 processor has three asynchronous reset inputs:

- **HRESETN**
- **SFRESETN**
- **NTRST**.

The **HRESETN** and **SFRESETN** ports do not have standard reset wrapper cells. The behavior is basically the same as shown in Figure 13-4, except that the **A1020DFTRESET** signal does not override these two signals until after the logic in the power manager block (see Figure 13-5 on page 13-14). The **HRESETN** pin must be controllable by an external pin to reset the power management block at the beginning of each test pattern. This signal must make single-cycle test timing to the flip-flops in the power manager.

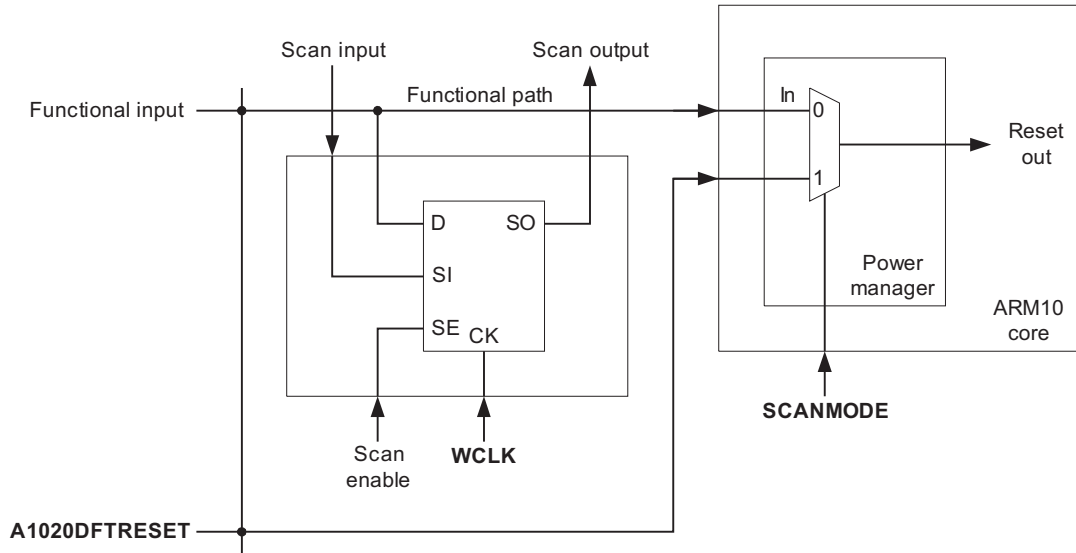


Figure 13-5 HRESET and SFRESET wrapper cell

13.4.3 Direct control of reset

The **A1020DFTRESET** port is a separate port that must be directly connected to a pin for direct control of the reset during test.

In internal test mode, **A1020SAFE** can be asserted so that the values at the output of the core are held in a steady state.

In external test mode, **A1020RSTSAFE** can be asserted, putting the **TCK** domain of the core into reset during external test mode.

13.4.4 Shared wrapper cell

Figure 13-6 on page 13-15 shows a shared wrapper cell. Shared wrapper cells can only be used on registered inputs or outputs, that is, on inputs or outputs on which registers are the closest element to the port. The shared cells in this wrapper are all controlled by **HCLK**.

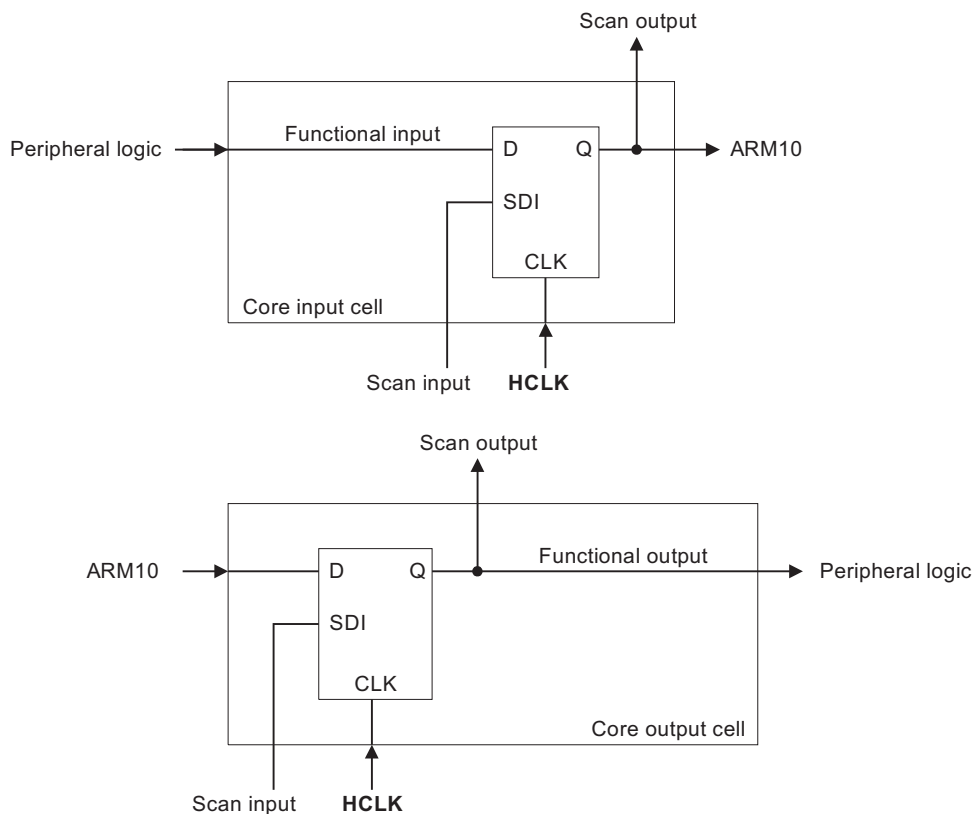


Figure 13-6 Shared wrapper cells

UDLTEST configures the wrapper chain so that only the wrapper cells connected to the **HCLK** domain, all shared, are used, as shown in Figure 13-7 on page 13-16. This provides a shorter wrapper chain while testing unwrapped logic connected to the **HCLK** domain of the ARM10 core.

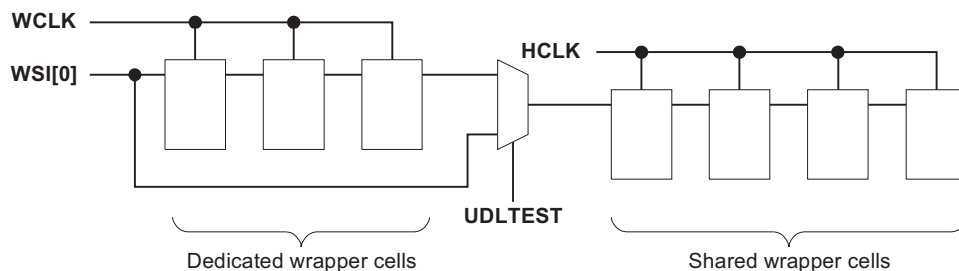


Figure 13-7 HCLK domain wrapper chain isolation

———— **Caution** ————

The following input ports do not have wrapper cells:

- **HRESPI[1:0]**
- **HRESPD[1:0]**.

Wrapper cells are for observing logic external to the core during external scan test mode. If the wrapper cells are not there, and the wrapper is used during test, any logic connected to these ports cannot be observed, and test coverage is affected.

A workaround is to register any external logic connected to these inputs.

13.5 Reset

The ARM10 processor has three asynchronous reset inputs:

- **SFRESETN**
- **HRESETN**
- **NTRST**.

The reset sequence for testing external logic using the ARM10 processor wrapper requires the use of **A1020SCANMODE** and **A1020DFTRESET**.

A1020SCANMODE must be set (see Table 13-24 on page 13-44 and Table 13-25 on page 13-45 for recommended test signal configurations during external testing), and **A1020DFTRESET** must toggle at the beginning of each pattern that uses the ARM10 wrapper to prevent bus contention in the core during external testing.

13.6 Memories

The ARM10 processor memories are all tested with memory BIST. There is also a test mode that allows the cache to have data loaded directly into it to operate the core. The configuration port values for these test modes are in Table 13-8.

13.6.1 Memory BIST and cache upload/download testing

The ARM10 processor supports memory BIST for RAM/PA/flag/CAM arrays inside the ICache, DCache, IMMU, and DMMU blocks. Industry standard patterns and an ARM-specific pattern are available to the user, enabling specific controls of sequences to support textbook fault models as well as high-performance cache RAM failure mechanisms. Insertion of test logic occurs away from the physical cache, piggybacking preexisting data paths. This allows for zero test timing impact on the cache signal interface while supporting full speed test.

The ARM10 processor also supports an extended feature of BIST that enables the user to upload binaries into the ICache and DCache for test execution. This allows for native code based testing in an SoC where specific I/Os are not necessarily available to outside interfaces.

BIST test execution and cache download use the **A1020SCANOUT[15:0]** bus. This bus delivers data from cache downloads and provides real-time BIST execution information.

The ARM10 processor is a hard core. BIST and cache upload execution patterns are delivered in *Condensed Reference Format* (CRF) and are supplied with a recommended test suite.

13.6.2 Test port signal configuration summary

SCANIN[15:0] and **A1020SCANOUT[15:0]** are used for BIST setup and data transfer during upload and download. Table 13-8 shows the **A1020TESTCFG[2:0]** values for uploads and downloads. The wrapper must be initialized before the ICache upload is started.

Table 13-8 Test pin configuration for upload, download, and BIST

A1020TESTCFG[2:0]	Description
000	ICache download
001	DCache download
010	ICache upload

Table 13-8 Test pin configuration for upload, download, and BIST (continued)

A1020TESTCFG[2:0]	Description
011	DCache upload
100	CAMs/flags/PAs upload
101	CAMs/flags/PAs download
110	BIST controller reset and instruction load
111	BIST test

13.6.3 Memory BIST test execution

The test sequence is as follows:

1. Perform a hard reset and then initialize the signals as described in Table 13-22 on page 13-42.
2. With **A1020TESTCFG[2:0]** = 0x6, set **SCANIN[15:0]** to load BIST instruction.
3. With **A1020TESTCFG[2:0]** = 0x7, continually monitor **A1020SCANOUT[15:0]**.
4. Repeat steps 2 and 3 for the next test.

13.6.4 BIST instruction format

The BIST instruction register configures the BIST engine for operation. Writing 0x6 to **A1020TESTCFG[2:0]** at the start of a test sequence asserts BIST engine reset and loads the BIST instruction register from the **SCANIN[15:0]** bus. The last positive edge of **GCLK** delivered to the ARM10 processor during BIST engine reset loads the instruction register. Allow a setup and hold time of more than two **GCLK** cycles for BIST instruction register loading before starting execution.

Table 13-9 shows the BIST instruction fields captured from **SCANIN[15:0]** when **A1020TESTCFG[2:0] = 0x6**.

Table 13-9 Encoding of BIST instruction fields

SCANIN bits	Description
[15:12]	Engine control
[11:8]	Block under test
[7:4]	Data word
[3:0]	BIST pattern

Engine control description

Table 13-10 describes the BIST instruction register control field.

Table 13-10 Encoding of BIST engine control field

SCANIN[15:12]	Description
0000	Normal BIST test execution, runs to completion. Used during BIST test.
0001	Stop on error, stops 2-3 cycles after error detection. Used during upload tests.
1111	Run cache test, executes native code on completion of upload. Used during upload tests.

Block description address size

Table 13-11 on page 13-21 shows how the BIST block under test field selects blocks in terms of x and y coordinates.

———— **Note** —————

CAM BIST does not include compare logic.

Table 13-11 Encoding of BIST block under test field

SCANIN[11:8]	Block	Address size (x, y)
0000	Cache CAM	2^{11} ($2^6, 2^5$)
0001	Cache RAM	2^{11} ($2^6, 2^5$)
0010	Cache PA, flags	2^{11} ($2^6, 2^5$)
1000	MMU CAM	2^6 ($2^6, 2^0$)
1001	MMU RAM	2^6 ($2^6, 2^0$)
1010	MMU PA	2^6 ($2^6, 2^0$)

Data word description

Table 13-12 shows how the BIST data word field selects the data word.

Table 13-12 Encoding of BIST data word field

SCANIN[7:4]	Test	Description
xxxx	BIST	Root data word.
xxx0	Upload	Parallel upload, instruction and data side cache upload.
xxx1	Upload	Serial upload, instruction or data side cache upload.

BIST patterns

Table 13-13 shows how the BIST pattern field selects test patterns. N is the number of times each memory cell is accessed.

Table 13-13 Encoding of BIST pattern field

SCANIN[3:0]	Description	N
0000	WriteSolids	1
0001	ReadSolids	1
0010	WriteCkdb	1
0011	ReadCkdb	1

Table 13-13 Encoding of BIST pattern field (continued)

SCANIN[3:0]	Description	N
0100	RowMarch, wordline fast	6
0101	ColMarch, bitline fast	6
0110	Bang, bitline fast write/read stress tests	18
1111	Bang, bitline fast write/read stress tests	2

BIST pattern descriptions

All BIST execution is performed with a physically mapped address space. This means that the least significant Xaddress switches between adjacent rows. For example, LSB + 1 switches between every second row. Yaddress space is also physically mapped for efficient and direct targeting of memory faults with the supplied patterns.

Table 13-14 lists the definitions for terms used in the BIST patterns.

Table 13-14 BIST pattern terms and definitions

Term	Definition
Column	Dimension in array parallel to bitlines on same sense amp
Row	Dimension in array parallel to wordlines
Row fast/Xfast	Target cell moves along bitlines before moving to next column
Col fast/Yfast	Target cell moves across bitline pairs before row/wordline
Xfast increment	Target cell begins nearest sense amp, moves away
Xfast decrement	Target cell begins furthest point from sense amp, moves closer
Yfast increment	Yaddr space moves from 0 to maximum, east-west relationship
Yfast decrement	Yaddr space moves from maximum to 0, opposite of increment

The following patterns are used:

WriteCkbd Is performed Xfast. This pattern is 1N, writing only. Data polarity is set by xor(Xaddr0, Yaddr0).

ReadCkbd Is performed Xfast. This pattern is 1N, reading only. Data polarity set by xor(Xaddr0, Yaddr0).

WriteSolids Is performed Xfast. This pattern is 1N, writing only. Data polarity = true.

ReadSolids Is performed Xfast. This pattern is 1N, reading only. Data polarity = true.

RowMarch Is performed Xfast. This 6N pattern has the following sequence:

1. WriteSolids, initialize array.
2. Read data/write databar increment.
3. Read databar/write data decrement.
4. Read data solid.

ColMarch Is 6N and performed Yfast with the same sequence as RowMarch.

PttnFail Is performed Xfast. It executes a WriteSolid pattern followed by a ReadSolid. Fails are injected by reversing data polarity on select addresses during ReadSolid. This pattern is required to insure BIST detection logic at the target array is functional.

Bang Is 18N, and performed Xfast, executing consecutive multiple writes and reads on a bitline pair.

The sequence is as follows:

1. WriteSolid, initialize array.
2. Read data target, write databar target, repeat write databar six times
This segment *bangs* bitline pairs insuring proper equalization after writes. Insufficient equalization or precharge causes slow reads when opposite data is read from the same bitline pair. Slow reads in self-timed caches result in functional failure not found in single-shot algorithms like March C-. This segment stresses bitline pullup and equalization so that a memory cell read may have to overcome an opposite bitline differential, missing critical sense-amp timing.
3. Repeat read databar target five times, write data row 0, read databar target, and write data target.
This segment walks down a bitcell, writes opposite data on that bitline pair, and reads target cell data. This failure mechanism is less common in 6T RAM cells compared to 4T or DRAM.
Using the sacrificial row also helps detect open decoder faults in the Xaddr space (Yaddr not subject to fault class architecturally) in the absence of Gray code pattern sequences. This pattern detects stuck-at faults, but its primary purpose is to address the analog characteristics of the memories. It is more effective in stressing bitline recovery than March C-.
4. Read data, verify array.

13.6.5 Mapping and description of memory BIST test monitors

The **A1020SCANOUT[15:0]** bus provides real-time data, allowing monitoring of test progress and pass/fail behavior. The bus becomes active for strobing after **A1020TESTCFG[2:0]** are changed from 0x6 (BIST reset) to 0x7 (BIST execute). On completion of the algorithm, the finished flag is set, and all **A1020SCANOUT[15:0]** outputs are sourced by registered sticky signals.

Table 13-15 A1020SCANOUT[15:0] mapping

Bits	Description
[15:10]	Unused
9	BIST done flag, current algorithm finished
8	Xaddr expire, set whenever Xaddr = maxAddr
7	Yaddr expire, set whenever Yaddr = maxAddr
6	Unused
5	I-side MMU failure
4	Data-side MMU failure
3	Instruction-side CAM/flag failure
2	Data-side CAM/flag failure
1	Instruction-side RAM failure
0	Data-side RAM failure

The **A1020SCANOUT** bus data meets timing requirements at the ARM10 processor interface. Because this bus can be routed throughout the SoC, timing failures might occur on the **A1020SCANOUT** strobe at the tester. Timing delay between the ARM10 processor interface and external pins must be accounted for in timing. Do not *set_false_path* this bus, even though scan is a substring of the net name.

———— **Note** —————

Failure flags toggle throughout test during normal BIST execution whenever a change in fail or pass status occurs. This information can be data logged to gain understanding of fail behavior. Once the BIST done flag has been set, fail flags are held if any failures were observed during test.

13.6.6 Memory BIST failure analysis

Direct bit mapping of array failures is not available in this version of the ARM10 processor. Understanding of the failure type can be obtained by:

- analyzing real time failure flags on **A1020SCANOUT**
- running a comprehensive BIST test suite, for example, using solids and dataword changes
- using the cache download mechanism described in *Cache upload/download, manufacturing test* on page 13-34.

Future ARM10 processors might have more direct bit mapping features installed. The cache dump mechanism does not support the MMU but does allow for determination of failing bits found during test.

Each real-time failure flag has a latency in relation to address expire flags due to internal pipelines. The information in Table 13-16 can be used to determine failure address. *Cycle#* is the cycle count between address expire and fail flag observations.

Table 13-16 Failure address formulas

Block	Latency	Xaddr formula	Yaddr Formula
ICache RAM	6	$(\text{cycle\#} - 64 \times \text{int}((\text{cycle\#} - 6) / 128)) / 2$	$\text{int}((\text{cycle\#} - 6) / 128)$
DCache RAM	8	$(\text{cycle\#} - 64 \times \text{int}((\text{cycle\#} - 8) / 128)) / 2$	$\text{int}((\text{cycle\#} - 8) / 128)$
ICache CAM	7	$(\text{cycle\#} - 64 \times \text{int}((\text{cycle\#} - 7) / 128)) / 2$	$\text{int}((\text{cycle\#} - 7) / 128)$
DCache CAM	7	$(\text{cycle\#} - 64 \times \text{int}((\text{cycle\#} - 7) / 128)) / 2$	$\text{int}((\text{cycle\#} - 7) / 128)$
ICache PA	7	$(\text{cycle\#} - 64 \times \text{int}((\text{cycle\#} - 7) / 128)) / 2$	$\text{int}((\text{cycle\#} - 7) / 128)$
DCache PA	7	$(\text{cycle\#} - 64 \times \text{int}((\text{cycle\#} - 7) / 128)) / 2$	$\text{int}((\text{cycle\#} - 7) / 128)$
MMU RAM	7	$(\text{cycle\#} - 64 \times \text{int}((\text{cycle\#} - 7) / 128)) / 2$	$\text{int}((\text{cycle\#} - 7) / 128)$
MMU CAM	7	$(\text{cycle\#} - 64 \times \text{int}((\text{cycle\#} - 7) / 128)) / 2$	$\text{int}((\text{cycle\#} - 7) / 128)$
MMU PA	7	$(\text{cycle\#} - 64 \times \text{int}((\text{cycle\#} - 7) / 128)) / 2$	$\text{int}((\text{cycle\#} - 7) / 128)$

Figure 13-11 on page 13-33 shows an example failure waveform highlighting Xaddr = 1, Yaddr = 0 failure in the ICache and DCache.

13.6.7 Memory BIST test suite

Vectors are provided in CRF format to exercise the defined test interface. The sequence 1-4 provides simple gross functional stuck-at tests. Patterns 1-5 establish fundamental cell integrity in a manner that provides gross functional yield data prior to engaging stress tests.

The pattern set comprises data words 0x9 and 0xA used in the following sequence:

1. WCKbd
Data word: 0x9
2. RCKbd 5s
Data word: 0x9
3. WCKbd As
Data word: 0xA
4. RCKbd
Data word: 0xA
5. Repeat patterns 1, 2, 3, 4 with 200ms extreme voltage pause to insure adequate data retention.
6. Y-fast March Decrement
Dataword: 0x6
This is a fundamental column fast pattern.
7. X-fast BANG
Dataword: 0x0
This provides bitline stress testing.

13.7 Memory BIST waveforms

The waveform diagrams associated with common memory BIST operations are shown in the following sections:

- *Reset followed by BIST test*
- *Test completion followed by new test* on page 13-28
- *Example of real-time failure* on page 13-30
- *Test termination, failure observed* on page 13-32.

13.7.1 Reset followed by BIST test

Figure 13-8 shows on release of reset assertions that the **SCANIN[15:0]** bus value of $0x02f0$ is captured while **A1020TESTCFG[2:0]** = $0x6$. BIST test execution is allowed once **A1020TESTCFG[2:0]** = $0x7$.

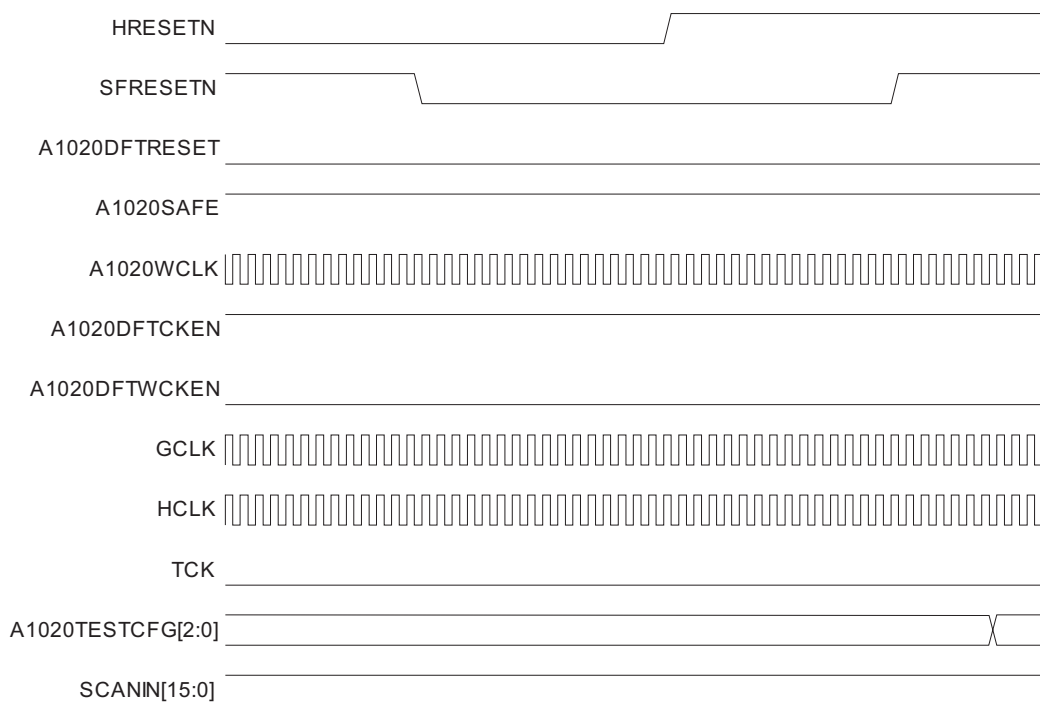


Figure 13-8 Reset followed by BIST test

Table 13-17 shows the **SCANIN[15:0]** values for reset followed by BIST test.

Table 13-17 Instruction fields for reset followed by BIST test

SCANIN bits	Value	Description
[15:12]	0000	Normal test execution
[11:8]	0010	Cache PA/flags
[7:4]	1111	Root data word
[3:0]	0000	WriteSolids pattern

13.7.2 Test completion followed by new test

In Figure 13-9 on page 13-29 completion of WriteSolids test occurs. Both Xaddr expire 8 and Yaddr expire 7 are set when the respective maxAddr occurs as defined in Table 13-15 on page 13-24. Completion flag 9 is set and no failures are observed. A second test is initiated by writing **A1020TESTCFG[2:0] = 0x6** and beginning the next test, column march.

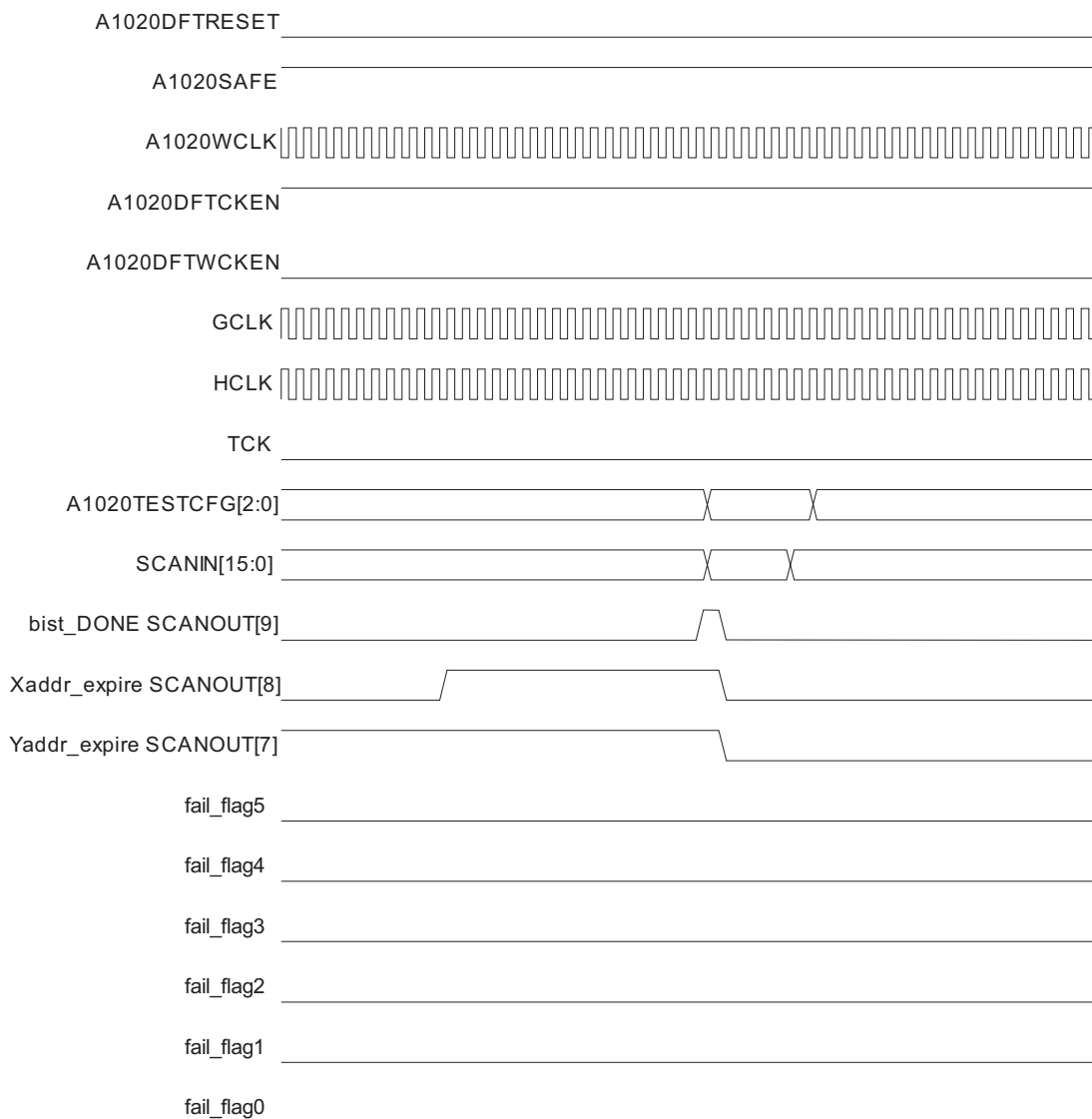


Figure 13-9 Test completion followed by a new test

Table 13-18 shows the values for the operations described in Figure 13-9 on page 13-29.

Table 13-18 Instruction fields for test completion followed by new test

SCANIN bits	Values	Description
[15:12]	0000	Normal test execution
[11:8]	0001	Cache RAM
[7:4]	1110	Root data word
[3:0]	0101	Column march, Yfast

13.7.3 Example of real-time failure

Figure 13-10 on page 13-31 shows a real-time failure flag being set. The fail was created for Xaddr = 0x1 and is shown for the RW increment portion of the test, pattern = 0xf is PtnnFail. The WriteSolids portion of the algorithm completed when both Xaddr and Yaddr expires were set.

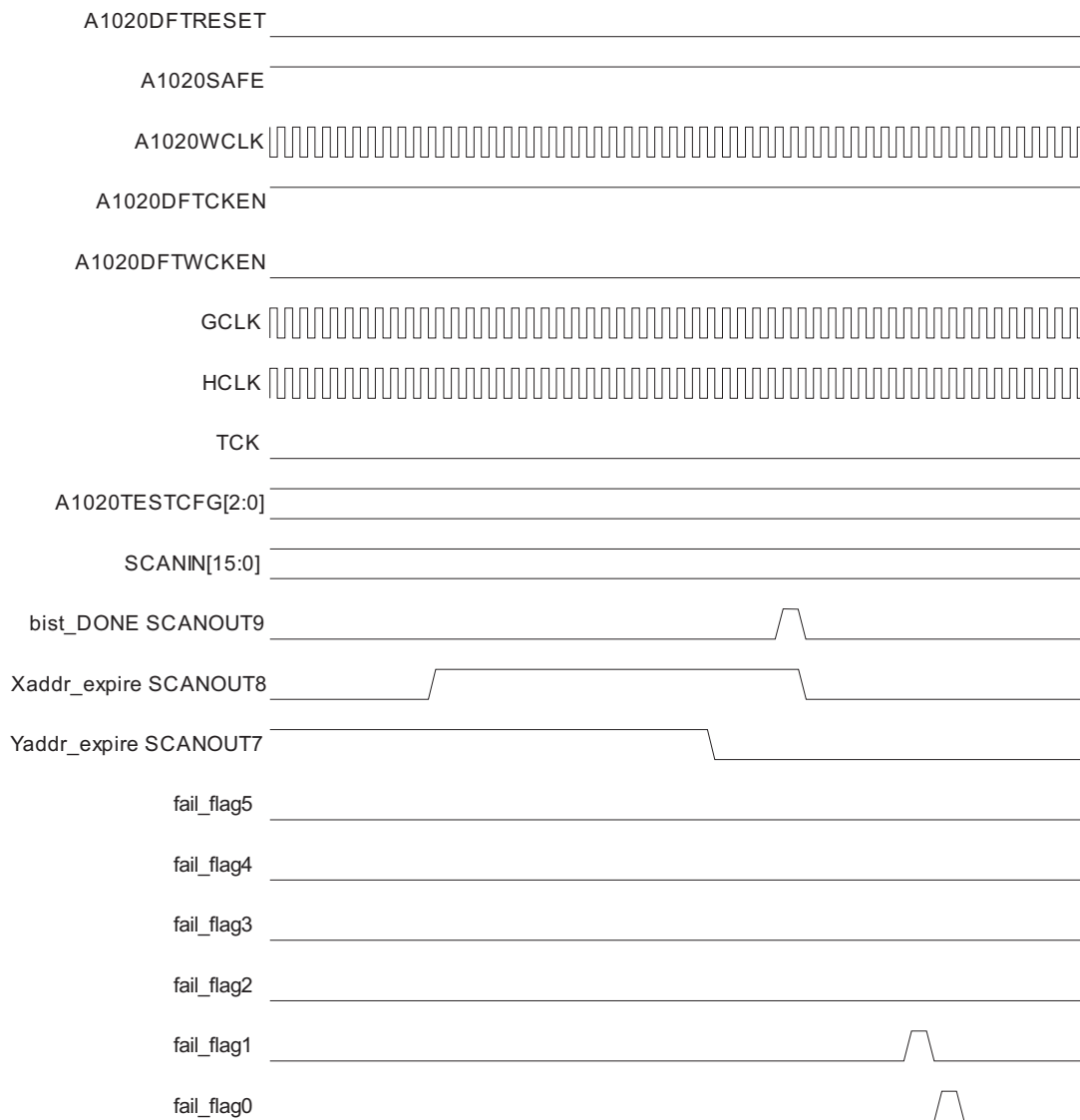


Figure 13-10 Setting a real time failure flag

13.7.4 Test termination, failure observed

Figure 13-11 on page 13-33 shows the completion of the pattern fail test. A sticky version of the failure flag is set when the **BIST_DONE9** signal is asserted. These values remain on the bus until a BIST engine reset is performed.

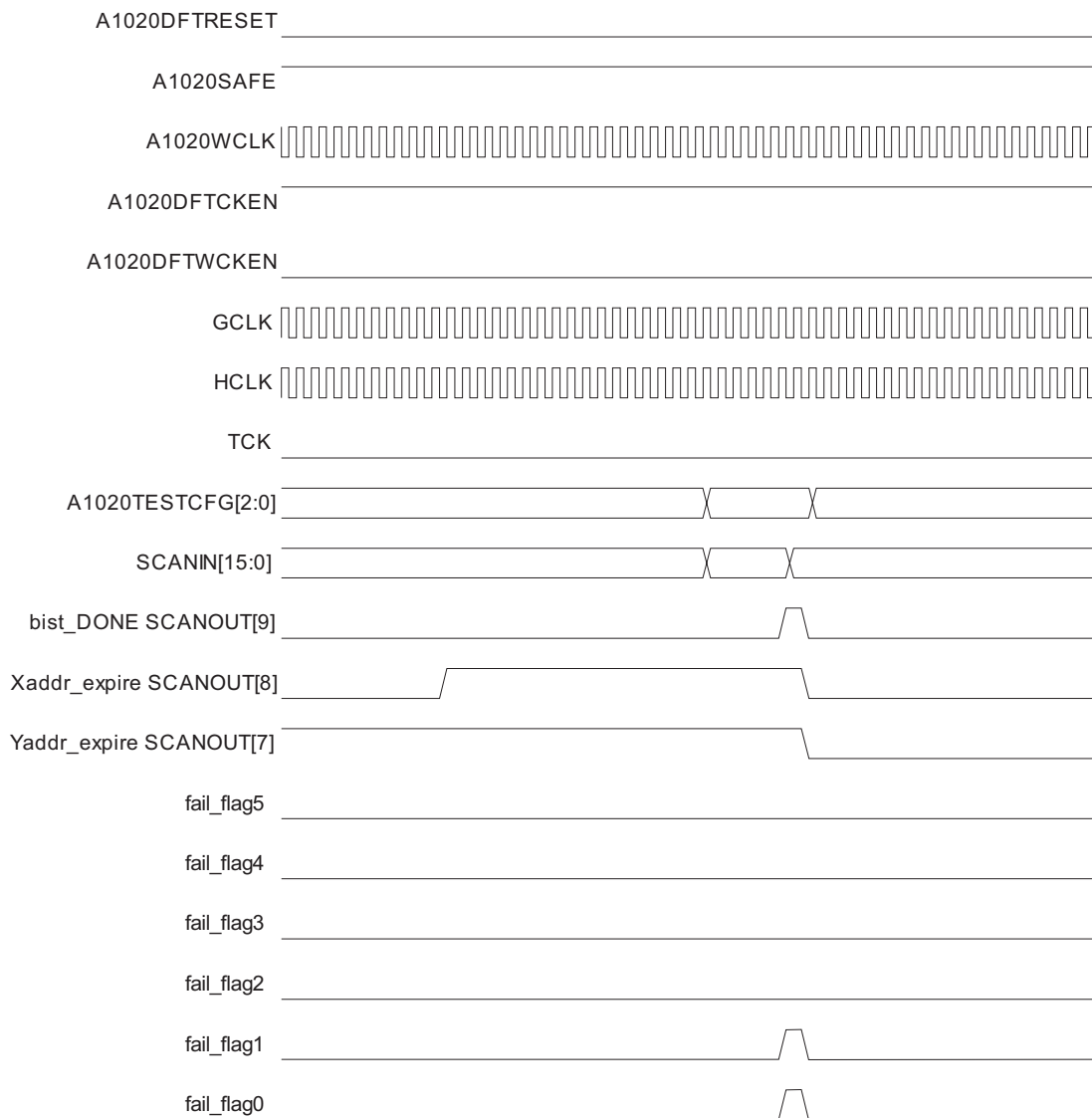


Figure 13-11 Completion of pattern fail test

13.8 Cache upload/download, manufacturing test

The ARM10 processor is partial scan and does not achieve high coverage by *Automated Test Pattern Generation* (ATPG) alone. The nonscan area of the processor is tested by partial scan ATPG and is supplemented with the cache upload test mechanism. To minimize design impact, the memory BIST logic was shared for this feature which allows us to directly load caches with functional test binaries in the ARM10 processor. This test can be performed in an SoC environment without external ARM10 processor bus transactions. External bus transactions and their supporting logic are fully scanned and are tested by ATPG. L1 caches and flags are loaded and downloaded by this mechanism but the MMU arrays are not supported. All tests supplied by ARM are locally resident in the L1 and self-deterministic. Expected results are also loaded for comparison against CPU-created results.

It is assumed that testing the ARM10 processor in an SoC environment occurs with no access to functional pins. Therefore, all functional patterns are self-contained (no external bus accesses are allowed) and self-deterministic.

All cache upload patterns are provided and fault graded by ARM Limited. The upload information described here is for information purposes only. The upload feature is designed to maximize ARM10 test coverage and cannot be used to test logic external to the ARM10 processor.

13.8.1 Test port signal configuration

Table 13-8 on page 13-18 shows the values for **A1020TESTCFG[2:0]** for cache upload and cache download tests. **SCANIN[15:0]** and **A1020SCANOUT[15:0]** are used as a data transfer bus. They are also used for monitoring of cache-loaded test patterns.

13.8.2 Cache upload test execution

Cache upload tests that use JTAG must be able to disable the wrapper during test in order for valid **TDO** to be created. During cache upload test execution, **A1020WMUXINSEL** and **A1020SAFE** must toggle. See the waveform in *Cache upload test execution* on page 13-36.

The sequence of operations is as follows:

1. Perform a hard reset.
2. Load the wrapper chain with the required values to prevent the external bus from disrupting execution.

NFIQ, **NIRQ**, **ISYNC**, and **CPBOUNCEE1** are set. All other input signals are cleared. **DBGEN** is set for some patterns.

3. Initialize **SCANIN[15:0]** to set BIST controls/target array with **A1020TESTCFG[2:0] = 0x6**.

The pattern selected must be WriteSolids/ReadSolids.

For the last array upload, write 0xF to the engine control field, **SCANIN[15:12]**, to execute code on completion. In this mode, the BIST controller performs a test-only soft reset to the ARM10 processor. This overrides default CP15 POR states to allow for immediate execution from caches. Writing 0xF to the engine control field before the last array being loaded causes Unpredictable behavior.

4. Write 0x0-0x5 to **A1020TESTCFG[2:0]** to upload or download values using the **SCANIN[23:0]** and **A1020SCANOUT[23:0]** buses.

The BIST controller increments address every fourth cycle when **A1020TESTCFG[2:0] = 0x0-0x5**. This allows 64-bit entries to be constructed. The BIST controller creates sequential addressing and enables paths to the arrays. See *BIST instruction format* on page 13-19 for encoding of BIST instructions.

5. Write **A1020TESTCFG[2:0] = 0x6** for early termination of upload for patterns less than array size.
6. Repeat steps 3-5 for next array.

———— **Note** —————

The upload/download configuration can be terminated at any time by setting **A1020TESTCFG[2:0] = 0x6**. This allows for reduced vector count when loading programs that do not require the entire address space. Perform an early termination only after the last required entry has been completely written. Termination during the upload of the last address produces Undefined data for that cache line.

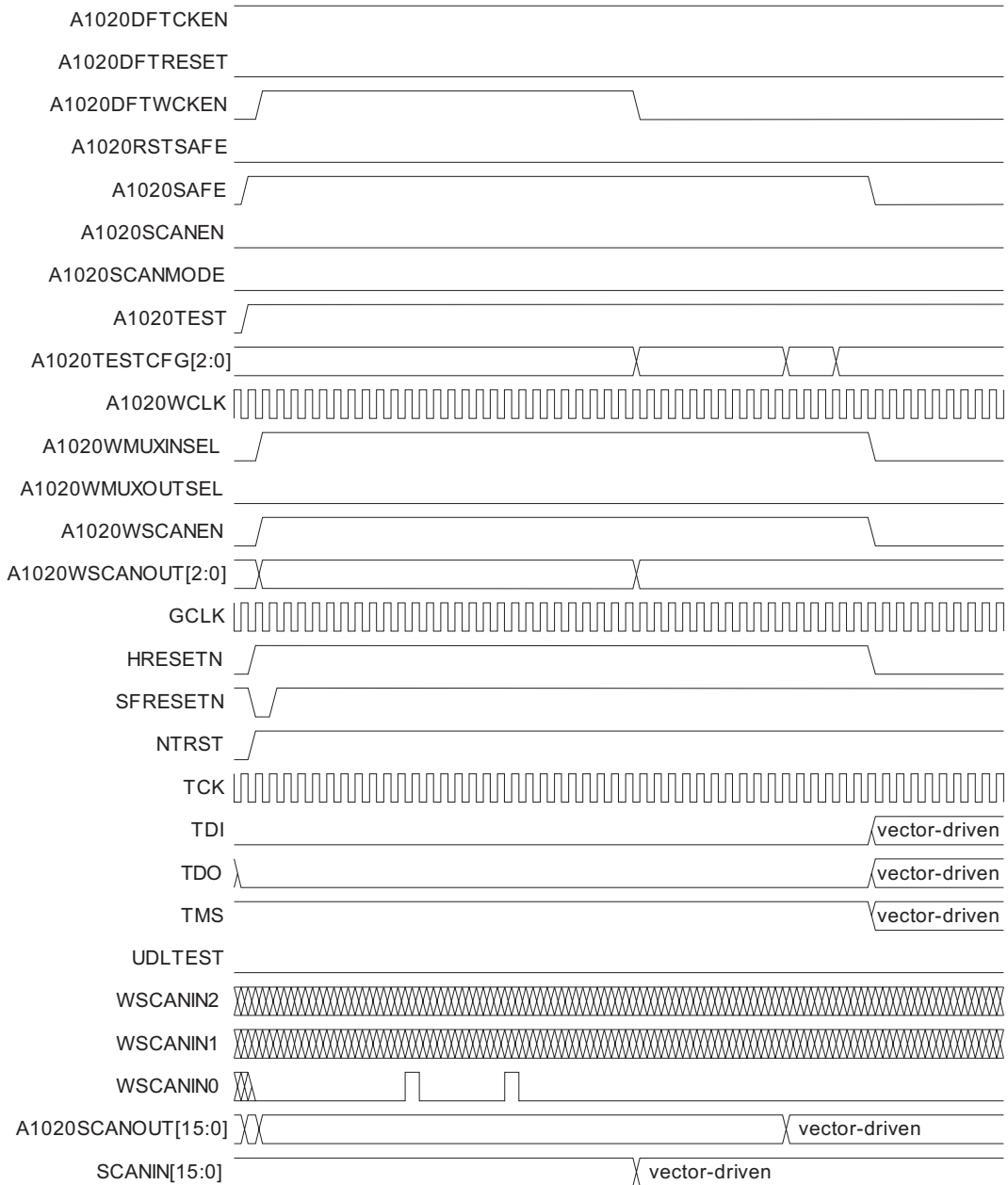


Figure 13-12 Cache upload test execution

13.8.3 Cache download test execution and waveforms

The cache download feature can be used for determining cache values at any time. Such reads of the caches are destructive and the device should be reset after data is downloaded. There are four pattern sets delivered for cache download. Datalogs of download tests can be used to bitmap failing bits. The downloads consist of reading all zeros, ones, and reading of checkerboard backgrounds produced by root datawords of 0xA and 0x5. The expected data pattern sets provided are those commonly found at the termination of provided BIST test patterns. A supplied README file describes cycle numbers where data entries appear on the **A1020SCANOUT** bus.

When using such patterns for debug, us care to insure not to cause a device reset between BIST test execution and download. Such resets invalidate cache entries.

13.8.4 Execution of binary test download

When **A1020TESTCFG[2:0]** moves from reset (0x6) to execution (0x0), ICache download begins, as shown in Table 13-8 on page 13-18. The first data read occurs 11 cycles later. Table 13-19 shows the cache download values for **A1020TESTCFG[2:0]**.

Table 13-19 Instruction fields for cache download

SCANIN bits	Values	Description
[15:12]	0000	Normal test execution
[11:8]	0001	Cache RAM
[7:4]	Don't care	-
[3:0]	0001	ReadSolids

Figure 13-13 on page 13-38 shows the execution of cache download start.

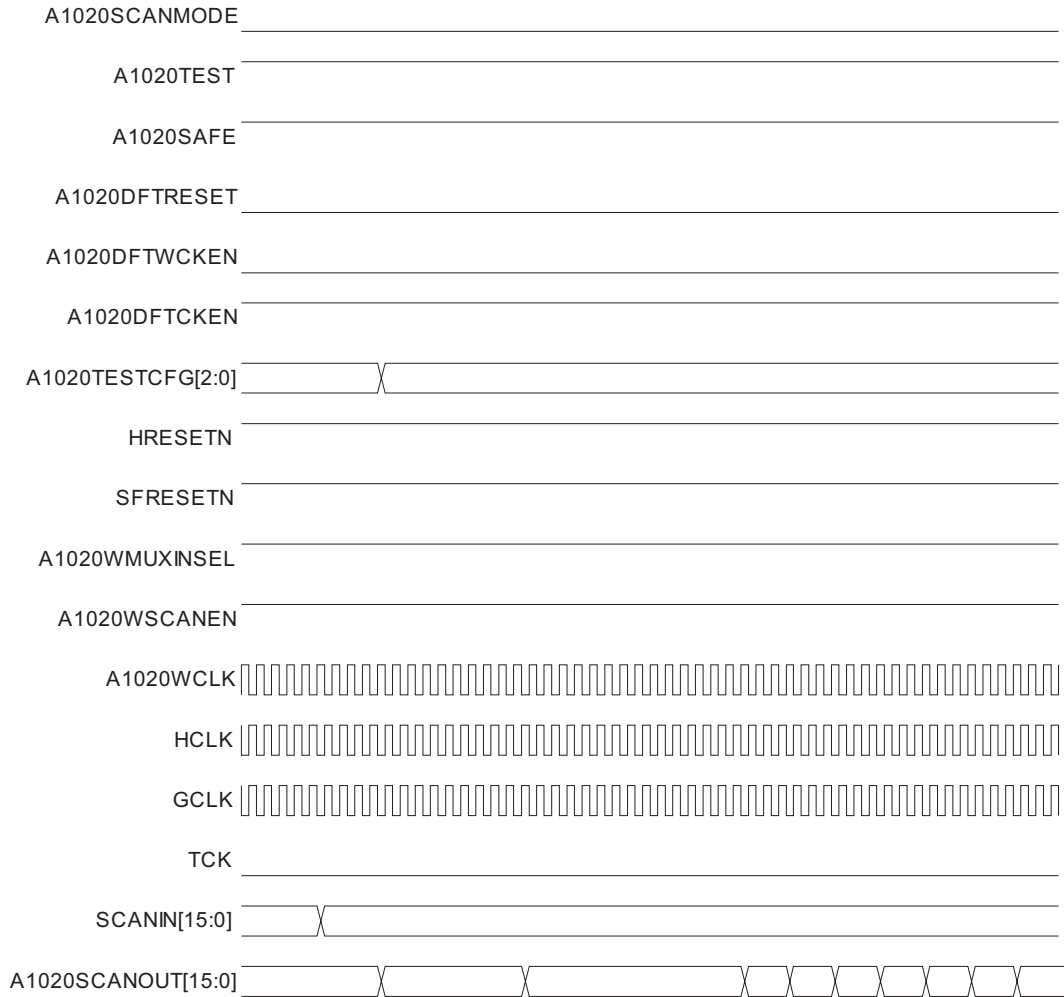


Figure 13-13 Execution of cache download start

13.8.5 Transition of download tests

Figure 13-14 on page 13-39 shows the transition of download tests. Completion of ICache load is followed by a BIST engine reset and a load of the same engine control register with settings $0x01C1$. **A1020TESTCFG[2:0]** = $0x1$, which defines DCache download. Other arrays are read by repeating the process with **A1020TESTCFG[2:0]** settings shown in Table 13-8 on page 13-18.

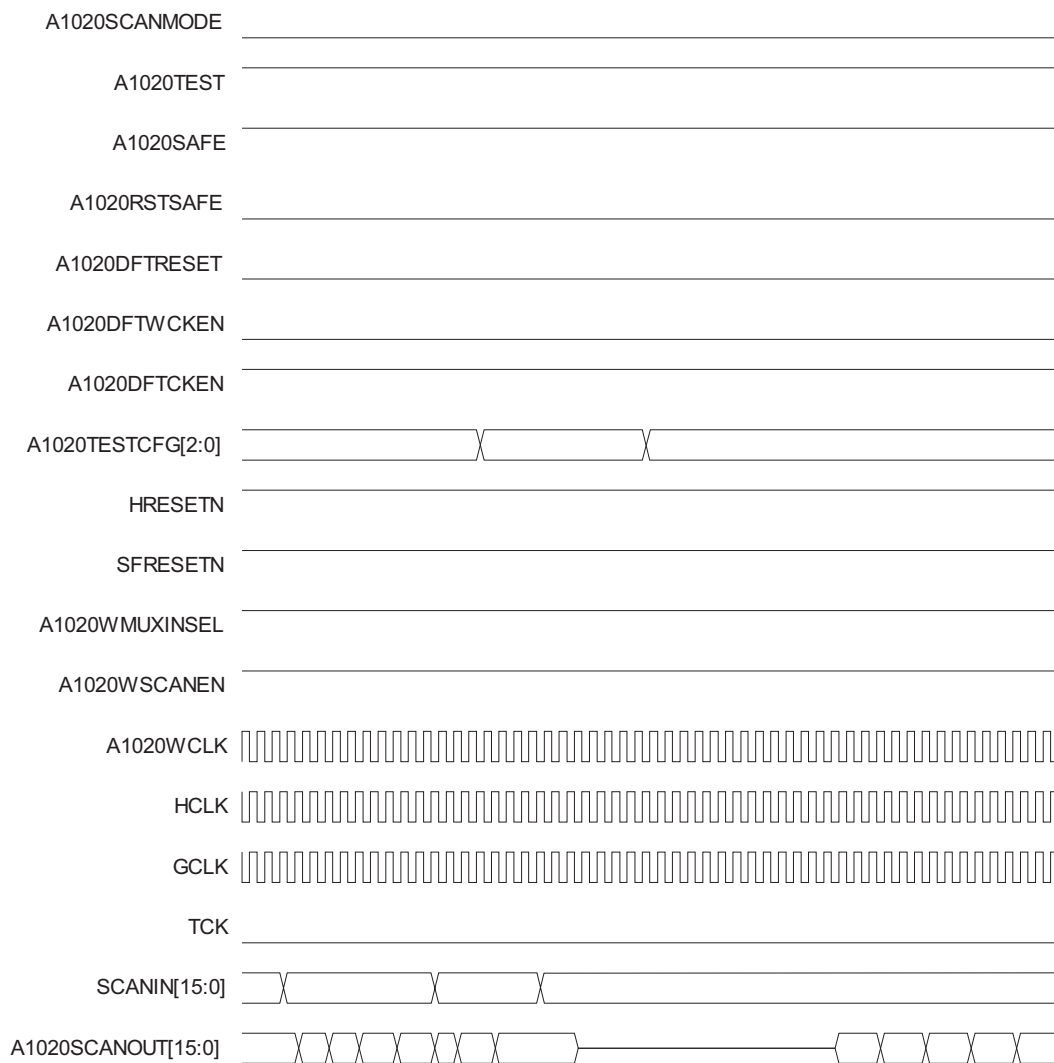


Figure 13-14 Execution of binary test download

Note
MMU download is not supported.

13.9 Test signal value tables

This section contains signal test value tables for the following test modes:

- *Test signals for ATPG testing*
- *Test signals in functional mode* on page 13-41
- *Test signals in cache upload mode* on page 13-42
- *Test signals in external test wrapper mode with one wrapper chain* on page 13-44
- *Test signals in external test wrapper mode with three wrapper chains* on page 13-45.

Table 13-20 shows the ARM10 test signal values for ATPG testing.

Table 13-20 Test signals for ATPG testing

Test signals	Connection
A1020TEST	1
A1020SCANMODE	1
A1020DFTCKEN	1
A1020DFTWCKEN	1
A1020SCANEN	Connect to an external pin
A1020WSCANEN	Connect to an external pin
A1020DFTRESET	Connect to an external pin
A1020MUXINSEL	1
A1020MUXOUTSEL	0
A1020SAFE	1 recommended
A1020RSTSAFE	0
SCANIN	Connect to external pins
A1020SCANOUT	Connect to external pins
UDLTEST	0 if 6-chain pattern or serial core test mode, else 1
WSCANIN	Connect to external pin

Table 13-20 Test signals for ATPG testing (continued)

Test signals	Connection
SCORETEST	0, unless serial scan pattern
SCANMUX6	Dependent upon pattern set, see <i>Scan chain configurations</i> on page 13-6
SCANMUX12	Dependent upon pattern set, see <i>Scan chain configurations</i> on page 13-6

Table 13-21 shows test signals values in functional mode.

Table 13-21 Test signals in functional mode

Test signals	Connection
A1020TEST	0
A1020SCANMODE	0
A1020DFTCKEN	1
A1020DFTWCKEN	0
A1020SCANEN	0
A1020WSCANEN	0
A1020DFTRESET	0 recommended
A1020MUXINSEL	0
A1020MUXOUTSEL	0
A1020SAFE	0
A1020RSTSAFE	0
SCANIN	0 recommended
A1020SCANOUT	-
UDLTEST	NA
SCANMUX6	NA
SCANMUX12	NA
WSCANIN	0 recommended

Table 13-22 shows the test signal values for memory BIST testing.

Table 13-22 Test signals during BIST testing

Signal	Value
A1020SCANMODE	0
A1020SCANEN	0
A1020DFTCKEN	1
A1020DFTRESET	0
A1020DFTWCKEN	0 recommended
A1020WSCANEN	0 recommended
A1020WMUXINSEL	0 recommended
A1020WMUXOUTSEL	0 recommended
A1020SAFE	1 recommended
A1020RSTSAFE	0
A1020TEST	1
A1020TESTCFG[2:0]	Connect to external pin
SFRESETN	Connect to external pin
HRESETN	Connect to external pin

Table 13-23 shows test signals values in cache upload mode.

Table 13-23 Test signals in cache upload mode

Test signals	Connection	Note
A1020TEST	1	
A1020SCANMODE	0	
A1020DFTCKEN	1	
A1020DFTWCKEN	Connect to external pin	
A1020SCANEN	Connect to external pin	
A1020WSCANEN	1	

Table 13-23 Test signals in cache upload mode (continued)

Test signals	Connection	Note
A1020DFTRRESET	0	
A1020MUXINSEL	Connect to external pin	
A1020MUXOUTSEL	0	
A1020SAFE	Connect to external pin	
A1020RSTSAFE	0	
SCANIN	Connect to external pins	
A1020SCANOUT	Connect to external pins	
UDLTEST	0	
SCANMUX6	1	
SCANMUX12	0	
SFRESETN	Connect to external pin	
HRESETN	Connect to external pin	
TDI	Connect to external pin	
TDO	Connect to external pin	
TMS	Connect to external pin	
NTRST	Connect to external pin	
A1020TESTCFG	Connect to external pin	
CPBUSYD1	0	
CPBUSYE1	0	If from the VFP10, CPBUSY signals can be constrained to the correct state by asserting VFP10SAFE . Tie unused CPBUSY signals to ground.
CPBUSYD2	0	
CPBUSYE2	0	
PMRXACK	0	Acknowledge signals from the power manager are anticipated to be zero after hard reset. Tie unused power manager acknowledge signals to ground.
PMTXACK	0	
FIFOFULL	0	If from the ETM10, FIFOFULL can be constrained to correct state by asserting ETM10SAFE . Tie unused FIFOFULL to ground.

Table 13-24 shows test signals values in external test wrapper mode with one wrapper chain.

Table 13-24 Test signals in external test wrapper mode with one wrapper chain

Test signal	Connection
A1020TEST	1
A1020SCANMODE	1
A1020DFTCKEN	0
A1020DFTWCKEN	1
A1020SCANEN	0
A1020WSCANEN	Connect to external pin
A1020DFTRESET	Connect to external pin
A1020MUXINSEL	0
A1020MUXOUTSEL	1
A1020SAFE	0
A1020RSTSAFE	1 recommended
SCANIN	0
A1020SCANOUT	Not needed
A1020WSCANOUT	Connect to external pin or another scan chain
WSCANIN	Connect to external pin
UDLTEST	0
SCANMUX6	1
SCANMUX12	0
SFRESETN	Connect to external pin
HRESETN	Connect to external pin

Table 13-25 shows test signals values in external test wrapper mode with three wrapper chains.

Table 13-25 Test signals in external test wrapper mode with three wrapper chains

Test signals	Connection
A1020TEST	1
A1020SCANMODE	1
A1020DFTCKEN	0
A1020DFTWCKEN	1
A1020SCANEN	0
A1020WSCANEN	Connect to an external pin
A1020DFTRESET	Connect to an external pin
A1020MUXINSEL	0
A1020MUXOUTSEL	1
A1020SAFE	0
A1020RSTSAFE	1 recommended
SCANIN	0
A1020SCANOUT	Not needed
A1020WSCANOUT	Connect to external pin or another scan chain
WSCANIN	Connect to external pin
UDLTEST	0
SCANMUX6	1
SCANMUX12	0
SCORETEST	0
SFRESETN	Connect to external pin
HRESETN	Connect to external pin

Chapter 14

Power Manager

This chapter describes the power manager and its extensible, memory map independent ARM10 processor interface. It contains the following sections:

- *About the power manager* on page 14-2
- *ARM10 processor power modes* on page 14-3
- *System control coprocessor* on page 14-8
- *Programming examples* on page 14-14
- *Power manager interface* on page 14-16
- *Timing* on page 14-17
- *Software example code sequences* on page 14-21.

14.1 About the power manager

Typical system-level power manager functions are built as application-specific hardware. For example, memory-mapped hardware registers are programmed to turn off subsystem clocks. In high-performance processes, however, leakage can be significant even when clocks are stopped, and a generic power management interface is required.

The ARM10 power manager interface is not memory-mapped and is extensible to accommodate process-driven voltage ranges and frequencies.

The NORMAL and OFF states are the minimum state set required to support power management.

14.1.1 Power management hardware requirements

In a system that includes a single or multiple processors, each ARM10 processor must have a power management isolation layer. The *lock-out* layer isolates the ARM10 processor from the system bus, placing the ARM10 processor bus in the IDLE state. The lock-out layer is similar to the layer that is provided in the ARM10 processor cache for isolation of clock, reset, and control signals from ARM10 processor signals.

14.2 ARM10 processor power modes

The processor supports the power modes listed in Table 14-1.

———— **Note** ————

In this chapter, the term *processor core state* refers to the state of:

- all banked registers
- the CPSR
- the MMU TLB
- the system control coprocessor, CP15
- the debug coprocessor, CP14
- the VFP10 coprocessor
- the ETM10.

Table 14-1 ARM10 processor power modes

Mode	Description	Recovery time
RUN	Processor executing instructions and able to program the power manager.	≥ 1 cycle
STANDBY	Processor clocks stopped. Return to RUN mode on interrupt request or external debug request.	$\geq 10^1$ cycle
DORMANT	Processor core state must be saved in external memory. If processor and caches have separate power rails, caches held in reduced-leakage state. Return to RUN mode on soft reset or power-on reset.	$\geq 10^2$ cycle
SHUTDOWN	Processor core state and cache states must be saved in external memory. Processor and caches powered down. Return to RUN mode on hard reset.	$\geq 10^4$ cycles

The *recovery time* is the time it takes the processor to reenter RUN mode and resume executing instructions. While in RUN mode, the recovery time is the time it takes the processor to change from one system power mode to another.

All normal transitions from RUN mode are due to commands written under program control. The power manager controls the sequencing back to RUN mode from any of the power-saving modes so that voltage supply rails and clocks are running after the appropriate wakeup or reset condition.

Figure 14-1 is a state diagram of the four power modes of the ARM10 processor.

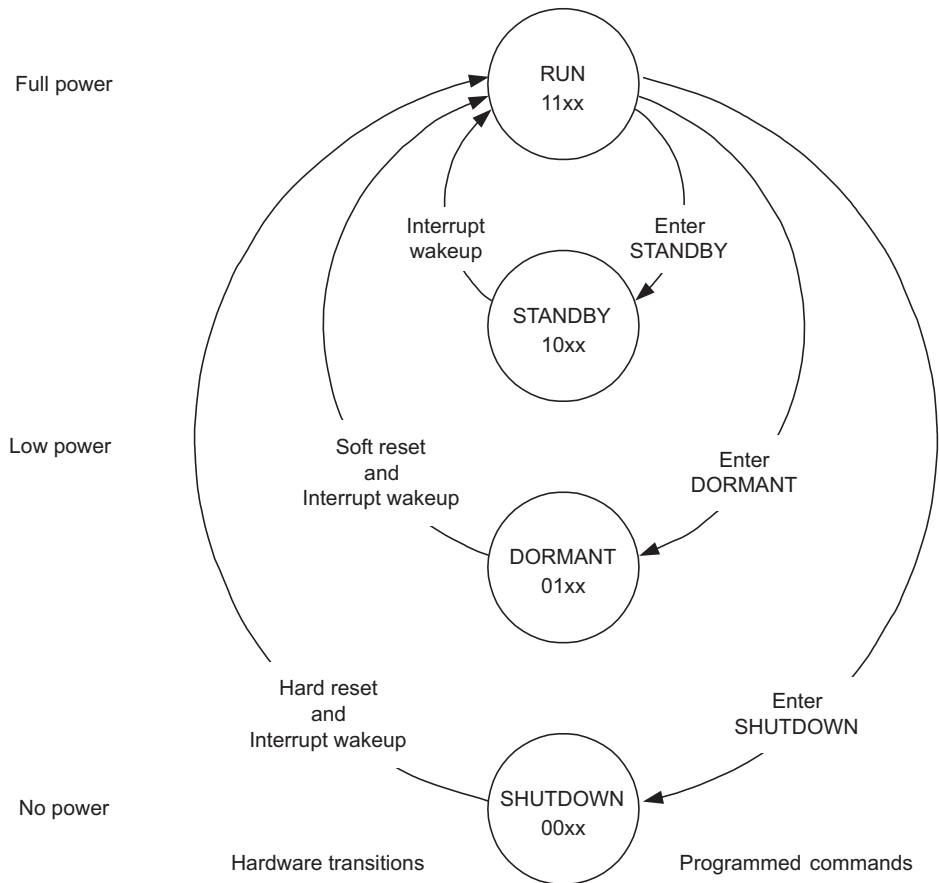


Figure 14-1 Power manager state diagram

Hardware transitions are caused by power supply problems such as low battery reserves or power supply regulation failure.

Table 14-2 summarizes the effects of the ARM10 processor and system power modes.

Table 14-2 Power mode V_{DD} states

ARM10 state	System power state PMTDR/PMRDR[7:4]	Clock	CPU V _{DD}	Cache V _{DD}	Power manager	Description
RUN	11xx	On	On	On	On	Operating speed depends on clock frequency and voltage level of V _{DD} .
STANDBY	10xx	Stop	On	On	On	Processor clocks stopped. Minimal dynamic current.
DORMANT	01xx	None	Off	On	On	Processor core state must be saved in external memory. Leakage current only.
SHUTDOWN	00xx	None	Off	Off	On	Processor core state must be saved in external memory.
		None	Off	Off	On	
		None	Off	Off	On	No leakage from processor or cache.
		None	Off	Off	Off	Processor core state must be saved in external memory. No power.

14.2.1 RUN mode

To determine how to return to RUN mode and resume execution, the ARM10 processor first requests the previous state from the power manager. Table 14-3 shows the restart conditions for each previous state.

Table 14-3 Reentering RUN mode

Previous state	Restart
STANDBY	Processor core state and cache states intact. If interrupt request wakes processor, interrupt vector points to execution entry point. If external debug request wakes processor, execution entry point is the instruction after the one that initiated STANDBY.
DORMANT	Cache states intact. Processor core state must be reloaded. Reset vector points to execution entry point.
SHUTDOWN	Processor core state and cache states must be reloaded. Reset vector points to execution entry point.

14.2.2 STANDBY mode

Do either of the following to put the processor in STANDBY mode:

- program the power manager to the IDLE state
- use the system control coprocessor, CP15, to issue a wait-for-interrupt command.

———— **Note** ————

Before entering STANDBY mode, software must enable wakeup by interrupt request or external debug request.

When exiting STANDBY mode, the processor resumes program execution at one of the following:

- the address pointed to by an interrupt vector if an interrupt request woke the processor
- the address after the instruction that initiated STANDBY mode if an external debug request woke the processor.

Software does not have to check the previous state of the power manager because no hardware state has to be restored.

14.2.3 DORMANT mode

To put the processor in DORMANT mode:

1. Save the processor core state.
2. Use the system control coprocessor, CP15, to issue a request to enter DORMANT mode.

In DORMANT mode, hardware removes power from the processor, leaving the caches powered.

To exit DORMANT mode, do one of the following:

- issue a soft reset
- issue a power-on reset.

A soft reset is the normal way to exit DORMANT mode, as it does not affect the cache state. The processor then vectors to the soft reset routine, which must get the previous state from the power manager so that it can restore the processor and MMU.

14.2.4 SHUTDOWN mode

To put the processor in SHUTDOWN mode:

1. Save the processor core state.
2. Save the cache state.
3. Use the system control coprocessor, CP15, to issue a request to enter SHUTDOWN mode.

To exit SHUTDOWN mode:

1. Issue a power-on reset.
2. Restore the processor core state.
3. Restore the cache state.

14.3 System control coprocessor

Coprocessor CP15 supports power management. CP15 has three registers to transmit and receive data from the power manager. The functionality is similar to the debug communications channel defined in CP14. The three registers are:

- *Power Manager Status Register*
- *Power Manager Receive Data Register* on page 14-9
- *Power Manager Transmit Data Register* on page 14-10.

14.3.1 Power Manager Status Register

The Power Manager Status Register, PMSR, is read-only. It controls synchronized handshaking between the processor and the power manager. Figure 14-2 shows the PMSR bit fields.



Figure 14-2 Power Manager Status Register

Table 14-4 describes the PMSR bit fields.

Table 14-4 PMSR bit fields

Bits	Meaning
[31:28]	Contain a fixed pattern that denotes the power manager architecture version number of the hardware. The code returned for revision 0001 is the first currently defined architecture.
[27:2]	Should Be Zero.
[1]	The W flag is set when the transmit channel is empty and available for a new power manager command. Writing a command to the transmit data register clears W until a handshake acknowledges receipt of the command. Reset sets W to indicate that the power manager transmit data register is ready to accept new data.
[0]	The R flag is set when the power manager receive data register is full and valid data can be read from the channel. Reading the receive data register clears R. Reset sets R to reflect the reason for waking up the processor.

Software can read the status register using the following instruction. Data is returned in register Rd:

```
MRC CP15, 0, Rd, c15, c14, 0
```

Writing to PMSR is Unpredictable.

14.3.2 Power Manager Receive Data Register

The Power Manager Receive Data Register, PMRDR, is read-only. When the R flag in PMSR is set, valid data can be read from PMRDR. An acknowledgement is sent to the power manager to indicate data acceptance. When the R flag in PMSR is cleared, reading PMRDR is Unpredictable. Figure 14-3 shows the bit fields of the PMRDR.

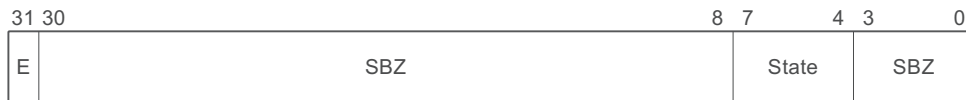


Figure 14-3 Power Manager Receive Data Register

Table 14-5 describes the PMRDR bit fields.

Table 14-5 PMRDR bit fields

Bits	Meaning
[31]	Emulation flag. When exiting a reset sequence, E reflects the last programmed state of the system: 1 = power manager issued a command in emulation mode 0 = power manager issued a command in normal mode.

Table 14-5 PMRDR bit fields

Bits	Meaning
[30:8]	Should Be Zero.
[7:4]	System power state. When exiting a reset sequence, this field reflects the last programmed state of the system: 1111 = TURBO 1110 = NORMAL 110x = SLOW 100x = IDLE 01xx = NAP 0011 = SLEEP 0010 = COMA 0001 = HIBERNATE 0000 = OFF.
[3:0]	Should Be Zero.

Software can read the receive data register using the following instruction. Data is returned in register Rd:

MRC CP15, 0, Rd, c15, c14, 1

Writing to PMRDR is Unpredictable.

14.3.3 Power Manager Transmit Data Register

The Power Manager Transmit Data Register, PMTDR, is write-only. When the W flag in PMSR is set, new data can be written to PMTDR. An acknowledgement following the write is sent to the power manager to indicate that new data is available. Writing to PMTDR clears W. Writing to PMTDR when W is clear is Unpredictable. Figure 14-4 shows the bit fields of the PMTDR.



Figure 14-4 Power Manager Transmit Data Register

Table 14-6 describes the PMTDR bit fields.

Table 14-6 PMTDR bit fields

Bits	Meaning
[31]	1 = power manager issued a command in emulation mode 0 = power manager issued a command in normal mode.
[30:8]	Should Be Zero.
[7:4]	System power state. When exiting a reset sequence, this value reflects the last programmed state of the system: 1111 = TURBO 1110 = NORMAL 110x = SLOW 100x = IDLE 01xx = NAP 0011 = SLEEP 0010 = COMA 0001 = HIBERNATE 0000 = OFF.
[3:0]	Should Be Zero.

Software can write the transmit data register using the following instruction. Data is written using register Rn:

```
MCR CP15, 0, Rn, c15, c14, 1
```

Reading PMTDR is Unpredictable.

14.3.4 Emulation mode

Emulation mode is used in a system to test both software and hardware behavior. Commands are issued in normal mode causing the power manager to change the power mode of the system and the voltages in the core. A typical normal mode command use is to change the mode from RUN to DORMANT to save power. This requires that the power manager tell the regulator controlling the voltage to the processor to lower the voltage from V_{dd} to 0. When a soft reset is issued, the power manager indicates that the voltage to the processor can be raised from 0 to V_{dd} .

To test software and hardware without testing the enabling and disabling of the voltage regulators, issue a command with the emulation bit (E) set. This signals the power manager to translate the command and change to the desired mode. The voltage regulator is never flagged to lower the voltage. When the command is transmitted and received, the power manager issues a soft reset sequence.

Note

The soft reset issued by the power manager occurs during emulation. All other forms of soft reset are done from an external source.

14.3.5 Transmission protocol

When issuing commands to the power manager, a specific sequence must be followed:

1. Verify that both PMTDR and PMRDR are empty by checking that the W flag is set and that the R flag is cleared where appropriate.
2. To transmit, write a command to PMTDR. This clears the W flag. Hardware then performs a handshake with the power manager, waiting for acceptance of the command using a double-ended handshake.
3. When the transmit data handshake is complete, hardware sets the W flag.

When receiving data, software must wait until the R flag is set. When R is set, new valid data is available in PMRDR.

Data transmit code

When data has to be transmitted to the power manager, software must always perform the code sequence shown below. The command is sent using register R1, while R0 reflects the status register contents:

tx_command:

```
MRC CP15, 0, R0, c15, c14, 0      ; check for outstanding commands
TST R0, #W_flag                   ; W flag clear indicates active command
BNE tx_command                    ; if command active, loop again
MCR CP15, 0, R1, c15, c14, 1      ; write new command to controller
```

Note

The W flag is polled until it is set. When W is set, the command can be sent to the power manager.

Data receive code

To wait until data has been received in the receive data register, software must always perform the code sequence shown below. The command is received into register R1, while R0 is used to reflect the status register contents:

rx_status:

```
MRC CP15, 0, R0, c15, c14, 0    ; check for incoming data
TST R0, #R_flag                 ; R flag clear indicates no data
BNE rx_status                   ; if no data, loop again
MRC CP15, 0, R0, c15, c14, 1    ; read in 'previous-state'
```

———— **Note** ————

The R flag is polled until it is cleared. When R is cleared, the command can be read.

14.4 Programming examples

This section contains examples of how to change the processor power mode.

14.4.1 RUN to STANDBY

This example changes the processor mode from RUN to STANDBY:

tx_command:

```
MRC CP15, 0, R0, c15, c14, 0    ; check for outstanding commands
TST R0, #W_flag                ; W flag clear indicates active command
BNE tx_command                  ; if command active, loop again
MOV R1, #PM_IDLE SHL 4          ; program IDLE state into 7:4, no emulation
MCR CP15, 0, R1, c15, c14, 1    ; write new command to controller
```

14.4.2 RUN to DORMANT

This example changes the processor mode from RUN to DORMANT:

;save all ARM1020E macrocell state here

tx_command:

```
MRC CP15, 0, R0, c15, c14, 0    ; check for outstanding commands
TST R0, #W_flag                ; W flag clear indicates active command
BNE tx_command                  ; if command active, loop again
MOV R1, #PM_NAP SHL 4           ; program NAP state into 7:4, no emulation
MCR CP15, 0, R1, c15, c14, 1    ; write new command to controller
B .                             ; branch to self to freeze core on this
                                ; instruction
```

14.4.3 RUN to SHUTDOWN

This example changes the processor mode from RUN to SHUTDOWN:

;no ARM1020E macrocell state needs to be saved since entering SHUTDOWN

tx_command:

```
MRC CP15, 0, R0, c15, c14, 0    ; check for outstanding commands
TST R0, #W_flag                ; W flag clear indicates active command
BNE tx_command                  ; if command active, loop again
MOV R1, #PM_SHUTDOWN SHL 4      ; put SHUTDOWN state into 7:4, no emulation
MCR CP15, 0, R1, c15, c14, 1    ; write new command to controller
B .                             ; branch to self to freeze core on this
                                ; instruction
```


14.4.4 Reset recovery

This example detects the previous state of the power manager before a power-on reset or soft reset:

```
B reset
;insert other code here
reset
MRC CP15, 0, R0, c15, c14, 0      ; check for incoming data
TST R0, #R_flag                  ; R flag clear indicates no data
BNE reset                        ; if no data, loop again
MRC CP15, 0, R0, c15, c14, 1     ; read in 'previous-state'
TST R0, #0xC0                   ; check to see if 'previous-state' RUN
BEQ last_state_run
TST R0, #0x80                   ; check to see if 'previous-state' STANDBY
BEQ last_state_standby
TST R0, #0x40                   ; check to see if 'previous-state' DORMANT
BEQ last_state_dormant
;execute default power-on reset code here
```

14.5 Power manager interface

Table 14-7 defines the interface between the power manager and the ARM10 processor.

Table 14-7 Power manager/processor interface signals

Signal	Direction	Description
PMEXISTS	To processor	Power manager active-HIGH signal to processor. If power manager not attached to processor, PMEXISTS must be at logic 0.
PMTXREQ	From processor	CPU request for power manager state change. PMTXREQ and PMTXACK provide a double-ended handshake in transmissions to the power manager.
PMTXACK	To processor	Power manager asserts PMTXACK to acknowledge processor state change on PMTX[3:0] .
PMTX[3:0]	From processor	CPU state change data.
PMTXEMUL	From processor	CPU state change request in emulation mode. Request for power manager to leave the voltage regulators unchanged.
PMRXREQ	From processor	CPU request for previous state of power manager. PMRXREQ and PMRXACK provide a double-ended handshake during power manager reception.
PMRXACK	To processor	Power manager acknowledgement of PMRXREQ . Signals valid data on PMRX[3:0] .
PMRX[3:0]	To processor	Power manager previous state data.
PMRXEMUL	To processor	Power manager previous state of emulation.
SFRESETN	To processor	Power manager active-LOW soft reset indicator.
HRESETN	To processor	Power manager active-LOW power-on or AHB bus reset.

14.6 Timing

The timing diagrams in this section illustrate the following:

- *ARM10 processor transmit*
- *ARM10 processor transmit with emulation*
- *ARM10 processor previous-state request on page 14-18*
- *ARM10 processor previous-state request with emulation on page 14-18*
- *ARM10 processor hard reset on page 14-19*
- *ARM10 processor soft reset from powerdown timing on page 14-19.*

14.6.1 ARM10 processor transmit

In Figure 14-5 the processor sends **PMTXREQ** to the power manager. The power manager acknowledges with **PMTXACK** and puts the state entered on **PMRX[3:0]**.

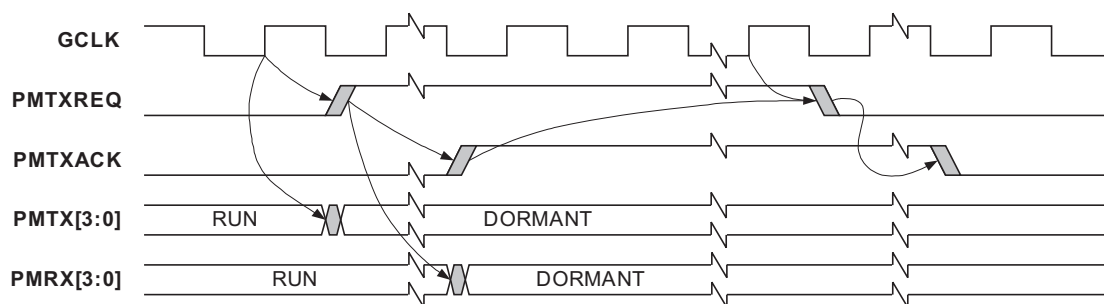


Figure 14-5 CPU transmit request timing

14.6.2 ARM10 processor transmit with emulation

In Figure 14-6 on page 14-18 the emulation bit is set. The processor sends **PMRXEQ** to the power manager. The power manager then acknowledges with **PMRXACK** and issues the requested state on **PMRX[3:0]**. In this case, the voltage regulators do not change.

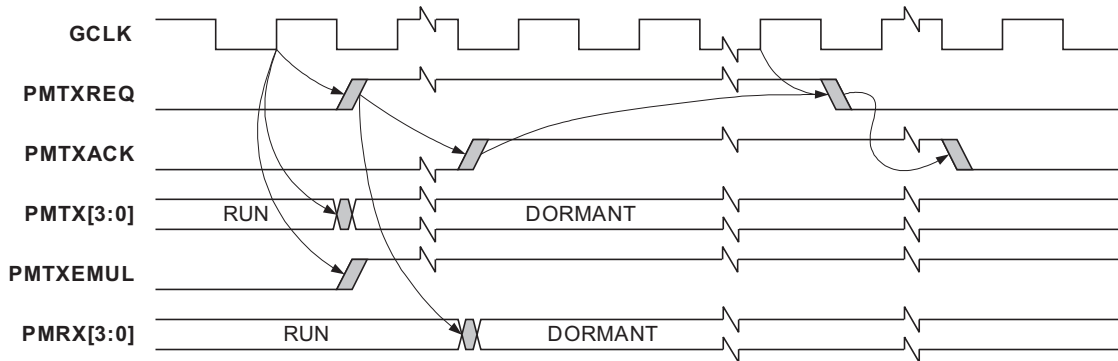


Figure 14-6 CPU transmit request timing with emulation bit set

14.6.3 ARM10 processor previous-state request

In Figure 14-7 the processor sends **PMRXEQ** to the power manager. The power manager then issues an acknowledge with the previous state on **PMRX[3:0]**.

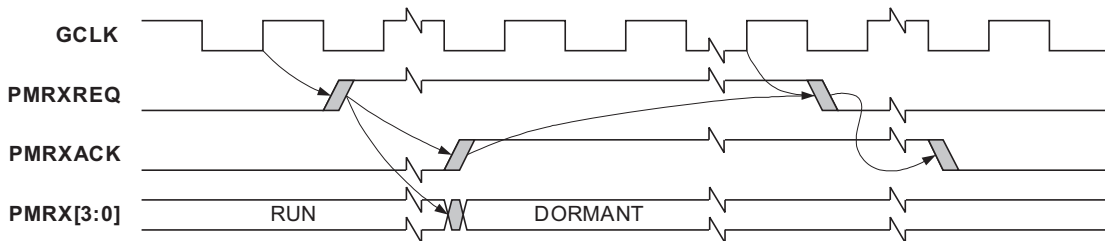


Figure 14-7 CPU previous state request timing

14.6.4 ARM10 processor previous-state request with emulation

In Figure 14-8 on page 14-19 the processor sends **PMRXEQ** to the power manager. The power manager issues an acknowledgment with the previous state on **PMRX[3:0]**. **PMRXEMUL** indicates that the previous state was in emulation mode.

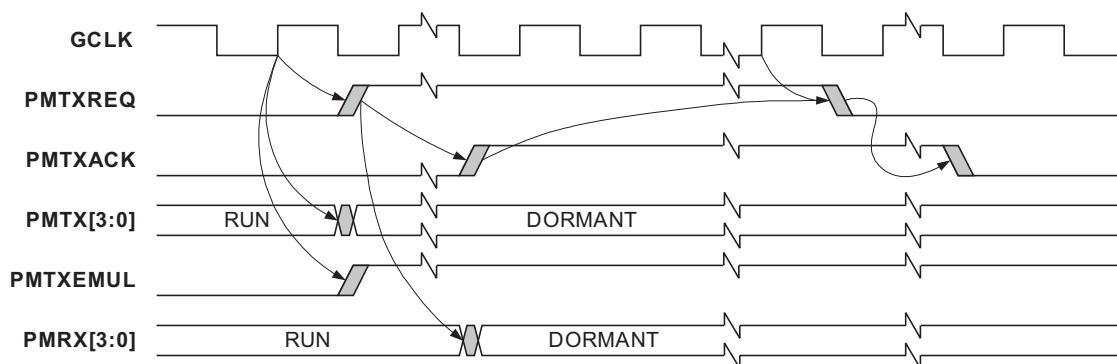


Figure 14-8 CPU previous state request timing with emulation bit set

14.6.5 ARM10 processor hard reset

Figure 14-9 shows that both hard reset, **HRESETN**, and soft reset, **SFRESETN**, must be issued to the processor in the same cycle.

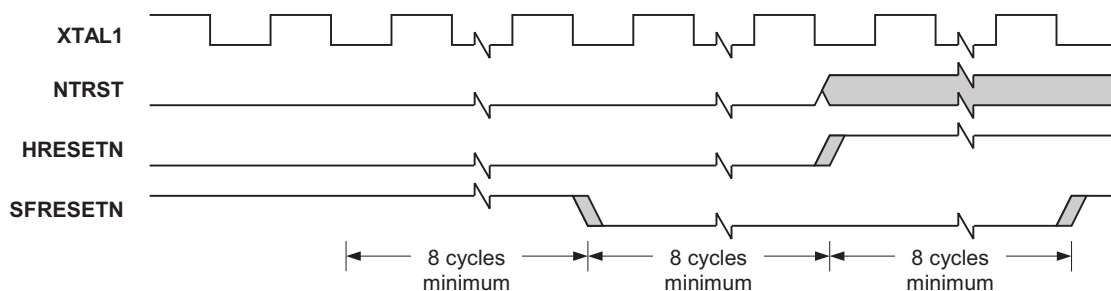


Figure 14-9 Hard reset timing

To guarantee that the processor properly resets all states, hard reset is asserted for a minimum of eight cycles, and soft reset must be extended a minimum of eight cycles beyond the point where **HRESETN** is deasserted.

14.6.6 ARM10 processor soft reset from powerdown timing

Figure 14-10 on page 14-20 shows how a soft reset can be issued following entry into DORMANT mode.

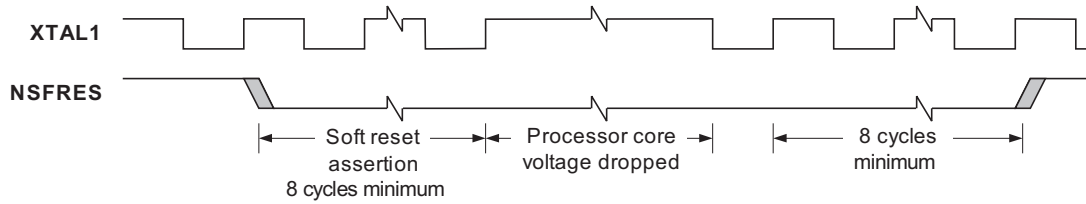


Figure 14-10 Soft reset from power-down timing

When the processor enters DORMANT and receives an acknowledgement from the power manager, the voltage can be removed from the processor, and the voltage to the processor caches can be lowered to the minimum value that retains state.

The soft reset signal, **SFRESETN**, must stay LOW when the processor voltage is taken away to ensure proper behavior when the processor voltage is returned.

When the processor and cache voltages are raised to the operational value, **SFRESETN** must be asserted at least eight more cycles to guarantee a proper exit from soft reset.

14.7 Software example code sequences

The precise definition of state to be saved and reloaded in a system is implementation-defined. The example routines in this section show the basic requirements and give a starting point for implementation.

14.7.1 Save_L0_state code sequence

```

        AREA |PowerDown|, CODE, READONLY
        KEEP
        EXPORT PwrMgt_Save_L0_State
PwrMgt_Save_L0_State
    ; On Entry: Processor must be in a privileged mode. R0 points to start of
    ; the data block in memory. This code disables virtual memory, so must be
    ; executed from a virtual address that is mapped to the same physical
    ; address.

    ; first save all the integer registers, CPSR & SPSRs

    STMIA R0!, {R1-R7}          ; save unbanked registers
    MRS R2, CPSR
    STMIA R0!, {R2}            ; save CPSR
    STMIA R0, {R8 - R14}^      ; save user mode banked registers
    ADD R0, R0, #28             ; increment base register

    BIC R3, R2, #0x1f          ; clear the mode bits from the CPSR value

    ; now roll through each of the privileged modes and save banked registers

    ORR R4, R3, #0x13          ; SVC mode
    MSR CPSR_cf, r4
    MRS R5, SPSR
    STMIA R0!, {R5, R13, R14}  ; save SPSR and banked registers

    ORR R4, R3, #0x1b          ; UNDEF mode
    MSR CPSR_c, r4
    MRS R5, SPSR
    STMIA R0!, {R5, R13, R14}  ; save SPSR and banked registers

    ORR R4, R3, #0x17          ; ABORT mode
    MSR CPSR_c, r4
    MRS R5, SPSR
    STMIA R0!, {R5, R13, R14}  ; save SPSR and banked registers

    ORR R4, R3, #0x12          ; IRQ mode
    MSR CPSR_c, r4
    MRS R5, SPSR
    STMIA R0!, {R5, R13, R14}  ; save SPSR and banked registers

```

```

ORR R4, R3, #0x11          ; FIQ mode
MSR CPSR_c, r4
MRS R5, SPSR
STMIA R0!, {R5, R8 - R14} ; save SPSR and banked registers

MSR CPSR_c, R2            ; and return to the original mode

; now do the CP15 registers

MRC p15, 0, R1, c1, c0, 0 ; Control register
MRC p15, 0, R2, c2, c0, 0 ; Translation Table Base
MRC p15, 0, R3, c3, c0, 0 ; Domain Access Control
MRC p15, 0, R5, c5, c0, 0 ; FSR
MRC p15, 0, R6, c6, c0, 0 ; FAR

STMIA R0!, {R1 - R3, R5, R6}

MOV R7, #0                ; dummy data
MCR p15, 0, R7, c7, c10, 4 ; Drain the write buffer

; Insert code here to power down ARM1020E macrocell
B .

END

```

14.7.2 Reload_L0_state code sequence

```

AREA |PowerDown|, CODE, READONLY
KEEP
EXPORT PwrMgt_Reload_L0_State
PwrMgt_Reload_L0_State
; On entry, the processor must be in a privileged mode. R0 points to start
; of the data block in memory. This code disables virtual memory, so must be
; executed from a virtual address that is mapped to the same physical
; address

; first clear the TLBs ready to turn on virtual memory

ADD R0, R0, #0xa0         ; size of the data block

MOV R7, #0                ; dummy data
MCR p15, 0, R7, c8, c7, 0 ; Invalidate ITLB/DTLB

; now do the CP15 registers

LDMDB R0!, {R1 - R3, R5, R6}

MCR p15, 0, R2, c2, c0, 0 ; Translation Table Base
MCR p15, 0, R3, c3, c0, 0 ; Domain Access Control

```



```

MCR p15, 0, R5, c5, c0, 0 ; FSR
MCR p15, 0, R6, c6, c0, 0 ; FAR
MCR p15, 0, R1, c1, c0, 0 ; Control register

MRS R2, CPSR
BIC R3, R2, #0x1f; clear the mode bits from the CPSR value

; now roll through each of the privileged modes and restore banked registers

ORR R4, R3, #0x11 ; FIQ mode
MSR CPSR_c, r4
LDMDB R0!, {R5, R8 - R14} ; restore SPSR and banked registers
MSR SPSR_cxsf, R5

ORR R4, R3, #0x12 ; IRQ mode
MSR CPSR_c, r4
LDMDB R0!, {R5, R13, R14} ; restore SPSR and banked registers
MSR SPSR_cxsf, R5

ORR R4, R3, #0x17; ABORT mode
MSR CPSR_c, r4
LDMDB R0!, {R5, R13, R14} ; restore SPSR and banked registers
MSR SPSR_cxsf, R5

ORR R4, R3, #0x1b ; UNDEF mode
MSR CPSR_c, r4
LDMDB R0!, {R5, R13, R14} ; restore SPSR and banked registers
MSR SPSR_cxsf, R5

ORR R4, R3, #0x13; SVC mode
MSR CPSR_cf, r4
LDMDB R0!, {R5, R13, R14} ; restore SPSR and banked registers
MSR SPSR_cxsf, R5

; now restore all the integer registers, CPSR & SPSRs

LDMDB R0, {R8 - R14}^ ; restore user mode banked registers
SUB R0, R0, #28 ; decrement base register

LDMDB R0!, {R2} ; restore CPSR
MSR CPSR_cxsf, R2
LDMDB R0!, {R1-R7} ; restore unbanked registers
END

```


Chapter 15

Clock Generator

This chapter describes the operation of a *Phase-Locked Loop* (PLL) using the clock generator. This chapter contains the following sections:

- *Features* on page 15-2
- *About the clock generator* on page 15-3
- *Interface description* on page 15-6
- *Output clock behavior* on page 15-9
- *PLL Configuration Register* on page 15-11
- *Output clock balancing* on page 15-13.

15.1 Features

The clock generator synthesizes two programmable clocks. It contains analog circuitry with enough flexibility to cover a range of applications while placing minimum restrictions on the remainder of the test chip.

The key features include:

- two synchronized, frequency-programmable clock outputs
- internal loop filter
- output duty cycle from 48% to 52%
- power-down and *Voltage-Controlled Oscillator (VCO)* bypass modes
- partner-specific mode support
- integrated crystal oscillator option
- testable design.

On and during reset, the **XTAL1** reference clock drives **HCLK** and **GCLK** directly, bypassing the VCO. After reset, when the PLL is configured and has achieved lock, the control interface can turn off VCO bypass mode. The output clocks must switch seamlessly from **XTAL1** to the VCO output without exceeding the frequency or minimum phase time of the faster clock.

Whether **HCLK** and **GCLK** are derived from the VCO or from a 50% duty cycle reference clock at less than 100MHz, the duty cycle degradation must be minimal. Phase lock with external signals and zero insertion delay are not required. The PLL feedback path, including the loop filter, is completely internal to the clock generator.

Dedicated V_{DDA} and V_{SSA} pins supply power to both the analog and digital portions of the clock generator. The clock generator must have its own power supply so that it does not affect power measurements made on the test chip.

The following equations show the derivations of **GCLK** and **HCLK**:

$$\text{GCLK} = \text{XTAL1} \times \text{MDIV}[7:0] + 1$$

$$\text{HCLK} = \text{XTAL1} \times \frac{\text{MDIV}[7:0] + 1}{\text{HDIV}[3:0] + 1}$$

Table 15-1 shows **GCLK** and **HCLK** frequencies with **XTAL1** at 20MHz.

Table 15-1 GCLK/HCLK frequencies with XTAL1 = 20MHz

GCLK/HCLK	MDIV[7:0]				
	0	1	2	3	...
0	20/20	40/40	60/60	80/80	...
1	20/10	40/20	60/30	80/40	...
2	20/6.67	40/13.3	60/20	80/26.7	...
...
14	20/1.33	40/2.67	60/4	80/5.33	...
15	20/1.25	40/2.5	60/3.75	80/5	...

You can program values of **MDIV[7:0]** with a given **XTAL1** that would cause the clock generator to operate outside of the VCO functional operating range. You must supply an addendum that states the restrictions placed on **MDIV[7:0]** for various **XTAL1** inputs. Here are example addendum restrictions on **MDIV[7:0]** for a VCO_{max} of 800MHz:

$$VCO_{max} = 800\text{MHz} = \text{XTAL1} \times \text{MDIV}[7:0]$$

- If **XTAL1** = 5MHz, then the maximum value for **MDIV[7:0]** is 159.
- If **XTAL1** = 100MHz, then the maximum value for **MDIV[7:0]** is 7.

15.3 Interface description

This section describes the clock generator input and output signals.

XTAL1 This is the reference clock input. During reset it drives the two output clocks, **HCLK** and **GCLK**. If an integrated crystal oscillator is used, it is one of the connections to the crystal. If an integrated crystal oscillator is not used, an external oscillator drives **XTAL1**.

NPORES This is the power-on reset input. During reset, **NPORES** is driven LOW for multiple **XTAL1** cycles and ensures that **XTAL1** drives **HCLK** and **GCLK** during this time.

XTAL2EN This output enables the external crystal oscillator. If the crystal oscillator is internal, then this is its output.

CLKTESTCTL[3:0]

As shown in Table 15-2, these test control inputs select clock generator internal clocks for viewing on **CLKTESTOUT**.

———— **Note** —————

The **CLKTESTCTL[3:0]** pins are not internally synchronized before use, meaning that entering a test mode might cause a VMUX, GMUX, or HMUX glitch.

Table 15-2 Test mode programming

CLKTESTCTL[3:0]	Test mode
0000	Normal mode of operation. CLKTESTOUT = 0. Crystal oscillator enabled.
0001	XTAL1 drives M divider. CLKTESTOUT = MCLK . Isolates design faults in M divider circuit.
0010	I_{DDQ} test mode. All circuits are silent. CLKTESTOUT = 0, HCLK = GCLK = XTAL1 . Apply patterns in VCO bypass mode. Then switch to I_{DDQ} test mode.
0011	VCO bypass mode. XTAL drives HCLK and GCLK directly.
01xx	CLKTESTCTL[1] drives GCLK . CLKTESTCTL[0] drives HCLK . Bypasses clock generator due to extreme failure.

Table 15-2 Test mode programming (continued)

CLKTESTCTL[3:0]	Test mode
1000	CLKTESTOUT is the GCLK output. Tests for defects in PLL, H divider, and crystal oscillator.
1001	CLKTESTOUT is the HCLK output. Tests for defects in PLL, H divider, and crystal oscillator.
1010	CLKTESTOUT is the VCO output. Tests for defects in PLL, H divider, and crystal oscillator.
1011	CLKTESTOUT is the crystal oscillator output. Tests for defects in PLL, H divider, and crystal oscillator.
110x	Partner-specific test modes.
111x	RESERVED.

———— **Note** —————

The rising edge of **GCLK** must be synchronous with the rising edge of **HCLK**. There must be zero delay between **GCLK** and **HCLK** for any given clock input.

CLKTESTOUT

This clock test output is for viewing **GCLK**, **HCLK**, **MCLK**, **XTAL1**, or **VCO**.

PCONFIGIN[5:0]

These are configuration inputs for PLL-specific control signals.

PCONFIGIN[5:0] are cleared by reset and can be programmed with a CP15 instruction.

PCONFIGOUT[1:0]

These are configuration outputs for PLL-specific control signals. If the PLL has a lock-detect signal, it must be tied to **PCONFIGOUT[0]**. Any other PLL outputs must use **PCONFIGOUT[1]**.

POWERDN This is the powerdown input. When **POWERDN** is HIGH, the PLL shuts down and draws the minimum leakage current. In a typical operating configuration, the VCO must first be bypassed so the ARM10 processor can continue to run from **XTAL1**. **POWERDN** is set by reset and can be programmed with a CP15 instruction.

BYPASS[1:0]

These are inputs that control selection of the VMUX, GMUX, and HMUX multiplexors. **BYPASS[1:0]** are set by reset and can be programmed with a CP15 instruction.

MDIV[7:0]

These inputs select the PLL multiplier. The value programmed is **MDIV[7:0]** + 1. **MDIV[7:0]** are cleared by reset and can be programmed with a CP15 instruction.

HDIV[3:0]

These inputs select the H divider. **HDIV[3:0]** are set by reset and can be programmed with a CP15 instruction.

GCLK

This output is the the primary clock of the ARM10 processor. The clock generator must be able to drive **GCLK** at maximum frequency under all process conditions. During reset, **GCLK** must be driven by the **XTAL1** input.

HCLK

This output is the primary AHB clock and is also an input to the ARM10 processor. The clock generator must be able to drive **HCLK** at maximum frequency under all process conditions. During reset, **HCLK** must be driven by the **XTAL1** input.

15.4 Output clock behavior

The clock generator output clocks, **HCLK** and **GCLK**, are defined by the inputs **HDIV[3:0]**, **BYPASS[1:0]**, and **POWERDN**. It is a strict requirement that the output clocks are driven by **XTAL1** during reset so that reset is propagated throughout the ARM10 processor. It is also a requirement that both **HCLK** and **GCLK** have no glitches, have synchronous rising edges, and have approximately 50% duty cycles.

When multiplexing from one input clock to the other, the resultant output clock must not have a pulse smaller than either of the input clocks. An output clock pulse smaller than either of the input clocks is a glitch. Clock switching must be done so that **HCLK** and **GCLK** remain glitch-free.

HCLK and **GCLK** must have synchronous rising edges. When reprogramming the H divider, take care to ensure that:

- the **HCLK** and **GCLK** rising edges are synchronous
- the resultant clocks have approximately 50% duty cycles
- no glitches occur.

Table 15-3 on page 15-10 shows the behavior of **HCLK** and **GCLK**.

———— **Note** —————

The inputs to the table are the outputs of the synchronizers from the CP15 coprocessor. It is strongly recommended that all inputs from the ARM10 processor go through a full synchronizer before being used in any logic in the clock generator.

Table 15-3 does not account for lock detection. PLL lock does not factor into the output **GCLK** and **HCLK** multiplexor selection logic.

Table 15-3 GCLK and HCLK behavior

HDIV[3:0]	BYPASS[1:0]	POWERDN	Internal VCO	HCLK	GCLK
> 0	00	0	Active	H divider output. VCO output drives H divider.	VCO output
	00	1	0		
	01	0	Active	H divider output. XTAL1 drives H divider.	
	01	1	0		XTAL1
	1x	0	Active		
	1x	1	0	XTAL1	
= 0	00	0	Active	VCO output	VCO output
	00	1	0		
	01	0	Active		
	01	1	0	XTAL1	XTAL1
	1x	0	Active		
	1x	1	0		

The only time that the behavior in Table 15-3 does not apply is when the clock generator is in one of the test modes defined in Table 15-2 on page 15-6.

Most of the clock generator inputs and outputs come from a CP15 register within the ARM10 processor. This register controls both dividers, PLL power-down enable, VMUX, HMUX, GMUX, and special partner-specific configuration inputs. The register can be read and written under software control in supervisor mode only.

15.5 PLL Configuration Register

Figure 15-2 shows the PLL configuration register. All signals within this register are all active HIGH. These instructions are defined to operate in supervisor mode only. Any other mode of operation in the ARM10 processor bounces the instruction, causing an exception to be taken.

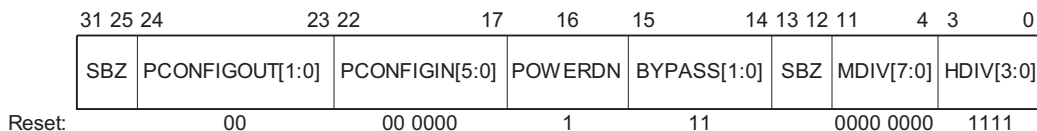


Figure 15-2 PLL Configuration Register

The instructions used to access the CP15 PLL configuration register are:

- write: MCR p15,0,Rd,c15,c12,0
- read: MRC p15,0,Rd,c15,c12,0.

The SBZ fields in must always be written as zeros.

15.5.1 Programming the PLL Configuration Register

Examples for reprogramming the CP15 PLL configuration register appear in the following sections:

- *After reset*
- *Entering powerdown state* on page 15-12
- *With no lock hardware* on page 15-12.

The examples are based on the following frequencies:

- **XTAL1** frequency = 20MHz
- **GCLK** frequency = 60MHz
- **HCLK** frequency = 30MHz.

After reset

Program the CP15 PLL configuration register, assuming that a lock indicator exists.

```

LDR r0, = 0x0000C021
MCR p15,0,r0,c15,c12,0 ; write new contents
Loop MRC p15,0,r1,c15,c12,0 ; reread the contents
      TST r1, #0x00800000 ; check to see if lock bit is set
      BNE Loop ; if lock bit not set, recheck

```

Entering powerdown state

When the clock generator is programming the **POWERDN** bit so that the PLL VCO is silent, software must simply read the CP15 PLL configuration register, set the **POWERDN** bit to one, and rewrite the CP15 PLL configuration register:

```
MRC p15,0,r1,c15,c12,0; ; read state of CP15 register
ORR r1, r1, #0x00010000 ; set the POWERDN bit
MCR p15,0,r1,c15,c12,0 ; reprogramming with POWERDN bit set
```

With no lock hardware

When the PLL is programmed and no lock hardware exists, the software must calculate how much time must be allocate to waiting. This is done by assuming a fixed value for the lock time, 150 μ s, and calculating the wait as a function of the input frequency,

XTAL1:

```

MOV r0, #150           ; lock time in us
MOV r1, #20            ; XTAL1 in MHz
MUL r2, r0, r1         ; counter wait time
LDR r3, = 0x0000C021  ; value to write
MCR p15,0,r3,c15,c12,0 ; write new CP15 PLL register contents
Loop SUBS r2, r2, #0x1 ; decrement wait counter
    BNE Loop           ; if count not zero, wait is not done so loop
```

15.6 Output clock balancing

HCLKOUT is a tap from the internal **HCLK** balanced clock tree. You can use it for:

- measuring the clock tree insertion delay on **HCLK**
- synchronizing the external **HCLK** in a memory system.

Figure 15-3 shows the requirements for balancing **HCLK** in a system.

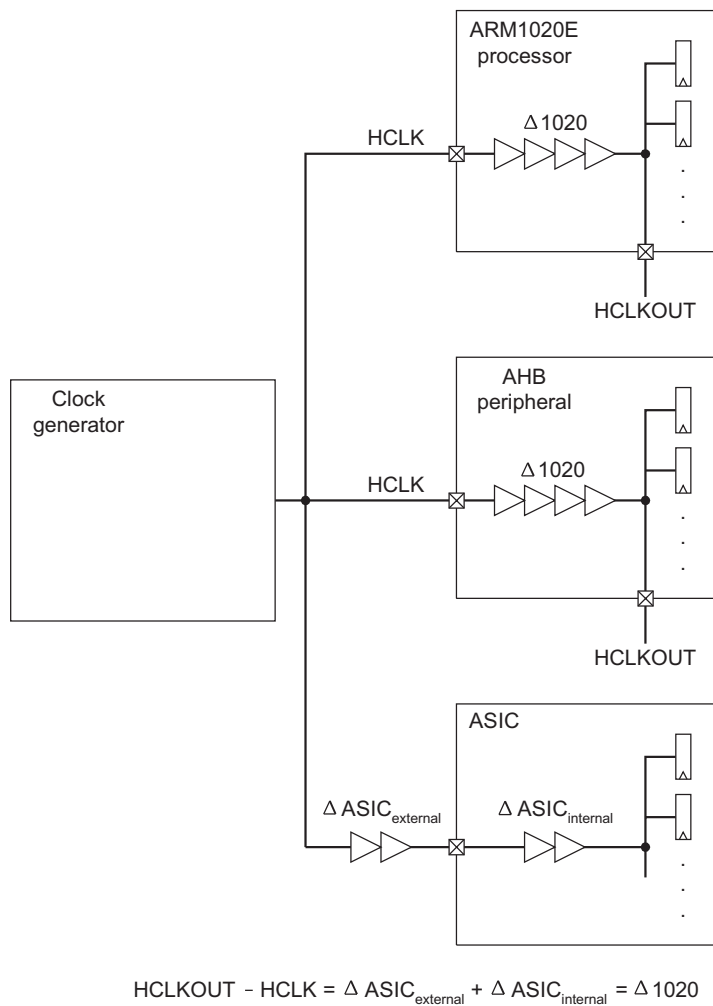


Figure 15-3 Output clock balancing

Appendix A

Signal Descriptions

This appendix describes the ARM10 processor signals. It contains the following sections:

- *Global control signals* on page A-2
- *AHB signals in normal mode* on page A-3
- *PLL signals* on page A-6
- *Power manager signals* on page A-7
- *JTAG and TAP controller signals* on page A-8
- *Debug signals* on page A-9
- *Coprocessor signals* on page A-10
- *Design for test signals* on page A-12
- *ETM signals* on page A-14.

A.1 Global control signals

Table A-1 shows all the processor input signals used to set clocks, memory configuration, vector table location, and external interrupt features.

Table A-1 Global control signals

Signal	I/O	Description
BIGENDINIT	I	Configures processor to treat memory bytes as big-endian or little-endian: 1 = big-endian format 0 = little-endian format.
GCLK	I	Global clock. Drives processor. Can be stopped in either clock phase.
HIVECSINIT	I	Configures the vector table location coming out of reset: 1 = 0xFFFF0000 0 = 0x00000000.
ISYNC	I	Set ISYNC HIGH if NFIQ and NIRQ are synchronized to the processor clock. Set it LOW for asynchronous interrupts.
NFIQ	I	Fast interrupt request signal. Active-LOW.
NIRQ	I	Interrupt request signal. Active-LOW.

A.2 AHB signals in normal mode

Table A-2 list the AHB signals divided by function.

Table A-2 AHB signals

Signal	I/O	Description
HADDRI[31:0]	O	IBIU address bus.
HADDRD[31:0]	O	DBIU address bus.
HBURSTI[2:0]	O	IBIU burst transfer type: 000 = single transfer 010 = four-beat wrapping burst.
HBURSTD[2:0]	O	DBIU burst transfer type: 000 = single transfer 010 = four-beat wrapping burst 011 = four-beat incrementing burst.
HCLK	I	Clock that times all bus transfers. All signals are related to the rising edge of HCLK .
HCLKOUT	O	External tap from internal HCLK balanced clock tree.
HPROTI[3:0]	O	IBIU protection control. Transfers are always opcode fetches: xxx0 = opcode fetch xxx1 = data access xx0x = user access xx1x = privileged access x0xx = not bufferable x1xx = bufferable 0xxx = not cachable 1xxx = cachable.
HPROTD[3:0]	O	DBIU protection control. Transfers are always data accesses: xxx0 = opcode fetch xxx1 = data access xx0x = user access xx1x = privileged access x0xx = not bufferable x1xx = bufferable 0xxx = not cachable 1xxx = cachable.
HRDATAI[63:0]	I	Read IBIU data bus. Transfers data and instructions from bus slaves to instruction-side bus master in read operations.

Table A-2 AHB signals (continued)

Signal	I/O	Description
HRDATAD[63:0]	I	Read DBIU data bus. Transfers data from bus slaves to data-side bus master in read operations.
HREADYI	I	Slave ready. HIGH means transfer finished. Can be driven LOW to extend transfer.
HREADYD	I	Slave ready. HIGH means transfer finished. Can be driven LOW to extend transfer.
HRESETN	I	Resets system and bus. It is the only active-LOW AHB signal.
HRESPI[1:0]	I	Slave response to IBIU. Reflects status of transfer: 00 = OKAY 01 = ERROR 10 = RETRY 11 = SPLIT.
HRESPD[1:0]	I	Slave response to DBIU. Reflects status of transfer: 00 = OKAY 01 = ERROR 10 = RETRY 11 = SPLIT.
HSIZEI[2:0]	O	Size of IBIU transfer: 000 = byte (8 bits) 001 = halfword (16 bits) 010 = word (32 bits) 011 = doubleword (64 bits) 100 = 4 words (128 bits) 101 = 8 words (256 bits) 110 = 16 words (512 bits) 111 = 32 words (1024 bits).
HSIZED[2:0]	O	Size of DBIU transfer: 000 = byte (8 bits) 001 = halfword (16 bits) 010 = word (32 bits) 011 = doubleword (64 bits) 100 = 4 words (128 bits) 101 = 8 words (256 bits) 110 = 16 words (512 bits) 111 = 32 words (1024 bits).

Table A-2 AHB signals (continued)

Signal	I/O	Description
HTRANSI[1:0]	O	Selects type of IBIU transfer: 00 = IDLE 01 = BUSY (This signal is not used) 10 = NONSEQUENTIAL 11 = SEQUENTIAL.
HTRANSD[1:0]	O	Reflects type of DBIU transfer: 00 = IDLE 01 = BUSY (This signal is not used) 10 = NONSEQUENTIAL 11 = SEQUENTIAL.
HWDATAD[63:0]	O	DBIU write data bus. Transfers data from master to slaves in write operations.
HWRITEI	O	IBIU transfer direction. HIGH means write transfer. LOW means read transfer.
HWRITED	O	DBIU transfer direction. HIGH means write transfer. LOW means read transfer.

Table A-3 lists arbiter signals.

Table A-3 Arbiter signals

Name	I/O	Description
HBUSREQD	O	Request line from DBIU.
HBUSREQI	O	Request line from IBIU.
HGRANTD	I	AHB mastership granted to DBIU.
HGRANTI	I	AHB mastership granted to IBIU.
HLOCKD	O	Indicates sequence of locked DBIU transfers in SWP operations.
HLOCKI	O	For AMBA compliance. Never asserted.

A.3 PLL signals

The signals described in Table A-4 are for test chip use only and must not be used for other designs.

Table A-4 PLL signals

Name	I/O	Description
BYPASS[1:0]	O	PLL bypass enable. Do not connect.
HDIV[3:0]	O	PLL HCLK divider. Do not connect.
MDIV[7:0]	O	PLL feedback divider. Do not connect.
PCONFIGOUT[1:0]	I	PLL output lines. Tie off LOW.
PCONFIGIN [5:0]	O	PLL configuration. Do not connect.
POWERDN	O	PLL powerdown. Do not connect.

A.4 Power manager signals

Table A-5 lists the power manager signals.

Table A-5 Power manager/processor interface signals

Signal	I/O	Description
PMEXISTS	I	Power manager active-HIGH signal to processor. If power manager not attached to processor, PMEXISTS must be at logic 0.
PMTXREQ	O	CPU request for power manager state change. PMTXREQ and PMTXACK provide a double-ended handshake in transmissions to the power manager.
PMTXACK	I	Power manager asserts PMTXACK to acknowledge processor state change on PMTX[3:0] .
PMTX[3:0]	O	CPU state change data.
PMTXEMUL	O	CPU state change request in emulation mode. Request for power manager to leave the voltage regulators unchanged.
PMRXREQ	O	CPU request for previous state of power manager. PMRXREQ and PMRXACK provide a double-ended handshake during power manager reception.
PMRXACK	I	Power manager acknowledgement of PMRXREQ . Signals valid data on PMRX[3:0] .
PMRX[3:0]	I	Power manager previous state data.
PMRXEMUL	I	Power manager previous state of emulation.
SFRESETN	I	Power manager active-LOW soft reset indicator.
HRESETN	I	Power manager active-LOW power-on or AHB bus reset.

A.5 JTAG and TAP controller signals

Table A-6 lists the JTAG signals.

Table A-6 JTAG signals

Name	I/O	Description
NTRST	I	Active-LOW test reset signal for boundary scan logic. LOW in normal operation.
TCK	I	Test (JTAG) clock.
TDI	I	JTAG test data input.
TDO	O	JTAG test data output.
TMS	I	Test mode select. Selects state of TAP controller state machine.

Table A-7 lists the TAP controller signals.

Table A-7 TAP controller signals

Name	I/O	Description
CLOCKDR	O	External boundary scan chain clock.
IR[3:0]	O	JTAG instruction register. Reflect current instruction in TAP controller instruction register.
NRSTOVR	O	Output TAP reset override. Used in boundary scan test. Active when scan chain 3 selected and IR = EXTEST, CLAMP or HIGHZ.
NTDOEN	O	Tristate enable for TDO output pin.
SCREG[4:0]	O	Scan chain selection register.
SDOUTBS	I	External or boundary scan out.
SHIFTR	O	Combinational decode of TAP state machine used as multiplexed external scan cell clock.
TAPID[31:0]	I	TAP ID number.
TAPSM[3:0]	O	Reflect current state of TAP controller state machine. Change on rising edge of TCK .
UPDATEDR	O	Combinational decode of TAP state machine used as multiplexed external scan cell clock.

A.6 Debug signals

Table A-8 lists the debug signals.

Table A-8 Debug signals

Name	I/O	Description
COMMRX	O	HIGH when comms channel receive buffer has data for processor to read.
COMMTX	O	Comms channel transmit. HIGH when comms channel transmit buffer is empty.
DBGACK	O	Debug acknowledge. HIGH when processor is in debug state.
DBGEN	I	Debug enable. Setting DBGEN enables external debug.
EDBGRQ	I	External debug request. Setting EDBGRQ puts processor in debug after current instruction.

A.7 Coprocessor signals

Table A-9 lists the coprocessor (CP) signals.

Table A-9 Coprocessor signals

Name	I/O	Description
ACANCELCP	O	Currently executing instruction ignored because of failed condition. Leave unconnected if CP interface unused.
AFLUSHCP	O	Cancel instructions in CP Execute, Decode, Issue, and Fetch stages. Leave unconnected if CP interface unused.
ASTOPCPD, ASTOPCPE	O	Hold CP pipeline in Decode stage. Driven from register after stalled stage. Hold CP pipeline in Execute stage. Driven from register after stalled stage. Leave both unconnected if CP interface unused.
CPBIGEND	O	Memory system is big-endian. When this signal is active, devices that support 64-bit data must assert CPLSSWP when loading or storing the 64-bit data for correct order when read/written. Leave unconnected if CP interface unused.
CPBOUNCEE1/2	I	Take undefined instruction trap for instruction in ARM Execute stage.
CPBUSYD1/2	I	Reserved for future expansion.
CPBUSYE1/2	I	Busy-waits the ARM Execute stage.
CPCLK	O	CP clock. Delayed version of GCLK . Must be provided at top-level clock distribution network. When VFP10 CP present, CPCLK is connected to VFPCLK . Leave unconnected if CP interface unused.
CPINSTR[25:0]	O	Instruction input from ARM10 processor. Valid at end of ARM Fetch stage. Validated by assertion of CPINSTRV. Leave unconnected if CP interface unused. Bits [27:26] always 11.
CPINSTRV	O	CP instruction on CPINSTR is valid new instruction. Leave unconnected if CP interface unused.
CPLSBUSY	CP to CP	Driven out of register on CP Issue/Decode boundary to signal other CPs that sender is doing a load or store multiple operation and is keeping control of STCMRCDATA bus.
CPLSDBL1/2	I	CP load/store request is for double word data.
CPLSLEN1/2[5:0]	I	Length of CP load/store transfer.
CPLSSWP1/2	I	Before writing, swap upper and lower data words on LDCMCRDATA or STCMRCDATA .

Table A-9 Coprocessor signals (continued)

Name	I/O	Description
CPRST	O	CP reset. Must be held for at least two cycles. Leave unconnected if CP interface unused.
CPSUPER	O	Supervisor mode. HIGH if ARM10 processor in Supervisor or interrupt-handling mode.
CPVALIDD	O	Valid CP instruction in ARM Decode stage.
LDCMCRDATA[63:0]	O	Carries data from ARM10 processor to CP. Leave unconnected if CP interface unused.
LSHOLDCPE	O	Hold CP pipeline in CP Execute stage. LSU stalled in ARM Execute stage. Leave unconnected if CP interface unused.
LSHOLDCPM	O	LSU stalled in ARM Memory stage in previous cycle. Leave unconnected if CP interface unused.
STCMRCDATA[63:0]	I	Carries data from CP to ARM10 processor.

A.8 Design for test signals

Table A-10 lists the DFT signals.

Table A-10 Design for test signals

Name	I/O	Description
A1020DFTCKEN	I	Enables the internal core GCLK .
A1020DFTRESET	I	Provides direct control over asynchronous resets in scan mode.
A1020DFTWCKEN	I	Enables the wrapper clock A1020WCLK to the dedicated test cells.
A1020RSTSAFE	I	Enables the reset to a portion of the core while testing external logic.
A1020SAFE	I	Forces safe values onto most core outputs. Used during core test.
A1020SCANEN	I	Scan enable for nonwrapper clock domains.
A1020SCANMODE	I	Puts the device into scan mode.
A1020SCANOUT[23:0]	O	Test bus input. Bits [15:0] required for cache upload or download. ATPG scan widths are user-configurable (24, 12, or 6).
A1020TEST	I	Enables cache upload or cache download and BIST test modes.
A1020TESTCFG[2:0]	I	Chooses which BIST or upload/download mode runs.
A1020WCLK	I	Wrapper clock for dedicated wrapper cells.
A1020WMUXINSEL	I	Selects core inputs (wrapper or external logic).
A1020WMUXOUTSEL	I	Selects core outputs (wrapper or external logic).
A1020WSCANEN	I	Scan enable for all wrapper cells.
WSCANIN[2:0]	I	Input ports for the wrapper scan chains.
A1020WSCANOUT[2:0]	O	Output ports for the wrapper scan chains.
HRESETN	I	Asynchronous reset.
SCANIN[23:0]	I	Test bus input. Bits [15:0] required for cache upload or download.
SCANMUX12	I	Setting both SCANMUX12 and SCANMUX6 enables access to 12 separate internal scan chains and 3 wrapper chains. Clearing both signals produces 24 separate internal scan chains and 3 wrapper chains.
SCANMUX6	I	Enables access to 6 separate internal scan chains and 1 wrapper chain.

Table A-10 Design for test signals (continued)

Name	I/O	Description
SCORETEST	I	Enables serial core test mode.
SFRESETN	I	Asynchronous reset.
UDLTEST	I	Enables the shared cells of the wrapper only. Must be enabled during 3-wrapper-chain mode, 12-chain mode and 24-chain mode.

A.9 ETM signals

Table A-11 lists the ETM10 signals.

Table A-11 ETM10 signals

Signal name	I/O	Description
ETMCORECTL[23:0]	O	Miscellaneous control signal inputs from the ARM10 processor.
ETMDA	O	The data address bus.
ETMDATA[63:0]	O	The load, store, and coprocessor data from the ARM10 processor.
ETMDATAVALID[1:0]	O	Valid signal for ETMDATA bus (one bit for each for HIGH and LOW word).
ETMIA	O	The instruction fetch address bus.
ETMR15BP[31:0]	O	The instruction address for branch phantom instructions.
ETMR15EX[31:0]	O	The instruction address for all nonbranch phantom instructions.
FIFOFULL	I	Indicates a request from the ETM10 for the ARM10 processor to stall execution to prevent the ETM10 from overflowing its FIFO.
ETMPWRDOWN	I	ETMPWRDOWN is asserted when the ETM10 is not enabled. The ARM10 processor uses the assertion ETMPWRDOWN to indicate that ETM10 bus inputs must be held stable.

Glossary

This glossary lists all the abbreviations used in the *ARM1020E Technical Reference Manual*.

Abort	An Abort is caused by an illegal memory access. Aborts can be caused by the external memory system or the MMU.
Access Permissions	The <i>Memory Management Unit</i> (MMU) determines the <i>Access Permissions</i> (AP) to regions of memory.
Advanced Microcontroller Bus Architecture	The ARM open standard for on-chip buses. AHB conforms to this standard.
AHB	<i>See</i> AMBA High-performance Bus.
ALU	<i>See</i> Arithmetic Logic Unit.
AMBA	<i>See</i> Advanced Microcontroller Bus Architecture.
AMBA High-performance Bus	The ARM processor interface to memory and peripherals.
AP	<i>See</i> Access Permissions.

Arithmetic logic unit	The component of the ARM processor that performs the arithmetic and logic operations.
ASIC	Application-Specific Integrated Circuit.
ATPG	Automated Test Pattern Generation.
Back of queue pointer	Points to the next location to read from when draining the write buffer <i>See also</i> Front of Queue pointer.
Big-endian	Memory organization in which the least significant byte of a word is at a higher address than the most significant byte. <i>See also</i> Little-endian.
BIU	<i>See</i> Bus Interface Unit.
BQ	<i>See</i> Back of Queue pointer <i>and also</i> Front of Queue pointer.
Branch folding	A branch can be predicted, pulled out of the normal instruction stream and effectively executed in parallel with the next instruction in the instruction stream.
Branch phantom	The condition codes of a predicted branch.
Breakpoint	If execution reaches this location, the debugger halts execution of the program. <i>See also</i> Watchpoint.
Bus interface unit	A <i>Bus Interface Unit</i> (BIU) that handles all data and/or instruction accesses across AHB.
C	Memory configuration Cachable <i>See also</i> NC, NCB <i>and</i> NCNB.
Cache hit	The instruction or data is found in the cache.
Cache miss	The instruction or data is not found in the cache.
CAM	<i>See</i> Content Addressable Memory.
Clock gating	Gating a clock signal for a macrocell with a control signal (such as PWRDOWN) and using the modified clock that results to control the operating state of the macrocell.
Condensed reference format	Condensed Reference Format. A vector file format proprietary to ARM Limited.
Content addressable memory	CAM includes comparison logic with each bit of storage. A data value is broadcast to all words of storage and compared with the values there. Words which match are flagged in some way. Subsequent operations can then work on flagged words. It is possible to read the flagged words out one at a time or write to certain bit positions in all of them.

Copy-back	<i>See</i> Write-back.
CPI	Clocks per instruction or cycles per instruction.
CPSR	Current program status register.
CPU core state	The state of: <ul style="list-style-type: none"> • all banked registers • the CPSR • the MMU TLB • the system control coprocessor, CP15 • the state of the debug coprocessor, CP14 • the VFP10 coprocessor • ETM10.
CRF	<i>See</i> Condensed Reference Format.
DA	Data address.
Data bus interface unit	The BIU that handles all data accesses across AHB.
Data physical address	The 32-bit address path between the DMMU and the DBIU.
Data streaming	The ability to return the second load miss data before a linefill completes.
Data transfer register	Two physically separate registers used to read and write to through the JTAG interface for debug.
DBIU	<i>See</i> Data bus interface unit.
DCache	Data cache (DCache) and associated write buffer. It has 1024 lines of 32 bytes arranged as a 64-way associative cache and uses virtual addresses from the ARM10E Integer Unit.
Debug ID Register	The DIDR contains details of implementer, architecture version and the number of watchpoints and breakpoints.
Debug Status and Control Register	The DSCR enables all debug functionality. It controls the debug modes, settings for catching ARM exceptions, and the comms channel.
DIDR	<i>See</i> Debug ID Register.
Dirty data	A data line that has been modified in the DCache and has not been written back to main memory.
DSCR	<i>See</i> Debug Status and Control Register.

DTR	<i>See</i> Data Transfer Register.
EmbeddedICE	The EmbeddedICE logic eases debugging in embedded systems. It contains watchpoint, control, and status registers.
Exception	An exception handles an event. For example, an external interrupt or an undefined instruction.
FAR	<i>See</i> Fault Address Register.
Fast context switch extension	The <i>Fast Context Switch Extension</i> (FCSE) enables cached processors with an MMU to present different addresses to the rest of the memory system for different software processes even when those processes are using identical addresses. <i>See also</i> MMU, MVA, PA, VA.
Fault	An abort that is generated by the MMU.
Fault Address Register	The FAR holds the virtual address of the access which was attempted when a fault occurred.
Fault Status Register	The FSR contains the source of the last fault. It indicates the domain and type of access being attempted when an abort occurred.
FCSE	<i>See</i> Fast context switch extension.
FIQ	Fast interrupt request. The exception for processing fast interrupts. <i>See also</i> IRQ.
FPGA	<i>See</i> Field programmable gate array.
Front of queue pointer	Points to the next entry to be written to in the write buffer. <i>See also</i> Back of queue pointer.
FSR	<i>See</i> Fault Status Register.
Gray code	Only one bit changes in the move from one state to the next state.
Halt mode	One of two mutually exclusive debug modes. In halt mode all processor execution halts when a breakpoint or watchpoint is encountered. All processor state, coprocessor state, memory and input/output locations can be examined and altered by the JTAG interface. <i>See also</i> Monitor mode.

Hit-under-miss	The HUM buffer enables program execution to continue even though there has been a data miss in the cache. If a load misses in the data cache, the outstanding request is moved into the HUM buffer. Other instructions, including loads, can continue to execute unless a second miss occurs or a dependency on the outstanding data is detected.
HUM	See Hit-under-miss.
IA	See Instruction address.
IBIU	See Instruction bus interface unit.
ICache	Instruction cache. It has 1024 lines of 32 bytes arranged as a 64-way associative cache. It uses virtual addresses from the ARM10E integer unit.
ICE	See In Circuit Emulator.
IDDQ	IDDQ refers to the quiescent current in CMOS integrated circuits. IDDQ is the IEEE symbol for the quiescent power supply current in a MOS circuit.
In-circuit emulator	A module for debugging in embedded systems.
Incoherence	When ICache or DCache copies of main memory and main memory get out of step with each other because one is updated and the other is not, the copies have become incoherent. See also memory coherence.
Instruction address	The 32-bit virtual address path between the ARM10E integer unit, IMMU, and ICache.
Instruction bus interface unit	The <i>Bus Interface Unit</i> (BIU) that handles all instruction accesses across AHB.
Instruction Transfer Register	The ITR sends instructions to the ARM10E processor during debug.
Instruction physical address	The MMU translates the modified virtual address to form the instruction physical address.
IPA	See Instruction Physical Address.
IRQ	Interrupt request. The exception for processing standard interrupts. See also FIQ, SWI.
ITR	See Instruction Transfer Register.
JTAG	Joint Test Action Group. The committee that defined the IEEE test access port and boundary-scan standard.

Leakage	The current each transistor takes even when it is not being switched.
Link register	Register 14 is the Link Register (LR). This register holds the address of the next instruction after a Branch and Link (BL) instruction, which is the instruction used to make a subroutine call.
Little-endian	Memory organization where the least significant byte of a word is at a lower address than the most significant byte. <i>See also</i> Big-endian.
Load/store unit	Part of the ARM10E integer unit that handles load/store transfers.
Lock-out layer	Isolates the CPU from the system bus, placing the CPU bus in the IDLE state.
LR	<i>See</i> Link Register.
LSU	<i>See</i> Load/Store Unit.
Memory coherency	Is the problem of ensuring that when a memory location is read either by a data read or an instruction fetch, the value actually obtained is always the value that was most recently written to the location. This can be difficult when there are multiple possible physical locations, such as main memory, a write buffer and/or cache. <i>See also</i> incoherence.
Memory management unit	An MMU controls caches and access permissions to blocks of memory, and translates virtual to physical addresses. The ARM processor has an IMMU for instructions and a DMMU for data. <i>See also</i> FCSE, MVA, TLB, PA, and VA,
Method of entry	In debug, bits [4:2] of the DSCR can be read to determine what caused an exception.
MMU	<i>See</i> Memory management unit.
Modified Virtual Address	Modified Virtual Address. A virtual address produced by the ARM10E integer unit can be changed by the current Process ID to provide a Modified Virtual Address(MVA) for the MMUs and caches. <i>See also</i> FCSE.
MOE	<i>See</i> Method of Entry.

Monitor mode	One of two mutually exclusive debug modes. In monitor mode the ARM processor enables a software abort handler provided by debug monitor or operating system debug task. When a breakpoint or watchpoint is encountered, this enables vital system interrupts to continue to be serviced while normal program execution is suspended. <i>See also</i> Halt mode.
Multi-ICE	Multi-ICE is a system for debugging embedded processor cores through a JTAG interface.
MVA	<i>See</i> Modified Virtual Address.
NB	Memory configuration, NonBufferable.
NC	Memory configuration, NonCachable.
NCB	Memory configuration, NonCachableBufferable.
NCNB	Memory configuration, NonCachableNonBufferable.
PA	Physical Address. The MMU performs a translation on Modified Virtual Addresses (MVA) to produce the Physical Address (PA) which is given to AHB to perform an external access. The PA is also stored in the DCache to avoid needing address translation when data is cast out of the cache. <i>See also</i> FCSE.
PA[7:0]	Physical Address (internal bus). The 8-bit data path between DMMU and DCache.
PC	Program Counter.
PDA	Personal Digital Assistant.
Penalty	the number of cycles in which no useful Execute stage pipeline activity can occur due to an instruction flow differing from that assumed or predicted. <i>See also</i> Branch folding, Branch phantom.
PLL	<i>See</i> Phase Locked Loop.
Phase Locked Loop	A clock synthesis device.
RDI	<i>See</i> Remote Debug Interface.
Remapping	Changing the address of physical memory or devices after the application has started executing. This is typically done to enable RAM to replace ROM once the initialization has been done.

Remote Debug Interface	Remote Debug Interface.
RISC	Reduced Instruction Set Computer.
ROM	Read Only Memory.
RTOS	Real Time Operating System.
Safe shift	values can be shifted from one scan cell to the next with no risk of error due to clock skew.
SDRAM	Synchronous Dynamic Random Access Memory.
SDT	Software Development Toolkit.
SP	Stack Pointer.
SPSR	Saved Program Status Register.
SRAM	Static Random Access Memory.
SWI	Software Interrupt. An instruction that causes the processor to call a programmer-specified subroutine.
TAP	Test Access Port.
TIC	Test Interface Controller.
TLB	Translation Look-aside Buffer. A cache of recently used page table entries that avoid the overhead of page-table-walking on every memory access. Part of the memory management unit.
TTBA	Translation Table Base Address. The starting point for the memory translation process. CP15 register r2 is the ARM1020E translation table base register.
Undefined	Indicates an instruction that generates an undefined instruction trap.
Unpredictable	Means the result of an instruction cannot be relied upon. Unpredictable instructions or results must not represent security holes. Unpredictable instructions must not halt or hang the processor, or any parts of the system.
VA	Virtual Address.

The MMU uses its page tables to translate a Virtual Address into a Physical Address. The processor executes code at the Virtual Address, which may be located elsewhere in physical memory.

See also FCSE, MVA, and PA.

Victim	The cache entry to be replaced.
Watchpoint	<p>A location in the program that is monitored. If the value stored there changes, the debugger halts execution of the program.</p> <p><i>See also</i> Breakpoint.</p>
Write-back	In a write-back cache, data is only written to main memory when it is forced out of the cache. Otherwise writes by the processor update only the cache. (Also known as copy-back).
Write buffer	Buffered writes can be written to memory by AHB while ARM10E continues reading instructions and data from ICache and DCache. ARM10E can also continue writing to DCache and the write buffer.
Write-through	In a write-through cache, data is written to main memory at the same time as the cache is updated.

Index

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